THE Σ^1_2 COUNTERPARTS TO STATEMENTS THAT ARE EQUIVALENT TO THE CONTINUUM HYPOTHESIS

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Abstract. We consider natural Σ^1_2 definable analogues of many of the classical statements that have been shown to be equivalent to CH. It is shown that these Σ^1_2 analogues are equivalent to that all reals are constructible. We also prove two partition relations for Σ^1_2 colourings which hold precisely when there is a non-constructible real.

§**1. Introduction.** In the mathematical literature, one finds a great number of statements that have been proved to be equivalent to the Continuum Hypothesis (CH). One such well-known equivalence is due to Sierpiński, and states that CH is equivalent to the statement that the plane $\mathbb{R}^2$ is the union of two sets $A, B \subseteq \mathbb{R}^2$ such that each horizontal section of A is countable, and each vertical section of B is countable. Another example is Davies' theorem, which states that CH is equivalent to that every function $f : \mathbb{R}^2 \to \mathbb{R}$ admits a representation

$$f(x, y) = \sum_{n=0}^{\infty} g_n(x) h_n(y),$$

where $g_n, h_n : \mathbb{R} \to \mathbb{R}$ are functions, and the sum above has only finitely many non-zero terms for every $(x, y) \in \mathbb{R}^2$.

In these types of theorems, usually the direct implication from CH is proved by a straight-forward inductive construction by well-ordering the reals in order type ω_1, and exploiting that each initial segment is countable. The result of the construction will usually be definable from the well-ordering. Perhaps it is no surprise then that if we work in Gödel's constructible universe L where there is a canonical choice of a well-ordering of $\mathbb{R}$, which moreover is Σ^1_2, then with some care it can be shown in many cases that there are Σ^1_2 definable witnesses to the direct implication.

On the other hand, the reverse implication often requires considerable ingenuity and does not at first seem to conform to a set pattern. In light of the above discussion about the situation in L, it is natural to ask what happens if we take a statement which implies CH, and replace it with a corresponding Σ^1_2 version. In [17] we considered the Σ^1_2 counterpart of Davies' theorem, and showed the following

"Σ_2^1 Davies' Theorem": All reals are constructible ($\mathbb{R} \subseteq L$) if and only if every Σ_2^1 function $f : \mathbb{R}^2 \to \mathbb{R}$ admits a representation

$$f(x, y) = \sum_{n=0}^{\infty} g(x, n)h(x, n),$$

where $g, h : \mathbb{R} \times \omega \to \mathbb{R}$ are Σ_2^1 functions, and the sum above has finitely many non-zero terms at each $(x, y) \in \mathbb{R}^2$.

It is natural to ask if this type of definable converse, which was found in the case of Davies' theorem, could hold for some of the many other statements that are equivalent to CH. However, the proof in [17] did not give a clear indication in this direction. In this paper we will prove that a number of the classical CH equivalents admit natural Σ_2^1 counterparts which turn out to be equivalent to that all reals are constructible. Specifically:

THEOREM 1.1. *The following statements are equivalent*:

(1) $\mathbb{R} \subseteq L$.
(2) *There are Σ_2^1 sets $A, B \subseteq \mathbb{R}^2$ such that $A \cup B = \mathbb{R}^2$, and all the sections $A_x = \{y \in \mathbb{R} : (x, y) \in A\}$ and $B^y = \{x \in \mathbb{R} : (x, y) \in B\}$ are countable.*
(3) *There are Σ_2^1 sets $A_1, A_2, A_3 \subseteq \mathbb{R}^3$ such that $A_1 \cup A_2 \cup A_3 = \mathbb{R}^3$, and every line l in the direction of the x_i-axis meets A_i in finitely many points.*
(4) *There are uncountable Σ_2^1 sets A_0 and A_1 such that $A_0 \cup A_1 = \mathbb{R}$ and for all $a \in \mathbb{R}$ the set $(a + A_0) \cap A_1$ is countable.*
(5) *The plane can be covered by three Σ_2^1 clouds[1] with centres in L.*
(6) *There is a Σ_2^1 surjection $f : \mathbb{R} \to \mathbb{R}^2 : x \mapsto (f_1(x), f_2(x))$ such that either $f_1'(x)$ or $f_2'(x)$ exists for all $x \in \mathbb{R}$.*

Here (2) and (3) correspond to CH equivalences proven by Sierpiński [14]; (4) to an equivalence due to Banach and Trzeciakiewicz, [1, 18]; (5) to an equivalence due to Komjath [6]; and (6) to an equivalence proven by Morayne [11].

The proofs of the above equivalences also offer an explanation for why and when a classical CH equivalence admits a Σ_2^1 counterpart. The reason that the above Σ_2^1 translations work can be found in the structure of the proofs of the corresponding classical CH equivalences. Though it is not always immediately clear from the literature, there is a common underlying structure of the proofs of CH from the given statement, and in fact of the statements themselves. Roughly speaking, the structure is as follows: The *statements* are of the form that there exists certain sets (or n-ary relations) $R_1, R_2, \ldots$ and functions $f_1, f_2, \ldots$ which satisfy some finiteness or countability requirement, and that *all* reals must satisfy some relations that are expressed in terms of the given sets and functions. The *proof* that such a statement implies CH then can be cast in the following general form: One fixes a set of reals of size $\aleph_1$, and forms a "hull" of reals that satisfies the relevant relations with this fixed set of reals. The countability condition on the sets and functions $R_1, R_2, \ldots$, $f_1, f_2, \ldots$ then implies that this "hull" must have size $\aleph_1$. The statement is then seen to imply that in fact *all* reals are in this hull, hence $2^{\aleph_0} = \aleph_1$.

In practice, one more often argues indirectly by assuming $\neg$CH, and then use this to produce a real which is "transcendental" in the sense that it fails to satisfy the

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[1]See §3.3 for the definition of clouds.

prescribed relations. In the Σ_2^1 translations we consider, this corresponds to assuming that there is a non-constructible real. In our proofs the finiteness/countability conditions are then used, in conjunction with the Mansfield–Solovay perfect set theorem (see Theorem 2.3), to prove that the constructible reals are indeed a suitable "hull". Another important tool is the Shoenfield absoluteness theorem (see [13] or [2, 25.20]), which allows us to work in a model of the form $L[x]$, $x \in \mathbb{R}$, which for the purpose of counting arguments can then be assumed to satisfy $\aleph_1^{L[x]} = \aleph_1^L$, see Lemma 2.2.

Using the same ideas we also prove the following two partition relations which hold for Σ_2^1 colourings precisely when there are non-constructible reals.

THEOREM 1.2. *The following are equivalent*:

(1) $\mathbb{R} \not\subseteq L$.
(2) *For every Σ_2^1 function $f : \mathbb{R} \times \mathbb{R} \to \omega$ there are sets $C, D \subseteq \mathbb{R}$ such that $|C| = |D| = \aleph_0$ and $f \restriction C \times D$ is monochromatic.*
(3) *For every Σ_2^1 colouring $g : \mathbb{R} \to \omega$ there are four distinct $x_{00}, x_{01}, x_{10}, x_{11} \in \mathbb{R}$ of the same colour such that*

$$x_{00} + x_{11} = x_{01} + x_{10}.$$

This theorem, as well as Theorem 1.1, naturally relativizes to $L[a]$ and $\Sigma_2^1(a)$, where $a \in \mathbb{R}$ is a parameter.

§2. **Definitions and preliminaries.** In this section, we collect various general definitions and preliminary observations that are needed in our proofs. For this purpose, it is immensely practical to follow the (effective) descriptive set-theoretic convention and use $\mathbb{R}$ to stand for *any* recursively presented uncountable Polish space (which is warranted since all such spaces are isomorphic by a Δ_1^1 bijection, see [12]). This convention will, however, cause problems later, where $\mathbb{R}$ will need to stand for the actual (linearly ordered field of) real numbers. Henceforth, we will use $\mathscr{R}$ to denote the descriptive set-theoretic *reals* and $\mathbb{R}$ for the actual real line.

We shall assume that the reader is familiar with the basic elements of (effective) descriptive set theory, as found in e.g., [2, 3, 9] or [12], though we briefly review the most important notions below. Our notation is, for the most part, in line with that of [12], and in particular, recursively presented Polish spaces are denoted with script letters $\mathscr{X}, \mathscr{Y}, \mathscr{Z}, \ldots$.

2.1. Σ_2^1 sets and functions. In this paper, a Σ_2^1 set is a set that can be defined by a Σ_2^1 predicate, a Π_2^1 set is a set that is the complement of a Σ_2^1 set, and a Δ_2^1 set is a set that is both Σ_2^1 and Π_2^1. We denote by $\Sigma_2^1(a)$, $\Pi_2^1(a)$, and $\Delta_2^1(a)$ the corresponding relativized pointclasses, where a is some real (i.e., $a \in \mathscr{R}$).

In this paper, we will say that a (total) function $f : \mathscr{X} \to \mathscr{Y}$ is Σ_2^1 (or, more generally, $\Sigma_2^1(a)$) if the graph of f is a Σ_2^1 ($\Sigma_2^1(a)$) subset of $\mathscr{X} \times \mathscr{Y}$. If a function has a Σ_2^1 graph then in fact it is a Δ_2^1 graph since if $\psi(x, y)$ is a Σ_2^1 predicate defining (the graph of) f then

$$f(x) = y \iff (\forall z)(\neg \psi(x, z) \vee z = y),$$

which shows that f has a Π_2^1 definition as well. We will say that a Σ_2^1 predicate $\psi(x, y)$ *defines a function* if there is a total function $f : \mathscr{X} \to \mathscr{Y}$ such that

$f(x) = y \iff \psi(x, y)$. The reader should be warned that this notion is sensitive to the model of set theory in which we work, since a predicate which defines a total function in one model may only define a partial function in another (for example, take $\psi(x, y)$ to be a Σ_2^1 predicate which says that $x = y$ and $x \in L$). Note, however, that

$$(\forall x)(\forall y, y')\psi(x, y) \wedge \psi(x, y') \implies y = y' \qquad (2.1)$$

is Π_2^1 and therefore absolute, and so if (2.1) is satisfied in one model, it is satisfied in all. In other words, a Σ_2^1 predicate which defines a partial function will do so in any model, but it may fail to define a total function in all models even if it does so in one.

2.2. Coding the L_α. Our notation follows that of [3, p. 167ff.], with very few differences. For convenience we recall the definitions and facts that are most important for the present paper.

The canonical well-ordering of L will be denoted $<_L$. The language of set theory (LOST) is denoted $\mathcal{L}_\varepsilon$. If $x \in 2^\omega$ then we define a binary relation on ω by

$$m \, \epsilon_x \, n \iff x(\langle m, n \rangle) = 1,$$

where $\langle \cdot, \cdot \rangle$ refers to some (fixed) standard Gödel pairing function coding a pair of integers by a single integer. We let

$$M_x = (\omega, \varepsilon_x)$$

be the $\mathcal{L}_\varepsilon$ structure coded by x. If M_x is well-founded and extensional then we denote by $\operatorname{tr}(M_x)$ the transitive collapse of M_x, and by $\pi_x : M_x \to \operatorname{tr}(M_x)$ the corresponding isomorphism.

The following proposition encapsulates the basic descriptive set-theoretic correspondences between x, M_x, and the satisfaction relation. We refer to [3, 13.8] and the remarks immediately thereafter for a proof.

PROPOSITION 2.1. (a) *If $\varphi(v_0, \ldots, v_{k-1})$ is a LOST formula with all free variables shown then*

$$\{(x, n_0 \ldots, n_{k-1}) \in 2^\omega \times \omega \times \cdots \times \omega : M_x \models \varphi[n_0, \ldots, n_{k-1}]\}.$$

is arithmetical.

(b) *For $x \in 2^\omega$ such that M_x is well-founded and extensional, the relation*

$$\{(m, y) \in \omega \times \mathscr{R} : \pi_x(m) = y\}$$

is arithmetical in x.

(c) *There is a LOST sentence σ_0 such that if $M_x \models \sigma_0$ and M_x is well-founded and extensional, then $M_x \simeq L_\delta$ for some limit ordinal $\delta < \omega_1$.*

(d) *There is a LOST formula $\varphi_0(v_0, v_1)$ which defines the restriction of the canonical well-ordering $<_L$ of L to L_δ for any limit $\delta > \omega$.*

Define as in [3, p. 170] the restriction $M_x \restriction k$, for $x \in 2^\omega$ and $k \in \omega$, to be the $\mathcal{L}_\varepsilon$ structure

$$M_x \restriction k = (\{n : n \, \varepsilon_x \, k\}, \varepsilon_x).$$

For $\mathscr{X}$ a recursively presented Polish space, let $R^{\mathscr{X}} \subseteq \mathscr{X} \times 2^\omega$ be defined by

$$R^{\mathscr{X}}(y, x) \iff M_x \text{ is well-founded, extensional, and } M_x \models \sigma_0 \wedge (\exists n)y = \pi_x(n) \wedge$$
$$(\forall z <_L x)(M_z \text{ is well-founded, extensional, and } M_z \models \sigma_0) \implies (\forall k)\pi_z(k) \neq y$$

In other words, $R^{\mathscr{X}}(y, x)$ holds iff x is the least code for an L_α, α a limit, such that $y \in L_\alpha$. The relation $R^{\mathscr{X}}$ is Δ_2^1.

2.3. Coding initial segments. Let $\prec$ denote $<_L \upharpoonright \mathscr{R}$, the canonical well-ordering of $\mathscr{R}$ in L. This is a Δ_2^1 well-ordering which has a good coding of initial segments. More precisely, $\prec$ is a *strongly Δ_2^1 well-ordering*, which means that $\prec$ has length ω_1 and $\text{IS} \subseteq \mathscr{R} \times \mathscr{R}^{\leq \omega}$ defined by

$$\text{IS}(x, v) \iff (\forall n)v(n) \prec x \wedge (\forall z \prec x)(\exists n)v(n) = z \wedge (\forall i, j)i = j \vee v(i) \neq v(j)$$

is Δ_2^1. The point is that quantifications over an initial segment of $\prec$ can be replaced by a quantifier over ω in hierarchy calculations, see [12, 5A.1] for details. We also define a function $\text{IS}^* : \mathscr{R} \to \mathscr{R}^{\leq \omega}$ and a partial function $\text{IS}^\# : \mathscr{R} \times \mathscr{R} \to \omega$ by

$$\text{IS}^*(x) = v \iff \text{IS}(x, v) \wedge (\forall w \prec v)\neg \text{IS}(x, w),$$
$$\text{IS}^\#(x, y) = n \iff \text{IS}^*(x)(n) = y.$$

These are Σ_2^1.

2.4. The size of $L \cap \mathbb{R}$. There are several counting arguments below that rely on having some information about the cardinality of sets of reals in L. The following simple observations are extremely useful for this purpose:

LEMMA 2.2. (1) *If there is a non-constructible real in V, then there is a non-constructible real $x \in V$ such that $\aleph_1^{L[x]} = \aleph_1^L$.*

(2) *Suppose ψ is a $\Sigma_2^1(a)$ predicate defining the set A, where $a \in L$. Then if A is uncountable, then $A \cap L$ is uncountable in L.*

PROOF. (1) If $\aleph_1^V = \aleph_1^L$, then any non-constructible $x \in V$ will do. If $\aleph_1^L$ is countable in V, then there must be a real $x \in V$ which is Cohen over L. For any such x it holds that $\aleph_1^{L[x]} = \aleph_1^L$.

(2) If $A \cap L$ is countable in L then there is some $v : \mathbb{N} \to \mathscr{R}$ in L such that

$$(\forall x)(\psi(x) \longrightarrow (\exists n)v(n) = x)$$

holds. Since this is $\Pi_2^1(a, v)$ it is absolute, and so A is countable. $\dashv$

The typical application of (1) above will be that if we know that some statement which is downwards absolute holds in V, and $\mathbb{R} \not\subseteq L$, then the statement holds in some $L[x]$ where $x \notin L$, and the constructible reals have cardinality $\aleph_1$ in $L[x]$.

Finally, we recall the perfect set theorem for Σ_2^1 sets by Mansfield and Solovay which will be used often:

THEOREM 2.3 (Mansfield [8], Solovay [15]). *Let A be a $\Sigma_2^1(a)$ set. Then either $A \subseteq L[a]$, or else A contains a perfect set. In particular, if a Σ_2^1 set contains a non-constructible real then it is uncountable.*

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§3. **Results.**

3.1. Sierpiński's equivalences. In this section we consider the Σ_2^1 counterparts of two of Sierpiński's classical CH equivalences (see e.g., [14]). The first is the counterpart to: CH is equivalent to the existence of two sets $A, B \subseteq \mathbb{R}^2$ with $A \cup B = \mathbb{R}^2$ such that all vertical sections of A are countable and all horizontal sections of B are countable.

We include a version of this that is stated in terms of covering the plane by graphs of countably many functions, since this is needed later in Section 3.4 below.

THEOREM 3.1. *The following are equivalent*:

(1) $\mathbb{R} \subseteq L$.
(2) *There is a Σ_2^1 linear order $<$ of $\mathbb{R}$ such that for all $x \in \mathbb{R}$ the initial segment $\{y \in \mathbb{R} : y < x\}$ is countable.*
(3) *There are Σ_2^1 sets $A, B \subseteq \mathbb{R}^2$ such that $A \cup B = \mathbb{R}^2$, and all the sections $A_x = \{y \in \mathbb{R} : (x, y) \in A\}$ and $B^y = \{x \in \mathbb{R} : (x, y) \in B\}$ are countable.*
(4) *There are Σ_2^1 functions $F_A : \mathbb{R} \times \omega \to \mathbb{R}$ and $F_B : \mathbb{R} \times \omega \to \mathbb{R}$ such that $A = \{(x, F_A(x, n)) : x \in \mathbb{R}, n \in \omega\}$ and $B = \{(F_B(y, n), y) : y \in \mathbb{R}, n \in \omega\}$ satisfy (3).*

PROOF. (1) $\implies$ (4). Let z be the $\prec$-least element with an infinite initial segment. Let

$$F_A(x, n) = \mathrm{IS}^*(\max_\prec(x, z))(n)$$

and

$$F_B(x, n) = \begin{cases} x & \text{if } n = 0, \\ \mathrm{IS}^*(\max_\prec(x, z))(n - 1) & \text{if } n > 0 \end{cases}$$

where $\max_\prec(x, z)$ is the larger of x and z in $\prec$.

(4) $\implies$ (3) is clear. For (3) $\implies$ (1), suppose that there is $x_0 \in \mathbb{R} \setminus L$ but that (3) holds. By Lemma 2.2 we may assume that $\aleph_1^{L[x_0]} = \aleph_1^L$ and that $V = L[x_0]$, since if (3) holds it holds in $L[x_0]$. Since the section A_{x_0} is countable we can find $y \in (\mathbb{R} \cap L) \setminus A_{x_0}$, and so $(x_0, y) \in B$ since $A \cup B = \mathbb{R}^2$. But this means that B^y, which is a $\Sigma_2^1(y)$ set, contains a non-constructible real (namely x_0), and so since $y \in L$ it follows by the perfect set theorem (Theorem 2.3) that it must be uncountable, a contradiction.

Finally, (1) $\implies$ (2) is clear, since the canonical well-ordering of $\mathbb{R}$ in L satisfies (2), and (2) $\implies$ (3) follows since defining $A = \{(x, y) \in \mathbb{R}^2 : y < x\}$ and $B = \{(x, y) : x \leq y\}$ clearly works. ⊣

Next we consider the Σ_2^1 counterpart to the following CH equivalence due to Sierpiński (see [14]): CH holds iff there are sets $A_1, A_2, A_3 \subseteq \mathbb{R}^3$ such that $A_1 \cup A_2 \cup A_3 = \mathbb{R}^3$, and every line l in the direction of the x_i-axis meets A_i in finitely many points.

THEOREM 3.2. *All reals are constructible if and only if there are Σ_2^1 sets $A_1, A_2, A_3 \subseteq \mathbb{R}^3$ such that $A_1 \cup A_2 \cup A_3 = \mathbb{R}^3$, and every line l in the direction of the x_i-axis meets A_i in finitely many points.*

PROOF. Suppose all reals are constructible. Define, for $i = 1, 2, 3$, the set $\tilde{A}_i$ by

$$(x_1, x_2, x_3) \in \tilde{A}_i \iff \text{if } x_j = \max_\prec\{x_1, x_2, x_3\} \text{ then } x_j \neq x_i,$$
$$\text{and if } k \neq i, j \text{ then } \text{IS}^\#(x_j, x_i) < \text{IS}^\#(x_j, x_k)$$
$$\iff (\forall j)(x_j = \max_\prec\{x_1, x_2, x_3\} \implies (x_j \neq x_i \wedge$$
$$((\forall k \leq 3)(k \neq i \wedge k \neq j) \implies \text{IS}^\#(x_j, x_i) < \text{IS}^\#(x_j, x_k))).$$

Clearly $\tilde{A}_i$ is Δ_2^1. Let $A_i = \tilde{A}_i \cup \{(x, y, z) : x = y = z\}$. Then $\mathbb{R}^3 = A_1 \cup A_2 \cup A_3$. If l is a line parallel to an axis, say $l = \{(x, b, c) : x \in \mathbb{R}\}$, then by definition there are only finitely many x such that $(x, b, c) \in A_1$.

For the converse, suppose there is $x_0 \in \mathbb{R} \setminus L$. As before, we may assume that $V = L[x_0]$ and that $\aleph_1^{L[x_0]} = \aleph_1^L$. If $(u, v) \in \mathbb{R}^2 \cap L$, then the line $\{(u, v, x) : x \in \mathbb{R}\} \cap A_3$ is a finite $\Sigma_2^1(u, v)$ set, and so by Theorem 2.3 it does not contain a non-constructible real. Thus $(u, v, x_0) \notin A_3$ for all $u, v \in \mathbb{R} \cap L$. For any $u \in \mathbb{Q}$ the set $\{(u, x, x_0) : x \in \mathbb{R}\} \cap A_2$ is finite, and so since $\aleph_1^L = \aleph_1$ there must be some $x_1 \in \mathbb{R} \cap L$ such that $(u, x_1, x_0) \notin A_2$ for all $u \in \mathbb{Q}$. Since $A_1 \cap \{(u, x_1, x_0) : u \in \mathbb{R}\}$ is finite, it follows that there is $x_2 \in \mathbb{Q}$ such that $(x_2, x_1, x_0) \notin A_1 \cup A_2 \cup A_3$.   ⊣

3.2. Banach–Trzeciakiewicz's equivalence. [1] and [18] contain the following equivalence: CH holds if and only if there are uncountable sets $A_0, A_1 \subseteq \mathbb{R}$ such that $A_0 \cup A_1 = \mathbb{R}$ and for each $a \in \mathbb{R}$ the set $(a + A_0) \cap A_1$ is countable. We have the following Σ_2^1 counterpart:

THEOREM 3.3. *All reals are constructible if and only if there are* uncountable Σ_2^1 *sets* A_0 *and* A_1 *such that* $A_0 \cup A_1 = \mathbb{R}$ *and for all* $a \in \mathbb{R}$ *the set* $(a + A_0) \cap A_1$ *is countable.*

PROOF. If $\mathbb{R} \subseteq L$, it is easy to see that there is a Δ_2^1 Hamel basis $H \subseteq \mathbb{R}$ for $\mathbb{R}$. (In fact, by [10] there even is a Π_1^1 Hamel basis for $\mathbb{R}$.) Define a function $f : \mathbb{R} \to \mathbb{R}^{<\omega}$ by

$$f(x) = (x_1, \ldots, x_n)$$
$$\iff x_1, \ldots, x_n \in H \wedge x_1 \prec \cdots \prec x_n \wedge (\exists q_1, \ldots, q_n \in \mathbb{Q} \setminus \{0\}) x = \sum_{i=1}^n q_i x_i.$$

Clearly f is Δ_2^1. Write $H = H_0 \cup H_1$, where H_0 and H_1 are disjoint uncountable Δ_2^1 sets, and define

$$x \in A_i \iff (\exists (x_1, \ldots, x_n) \in \mathbb{R}^{<\omega}) f(x) = (x_1, \ldots, x_n) \wedge x_n \in H_i.$$

Then A_i is Σ_2^1 (in fact, A_i is Δ_2^1) and $A_0 \cup A_1 = \mathbb{R}$. Fix $a \in \mathbb{R}$, and note that if $\max f(a) \prec \max f(x)$ and $x \in A_0$ then $a + x \in A_0$. Thus $(a + A_0) \cap A_1$ is countable since $\{x \in \mathbb{R} : \max f(x) \preceq \max f(a)\}$ is.

For the converse, suppose that there is a non-constructible real $x_0 \in \mathbb{R} \setminus L$. By Lemma 2.2.(1) we may assume that $\aleph_1^{L[x_0]} = \aleph_1^L$. From this and Lemma 2.2.(2) it follows that $A_0 \cap L$ and $A_1 \cap L$ are uncountable, and so $A_0 \cap L[x_0]$ and $A_1 \cap L[x_0]$ are uncountable in $L[x_0]$. By assumption, for each $a \in A_0$ we either have $a + x_0 \in A_0$ or $a + x_0 \in A_1$. If the latter held for uncountably many $a \in A_0 \cap L$ then $(x_0 + A_0) \cap A_1$ would be uncountable, contrary to our assumption. Thus we can find $a \in L \cap A_0$ such that $a + x_0 \in A_0$. Similarly, there is $b \in A_1 \cap L$ such that $b + x_0 \in A_1$. But since $b + x_0 = a + x_0 + (b - a)$ we now have that $b + x_0 \in ((b - a) + A_0) \cap A_1$,

and so this set, which is $\Sigma_2^1(b - a)$, contains the non-constructible real $b + x_0$, and so is uncountable.   ⊣

3.3. Komjath's clouds. A *cloud* in $\mathbb{R}^2$ is a set $A \subseteq \mathbb{R}^2$ such that for some point $\vec{x} \in \mathbb{R}^2$ (called a *centre* of A) it holds that each infinite ray from $\vec{x}$ meets A in at most finitely many points. In [6] the following was shown:

THEOREM 3.4 (Komjath). *CH is equivalent to that the plane can be covered by three clouds.*

THEOREM 3.5. $\mathbb{R} \subseteq L$ *is equivalent to that the plane can be covered by three* Σ_2^1 *clouds with centres in* L.

PROOF. Assume that $\mathbb{R} \subseteq L$. We will give Σ_2^1 definitions of clouds A_0, A_1, and A_2 centered at $a_0 = (0, 1)$, $a_1 = (1, 0)$ and $a_2 = (0, 0)$, respectively, such that $\mathbb{R} = A_0 \cup A_1 \cup A_2$. For $y \in \mathbb{R}^2 \setminus \{a_0, a_1, a_2\}$ let $\overline{a_i y}$ denote the infinite ray starting at a_i extending through y. Let $\mathscr{E}$ be the set of all infinite rays from a_0, a_1 or a_2. The set of $\mathscr{E}$ can be identified with the union of the three disjoint circles centered at a_0, a_1 and a_2, and so $\mathscr{E}$ is a recursively presented Polish space in a natural way. Let $\mathscr{E}_\alpha = \mathscr{E} \cap L_\alpha$.

We define the set $A_i' \subseteq \mathbb{R}^2 \times 2^\omega$ as follows: $(y, x) \in A_i'$ if and only if

(1) $R^{\mathscr{E}}(\overline{a_i y}, x)$, i.e., x is $\prec$-least such that $M_x \simeq L_\alpha$ for the smallest limit $\alpha > \omega$ such that $\overline{a_i y} \in L_\alpha$.

(2) If $(j_l)_{l \in \omega}$ is a strictly increasing sequence enumerating the set

$$\{j \in \omega : \pi_x(j) \in \mathscr{E}_\alpha \setminus \bigcup \{\mathscr{E}_\delta : \delta < \alpha, \delta \text{ a limit}\}\}$$

and the ray $\overline{a_i y}$ is $\pi_x(j_l)$, then y is a point of intersection of $\pi_x(j_l)$ and one of the rays $\pi_x(j_0), \ldots, \pi_x(j_{l-1})$ or $\overline{a_j a_k}$, $j \neq k$ and $j, k \neq i$.

Then A_i' is Σ_2^1 since (2) can (given that (1) holds) be expressed by saying (where $j, k \neq i$)

$(\exists l)[\pi_x(l) = \overline{a_i y} \wedge ((y \in \overline{a_i y} \cap \overline{a_{j_0} a_{j_1}}) \vee ((\exists j < k) R^{\mathscr{E}}(\pi_x(j), x) \wedge y \in \pi_x(j) \cap \overline{a_i y}))].$

Let $A_i = \{y \in \mathbb{R}^2 : (\exists x) A_i'(y, x)\}$, which clearly is a Σ_2^1 set, and note that if $y \in \mathbb{R}^2 \cap L$ then there must be some $i \in \{0, 1, 2\}$ such that $y \in A_i$, and so $A_0 \cup A_1 \cup A_2 = \mathbb{R}^2 \cap L$, as required.

For the converse, assume that there are Σ_2^1 clouds A_0, A_1, and A_2 with centres in L covering the plane. After possibly applying an affine transformation (defined in L), we may assume that A_0, A_1, and A_2 are centered at $(0, 1), (1, 0)$, and $(0, 0)$, respectively.

By the usual arguments, we can assume that $V = L[r]$ for some $r \in \mathbb{R} \setminus L$ and that $\aleph_1 = \aleph_1^L$. Define an equivalence relation in $(0, \frac{\pi}{4})$ by

$$\alpha \sim \alpha' \iff \frac{1 - \tan(\alpha')}{1 - \tan(\alpha)} \in \mathbb{Q}_+.$$

Then $\sim$ has countable classes and $\alpha \in L$ iff $[\alpha]_\sim \subseteq L$.

For $\alpha, \beta \in (0, \frac{\pi}{4})$, let l_α denote the straight line in the plane given by the equation $\tan(\alpha)x + y = 1$, and t_β be the line given by $x + \tan(\beta)y = 1$. Note that the intersection point (x, y) of l_α and t_β satisfies $\frac{y}{x} = \frac{1 - \tan(\alpha)}{1 - \tan(\beta)}$.

Consider $\alpha \in (0, \frac{\pi}{4}) \cap L$. Since $l_\alpha \cap A_0$ is a finite $\Sigma_2^1(\alpha)$ set, it cannot contain any non-constructible points (by Theorem 2.3, for example). Thus if $\beta \in (0, \frac{\pi}{4}) \setminus L$, then the intersection of l_α and t_β cannot be in A_0. So fix $\beta_0 \in (0, \frac{\pi}{4}) \setminus L$. Since $A_1 \cap \{t_\beta : \beta \in [\beta_0]_\sim\}$ is countable there must be some $\alpha_0 \in (0, \frac{\pi}{4}) \cap L$ such that $t_\beta \cap t_\alpha \nsubseteq A_1$ for all $\beta \in [\beta_0]_\sim$ and $\alpha \in [\alpha_0]_\sim$, whence $t_\beta \cap t_\alpha \subseteq A_2$ for such α and β. For $n \in \mathbb{N}$, choose $\alpha_n \in [\alpha_0]_\sim$ and $\beta_n \in [\beta_0]_\sim$ such that

$$\frac{1 - \tan(\alpha_0)}{1 - \tan(\alpha_n)} = n = \frac{1 - \tan(\beta_0)}{1 - \tan(\beta_n)}.$$

Then for all $n \in \omega$ the intersection point $(x_n, y_n) \in l_{\alpha_n} \cap t_{\beta_n}$ satisfies

$$\frac{y_n}{x_n} = \frac{1 - \tan(\alpha_n)}{1 - \tan(\beta_n)} = \frac{1 - \tan(\alpha_0)}{1 - \tan(\beta_0)}.$$

and so they are all on the same line through $(0, 0)$, and since $(x_n, y_n) \in A_2$ for all $n \in \mathbb{N}$ this contradicts that each ray from $(0, 0)$ meets A_2 in finitely many points. $\dashv$

REMARK 3.6. It is interesting to note that in the previous proof, the assumption that A_1 and A_2 are Σ_2^1 were never used. Thus we have:

COROLLARY 3.7. $\mathbb{R} \subseteq L$ is equivalent to that the plane can be covered by three clouds with centres in L, one of which is Σ_2^1.

3.4. Differentiable functions after Morayne. A *Peano function* is a surjection $f : \mathbb{R} \to \mathbb{R} \times \mathbb{R}$. In [11], Morayne proved that CH is equivalent to the existence of a Peano function $f(x) = (f_1(x), f_2(x))$ such that at every $x \in \mathbb{R}$ at least one of the derivatives $f_1'(x)$ or $f_2'(x)$ exists. We obtain the following corresponding Σ_2^1 version:

THEOREM 3.8. *The following are equivalent*:

(1) *All reals are constructible*
(2) *There is a Σ_2^1 surjection $f : \mathbb{R} \to \mathbb{R}^2 : x \mapsto (f_1(x), f_2(x))$ such that either $f_1'(x)$ or $f_2'(x)$ exists for all $x \in \mathbb{R}$.*

PROOF OF (1) $\implies$ (2). We will show that the construction from CH due to Morayne translates to the Σ_2^1 setting.

For this, first define $f_1(t) = t \sin(t)$ on $t \in (-\infty, 1) = I_1$ and $f_2(t) = t \sin(t)$ on $t \in (-1, \infty) = I_2$. The sets

$$C^i = \{(r, t) \in \mathbb{R} \times I_i : f_i(t) = r\},$$

$i = 1, 2$, are Δ_1^1 and for each $r \in \mathbb{R}$ the section $C_r^i = \{t \in I_i : (r, t) \in C^i\}$ is countably infinite. It follows from (the effective version of) the Lusin–Novikov Theorem [5, 18.10] that there are Δ_1^1 functions $g_i : \mathbb{R} \to I_i^\omega$ such that $g_i(r)$ enumerates C_r^i injectively. Now let F_A and F_B be the functions from Theorem 3.1.(4), and define for $t \in \mathbb{R} \setminus I_1$

$$f_1(t) = y \iff (\exists r \in \mathbb{R})(\exists n \in \omega) f_2(t) = r \wedge g_2(r)(n) = t \wedge F_B(r, n) = y$$

and for $t \in \mathbb{R} \setminus I_2$

$$f_2(t) = y \iff (\exists r \in \mathbb{R})(\exists n \in \omega) f_1(t) = r \wedge g_1(r)(n) = t \wedge F_A(r, n) = y.$$

Note that whenever $t \notin I_1$ and $f_2(t)$ assumes the value r for the nth time as enumerated by $g_2(r)$, then $(f_1(t), f_2(t)) = (F_B(r,n), r)$, and so the graph of $t \mapsto (f_1(t), f_2(t))$ covers $B = \{(F_B(r,n), r) : r \in \mathbb{R}, n \in \omega\}$ as t ranges in I_1. Similarly, the graph of $t \mapsto (f_1(t), f_2(t))$ covers $A = \{(r, F_A(r,n)) : r \in \mathbb{R}, n \in \omega\}$ as t ranges over I_2. Thus $t \mapsto (f_1(t), f_2(t))$ is a Σ_2^1 Peano function with f_1 differentiable on I_1 and f_2 differentiable on I_2. $\dashv$

The proof of (2) $\implies$ (1) in Theorem 3.8 requires several lemmata. We start with a general observation about open Π_2^1 sets. Recall that the class of Π_2^1 sets is ω-*parametrized*, meaning that for any recursively presented Polish $\mathscr{X}$, there is a Π_2^1 set $P^{(\mathscr{X})} \subseteq \omega \times \mathscr{X}$ such that

$$P_n^{(\mathscr{X})} = \{x \in \mathscr{X} : (n,x) \in P\}$$

enumerates the Π_2^1 sets in $\mathscr{X}$. In particular, there is such a set $P^{(\omega)} \subseteq \omega \times \omega$ parametrizing the Π_2^1 subsets of ω. We let

$$\mathfrak{a} = \{\langle n,m\rangle : (n,m) \in P^{(\omega)}\},$$

where $\langle \cdot, \cdot \rangle$ is some standard Gödel pairing function. Note that $\mathfrak{a} \in L$.

LEMMA 3.9. *Suppose $A \subseteq \mathscr{X}$ is an open Π_2^1 set. Then there is a $\Sigma_1^0(\mathfrak{a})$ predicate $\psi(x)$ such that $x \in A \iff \psi(x)$.*

PROOF. Let d be a compatible metric on $\mathscr{X}$ and let $(x_n)_{n\in\omega}$ be a dense sequence in $\mathscr{X}$ such that $(d, (x_n)_{n\in\omega})$ is a recursive presentation of $\mathscr{X}$. Let $(q_m)_{m\in\omega}$ enumerate (effectively) the positive rationals, and define

$$a = \{\langle n,m\rangle \in \omega : (\forall x)d(x,x_n) < q_m \implies x \in A\}.$$

Then the set $a \subseteq \omega$ is Π_2^1, and

$$x \in A \iff (\exists n,m)\langle n,m\rangle \in a \wedge d(x,x_n) < q_m$$

which gives a $\Sigma_1^0(a)$ definition of A, whence A is $\Sigma_1^0(\mathfrak{a})$. $\dashv$

LEMMA 3.10. *Let $\psi(x,y)$ be a Δ_2^1 predicate which defines a function $f : \mathbb{R} \to \mathbb{R}$. Then*:

(1) *There is a $\Pi_2^0(\mathfrak{a})$ predicate $\phi(x)$ such that in any model in which ψ defines a function we have: $\phi(x)$ holds if and only if x is a point of continuity of f.*

(2) *There is a $\Pi_2^0(b)$ predicate $\hat\psi(x,y)$ with parameter $b \in L$ such that in any model where ψ defines a function we have: $\hat\psi(x,y)$ if and only if $\psi(x,y) \wedge \phi(x)$.*

PROOF. (1) Recall that for $x \in \mathbb{R}$, the *oscillation* of f at x is defined as

$$\mathrm{osc}_f(x) = \inf\{\mathrm{diam}(f(U)) : x \in U \wedge U \subseteq \mathbb{R} \text{ is open}\},$$

and that x is a point of continuity precisely when $\mathrm{osc}_f(x) = 0$. Let $\phi(x,\varepsilon)$ be the following predicate:

$$(\exists q,r,\delta \in \mathbb{Q}_+)|x-q| < r \wedge [(\forall x_0,x_1)(\forall y_0,y_1)(f(x_0) = y_0 \wedge f(x_1) = y_1 \wedge$$
$$|x_0 - q| < r \wedge |x_1 - q| < r) \longrightarrow |y_0 - y_1| < \varepsilon - \delta].$$

This is Π_2^1 and $\phi(x,\varepsilon)$ holds precisely when $\mathrm{osc}_f(x) < \varepsilon$. On the other hand, it is easy to see that $\{(x,\varepsilon) \in \mathbb{R} \times \mathbb{Q}_+ : \mathrm{osc}_f(x) < \varepsilon\}$ is open (when $\mathbb{Q}_+$ has the discrete topology), and so $\{(x,\varepsilon) \in \mathbb{R} \times \mathbb{Q}_+ : \hat\psi(x,\varepsilon)\}$ is an open Π_2^1 set. It follows

from Lemma 3.9 that there is a $\Sigma^0_1(\mathfrak{a})$ predicate $\hat{\phi}(x,\varepsilon)$ such that $\hat{\psi}(x,\varepsilon)$ iff $\hat{\phi}(x,\varepsilon)$. Thus if we let $\phi(x)$ be $(\forall \varepsilon \in \mathbb{Q}_+)\hat{\phi}(x,\varepsilon)$ then $\phi(x)$ is a $\Pi^0_2(\mathfrak{a})$ predicate which holds precisely when x is a point of continuity of f, and ϕ does so in any model where $\psi(x,y)$ defines a function.

(2) Fix a sequence $(x_n)_{n\in\omega}$ in $\mathbb{R} \cap L$ such that

$$L \models \text{``}(x_n)_{n\in\omega} \text{ is dense in } \{x \in \mathbb{R} : \phi(x)\}\text{''}.$$

To say that (x_n) is dense in $\{x \in \mathbb{R} : \phi(x)\}$ can be expressed as

$$(\forall \varepsilon \in \mathbb{Q}_+)(\forall x)(\phi(x) \longrightarrow (\exists n)|x_n - x| < \varepsilon),$$

which is $\Pi^1_1(\mathfrak{a}, (x_n)_{n\in\omega})$, and so this statement is absolute. Let $(y_n)_{n\in\omega}$ be the sequence in $\mathbb{R} \cap L$ defined by $y_n = f(x_n)$, and let $\hat{\psi}(x,y)$ be the predicate

$$\phi(x) \wedge (\forall \varepsilon \in \mathbb{Q}_+)(\exists n)|x_n - x| < \varepsilon \wedge |y_n - y| < \varepsilon.$$

Then $\hat{\psi}(x,y)$ is $\Pi^0_2(\mathfrak{a}, (x_n)_{n\in\omega}, (y_n)_{n\in\omega})$ and since f is continuous on the set $\{x \in \mathbb{R} : \phi(x)\}$ it holds that

$$\hat{\psi}(x,y) \iff \phi(x) \wedge \psi(x,y)$$

in any model where ψ defines a function, as required. ⊣

LEMMA 3.11. *Let $f : \mathbb{R} \to \mathbb{R}$ be a function. Then:*
(1) *There is a Π^1_1 set $H \subset \mathbb{R}$ such that*

$$\{x \in \mathbb{R} : f'(x) \text{ exists}\} \subseteq H$$

and $\{y \in \mathbb{R} : f^{-1}(y) \cap H \text{ is uncountable}\}$ is Lebesgue null.

(2) *If f is defined by the Δ^1_2 predicate $\psi(x,y)$ then there is a $\Pi^1_1(b)$ predicate $\chi(x)$ with a parameter $b \in L$ such if we let $H = \{x \in \mathbb{R} : \chi(x)\}$ then (1) holds for this H and f defined by ψ in any model where $\psi(x,y)$ defines a function.*

PROOF. (1) Let C be the set of points of continuity of f. It is well-known that this is a G_δ set. Define

$$x \in H \iff x \in C \wedge (\exists y)(\forall \varepsilon > 0)(\exists \delta > 0)(\forall z \in C \setminus \{x\})$$
$$\left[|x - z| < \delta \implies |\frac{f(x) - f(z)}{x - z} - y| < \varepsilon\right]. \tag{3.1}$$

It is clear that if $f'(x)$ exists then $x \in H$. Permitting a slight *abus de langage* here and below, we extend the definition of $f'(x)$ to all of H be letting $f'(x) = y$ where y is such that (3.1) is satisfied.

CLAIM 3.12. *H is Π^1_1.*

PROOF. Let $\bar{f} : \mathbb{R} \to \mathbb{R}$ be a Borel function such that $\bar{f} \restriction C = f \restriction C$. We claim that $x \in H$ if and only if

$$x \in C \wedge (\forall \varepsilon > 0)(\exists q \in \mathbb{Q})(\exists \delta > 0)(\forall z \neq x)$$
$$\left[(z \in C \wedge |x - z| < \delta) \implies \left|\frac{\bar{f}(x) - \bar{f}(z)}{x - z} - q\right| < \varepsilon\right]. \tag{3.2}$$

If x is isolated in C then clearly (3.1) holds for x if and only if (3.2) holds. So assume that x is not isolated. If (3.2) holds for x, let q_n witness (3.2) with $\varepsilon = \frac{1}{2^{n+1}}$. Then $|q_{n+1} - q_n| \leq \frac{1}{2^n}$ so q_n is Cauchy, and if we let $y = \lim_{n\to\infty} q_n$ then y is easily

seen to be a witness to (3.1). Conversely, if (3.1) holds for x, and y is a witness to this, then let $q_n \in \mathbb{Q}$ be a sequence of rationals such that $q_n \to y$. Then it is clear that for all $\varepsilon > 0$ we can find some n such that q_n is a witness to that (3.2) holds. Since (3.2) is $\mathbf{\Pi}_1^1$, the claim is proved.   ⊣

CLAIM 3.13. $\{y \in \mathbb{R} : f^{-1}(y) \cap H \text{ is uncountable}\}$ is Lebesgue null.

PROOF. The proof uses the idea from [7, Ch. 5.15]. It clearly suffices to show for all $m \in \mathbb{N}$ that the sets
$$Y_m = \{y \in [-m, m] : f^{-1}(y) \cap H \text{ is uncountable}\}$$
are null; we will prove this for Y_1, from which the other cases follow by rescaling the codomain of f (or by an identical proof). For $y \in Y_1$, pick $t_y \in f^{-1}(y) \cap H$ such that $f'(t_y) = 0$. Such a t_y exists since when $y \in Y_1$ the set $f^{-1}(y) \cap H$ is uncountable and so it contains an accumulation point, and as f is constant on this set we must have $f'(t) = 0$ at any accumulation point. Let $T = \{t_y : y \in Y_1\}$, and note that $f(T) = Y_1$.

Let $\varepsilon > 0$, and for each $t \in T$ let $1 > \delta_t > 0$ be such that for all $z \in C$ with $|t - z| < \delta_t$ we have
$$\left| \frac{f(t) - f(z)}{t - z} \right| < \varepsilon,$$
and let $I_t = (t - \delta_t, t + \delta_t)$. Note that for any $z \in I_t \cap C$ we have $f(z) \in (f(t) - \varepsilon\delta_t, f(t) + \varepsilon\delta_t)$, and so we have $\mu(f(I_t \cap C)) \leq 2\varepsilon\delta_t$. Since the intervals I_t cover T, we can find $t_i \in T$, $i \in \mathbb{N}$, such that $U = \bigcup_{t \in T} I_t = \bigcup_{i \in \mathbb{N}} I_{t_i}$. We claim that $\mu(f(U \cap C)) \leq 4\varepsilon$. To see this it is enough to prove that $\mu(f(K \cap C)) \leq 4\varepsilon$ for all compact $K \subseteq U$. If $K \subseteq U$ is compact, then we can find $N \in \mathbb{N}$ such that $K \subseteq \bigcup_{i=1}^N I_{t_i}$. Moreover, after possibly going to a subcover, we can assume that each $x \in K$ is contained in at most two different intervals I_{t_i}, and so we have $\sum_{i=1}^N 2\delta_{t_i} = \sum_{i=1}^N \mu(I_{t_i}) \leq 4$. Thus
$$\mu(f(C \cap K)) \leq \mu(f(\bigcup_{i=1}^N I_{t_i} \cap C)) \leq \sum_{i=1}^N \mu(f(I_{t_i} \cap C)) \leq \sum_{i=1}^N 2\varepsilon\delta_{t_i} = 4\varepsilon,$$
as required. It follows that $\mu(f(U \cap C)) = 0$, and so since $T \subseteq U \cap C$ we have $\mu(Y_1) = \mu(f(T)) = 0$.   ⊣

(2) Let $\phi(x)$ and $\hat{\psi}(x, y)$ be the predicates defined in Lemma 3.10, and let $\chi(x)$ be the predicate
$$\phi(x) \wedge (\forall \varepsilon > 0)(\exists q \in \mathbb{Q})(\exists \delta > 0)(\forall z)(\forall y_0, y_1)$$
$$(\phi(z) \wedge z \neq x \wedge |x - z| < \delta \wedge \hat{\psi}(x, y_0) \wedge \hat{\psi}(z, y_1)) \longrightarrow |\frac{y_0 - y_1}{x - z} - q| < \varepsilon.$$

Then $\chi(x)$ is $\mathbf{\Pi}_1^1(b)$ (where $b \in L$ is the parameter in $\hat{\psi}$), and if $\psi(x, y)$ defines a function then the set $\{x \in \mathbb{R} : \chi(x)\}$ is equal to the set H defined in (3.1).   ⊣

PROOF OF THEOREM 3.8. We may assume that $\aleph_1^L = \aleph_1$. Fix $f : \mathbb{R} \to \mathbb{R}^2 : x \mapsto (f_1(x), f_2(x))$ as in the statement of the theorem. Applying Lemma 3.11 to f_1 and f_2, there are $\mathbf{\Pi}_1^1(b)$ ($b \in L$) sets H_1 and H_2 defined by $\mathbf{\Pi}_1^1(b)$ formulas $\chi_1(x)$ and $\chi_2(x)$ such that
$$Y_i = \{y \in \mathbb{R} : |f_i^{-1}(y) \cap H_i| > \aleph_0\}$$

is Lebesgue null, and such that the points of differentiability of f_i are contained in H_i. Let
$$Y_i^L = \{y \in \mathbb{R} \cap L : L \models |f_i^{-1}(y) \cap H_i| > \aleph_0\}.$$
We claim that $Y_i \cap L = Y_i^L$. Since $\aleph_1 = \aleph_1^L$ it is clear that if $y \in Y_i^L$ then $y \in Y_i$. On the other hand, note that the set $\Gamma_{i,y} = f_i^{-1}(y) \cap H_i$ is $\Delta^1_2(b, y)$, and so if $y \in L$ then by Lemma 2.2 the set $\Gamma_{i,y}$ is countable if $\Gamma_{i,y} \cap L$ is. So if $y \in L \setminus Y_i^L$ then $y \notin Y_i$, as required.

Let $\mathbb{R}^* = \mathbb{R} \setminus (Y_1 \cup Y_2)$, which has full measure, and let $A_1 = f(H_2)$ and $A_2 = f(H_1)$. Then A_1 and A_2 are $\Sigma^1_2(b)$ sets, and since either $f_1'(t)$ or $f_2'(t)$ exists for all $t \in \mathbb{R}$ we must have that $\mathbb{R} = A_1 \cup A_2$. For any $r \in \mathbb{R}^*$ the set $f_i^{-1}(r) \cap H_i$ is countable by the definition of $\mathbb{R}^*$, and so there are at most countably many $t \in H_i$ such that $f_i(t) = r$. Since $Y_1^L \cup Y_2^L$ is null in L there are uncountably many constructible reals $(x_\alpha : \alpha < \aleph_1)$ not belonging to $Y_1^L \cup Y_2^L$, and therefore not to $Y_1 \cup Y_2$. On the other hand, since $\mathbb{R}^*$ has full measure there is $r \in \mathbb{R}^* \setminus L$. The horizontal section $A_1^{x_\alpha}$ contains only constructible reals since $A_1^{x_\alpha}$ is $\Sigma^1_2(a, x_\alpha)$, and so if it contained a non-constructible real then it would be uncountable by Theorem 2.3. Since $A_1 \cup A_2$ cover $\mathbb{R}^2$ it must then be the case that the vertical section $(A_2)_r$ contains all the points of the form (r, x_α). But this contradicts that $(A_2)_r$ is countable. ⊣

3.5. Polarized partitions. Another type of statement that can be proved by counting arguments analogous to the above are polarized partition relations for Σ^1_2 colourings of $\mathscr{R} \times \mathscr{R}$ (where, as in Section 2, $\mathscr{R}$ refers to an uncountable recursively presented Polish space). These may be viewed as regularity properties that Σ^1_2 colourings have in the presence of a non-constructible real.

We have the following definable analogue of [7, 24.27]:

THEOREM 3.14. *The following are equivalent*:

(1) $\mathscr{R} \nsubseteq L$.
(2) *For every Σ^1_2-definable function $f : \mathscr{R} \times \mathscr{R} \to \omega$ there are sets $C, D \subseteq \mathscr{R}$ such that $|C| = |D| = 2$ and $f \restriction C \times D$ is monochromatic.*
(3) *For every Σ^1_2-definable function $f : \mathscr{R} \times \mathscr{R} \to \omega$ there are sets $C, D \subseteq \mathscr{R}$ such that $|C| = |D| = \aleph_0$ and $f \restriction C \times D$ is monochromatic.*
(4) *For every Σ^1_2-definable function $f : \mathscr{R} \times \mathscr{R} \to \omega$ there are countably infinite Σ^1_2 sets $C, D \subseteq \mathscr{R}$ such that $f \restriction C \times D$ is monochromatic.*

PROOF. $(4) \implies (3) \implies (2)$ is clear.

$(1) \implies (3)$: We may assume that $V = L[z]$ for some $z \notin L$ and that it holds that $\aleph_1^L = \aleph_1$. Assume (3) fails, and fix f witnessing this. For $s \in [\mathscr{R}]^\omega$, let
$$T(s, i) = \{y \in \mathscr{R} : (\forall x \in s) f(x, y) = i\}.$$
Since s is a countable sequence this quantification over s may be replaced by a number quantifier over the domain of s. Thus $T(s, i)$ is $\Sigma^1_2(s)$. By assumption we have that $|T(s, i)| < \aleph_0$ and so $T(s, i) \subseteq L$ by Theorem 2.3. Let $U_i = \{x \in \mathscr{R} \cap L : f(x, z) = i\}$. Then $|U_i| < \aleph_0$ since otherwise we could find $s_i \subseteq U_i$ of size $\aleph_0$ from which $z \in T(s_i, i)$ would follow, contradicting that $z \notin L$. But now we have
$$\mathscr{R} \cap L = \bigcup_{i \in \omega} U_i$$

so that $|\mathscr{R} \cap L|$ is countable, a contradiction.

(3) $\implies$ (4): Fix a Σ_2^1-definable function f and $i \in \omega$ such that there exists $C, D \in [\mathscr{R}]^\omega$ with $f(C \times D) = \{i\}$. Since f is Σ_2^1 the set

$$\{(C, D) \in [\mathscr{R}]^\omega \times [\mathscr{R}]^\omega : (\forall (x, y) \in C \times D) f(x, y) = i\}$$

is Σ_2^1, and it is non-empty by the above. Thus by Σ_2^1-uniformization (e.g., [12, 4E.4]) it contains a Σ_2^1 definable pair (C, D).

(2) $\implies$ (1): Suppose $\mathscr{R} \subseteq L$ and let $\preceq$ denote the usual Σ_2^1 well-ordering of $L \cap \mathscr{R}$. Recall $\mathrm{IS}^\#(x, y)$ from 2.3, and define

$$f(x, y) = \begin{cases} \mathrm{IS}^\#(y, x) + 1 & \text{if } x \prec y, \\ \mathrm{IS}^\#(x, y) + 1 & \text{if } y \prec x, \\ 0 & \text{if } x = y. \end{cases}$$

Let $\{x, x'\}, \{y, y\} \subseteq \mathscr{R}$ where $x \neq x'$ and $y \neq y'$, and assume that $x, x', y \preceq y'$. Then $f(x, y') \neq f(x', y')$, and so $f \restriction \{x, x'\} \times \{y, y'\}$ is not monochromatic. ⊣

3.6. A Schur type partition result. As an application of Theorem 3.14 we prove the following definable analogue of [7, 24.37].

THEOREM 3.15. *There is a non-constructible real if and only if for any Σ_2^1 colouring $g : \mathbb{R} \to \omega$ there are four distinct $x_{00}, x_{01}, x_{10}, x_{11} \in \mathbb{R}$ of the same colour such that*

$$x_{00} + x_{11} = x_{01} + x_{10}.$$

PROOF. Assume $\mathbb{R} \nsubseteq L$ and let $g : \mathbb{R} \to \omega$ be a colouring. By [5, 19.2] we can find a continuous $h : 2^\omega \to \mathbb{R}$ such that $h(2^\omega)$ is linearly independent over $\mathbb{Q}$. It may be shown using [4] that this h can be taken to be Δ_1^1. Now let $f : 2^\omega \times 2^\omega \to \omega :$ $(x, y) \mapsto g(h(x) + h(y))$. Then by Theorem 3.14 we can find $x_0 \neq x_1$ and $y_0 \neq y_1$ such that $f \restriction \{x_0, x_1\} \times \{y_0, y_1\}$ is monochromatic. If we let $x_{ij} = h(x_i) + h(y_j)$ for $0 \leq i, j \leq 1$ then clearly $x_{00} + x_{11} = x_{01} + x_{1,0}$ and these are distinct since $h(2^\omega)$ is linearly independent over $\mathbb{Q}$.

Conversely, assume that $\mathbb{R} \subseteq L$. We define a function $g : \mathbb{R} \to \omega$ by

$$g(x) = m \iff (\exists \epsilon) R^{\mathbb{R}}(x, \epsilon) \wedge \pi_\epsilon(m) = x,$$

where $R^{\mathbb{R}}$ is defined as in Section 2.2. Then g is Σ_2^1. Let $x_{00}, x_{01}, x_{10}, x_{11} \in \mathbb{R}$ be distinct, and let $\alpha < \omega$ be least such that $x_{i,j} \in L_\alpha$ for all $0 \leq i, j \leq 1$. It cannot be the case that three of the x_{ij} are already in some L_β where $\beta < \alpha$, since the L_α are closed under addition. Thus two of the $x_{i,j}$ are in L_α and not in any L_β for $\beta < \alpha$. But then these two $x_{i,j}$ are coloured differently by g. ⊣

REMARK 3.16. It is clear from the above that what is really needed to make all of the above theorems work for Σ_n^1 (or more generally, $\Sigma_n^1(a)$ versions) is an inner model relative to which we have a Σ_n^1 absoluteness principle and a perfect set theorem for Σ_n^1. If we have this, then we will be able to prove that the Σ_n^1 versions of the statements in Theorem 1.1 and Theorem 3.14 are equivalent to all reals being in that inner model.

For example, it is well-known (see [3, Section 15]) that if there is a measurable cardinal κ, and U is an ultrafilter witnessing this, then the inner model $L[U]$ has this relationship to the class of Σ_3^1 sets, provided that $0^\sharp$ does not exist. Thus in this

context we obtain Σ_3^1 versions of Theorem 1.1 and Theorem 3.14, with L replaced by $L[U]$.

Philip Welch has further pointed out to us that you can more generally do this using the core model K below one Woodin cardinal. Assume (i) there exists a measurable cardinal and (ii) sharps for reals. Then this model is Σ_3^1 correct, and so the above theorems work over this model.

REMARK 3.17. The referee has made the following remark, which expands on Philip Welch's observation above. Assuming the existence of appropriate large cardinals, the required correctness of the core model can be lifted throughout the projective hierarchy, replacing K with M_n (see [16, 7.2] for the definition of M_n), but this does not quite suffice to fully capture all reals in the analogues of the theorems of the paper. Thus the question of level by level equivalences throughout the full projective hierarchy appear interesting.

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