

DECIDABILITY AND CLASSIFICATION OF THE THEORY OF INTEGERS WITH PRIMES

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Abstract. We show that under Dickson's conjecture about the distribution of primes in the natural numbers, the theory $Th(\mathbb{Z}, +, 1, 0, Pr)$ where Pr is a predicate for the prime numbers and their negations is decidable, unstable, and supersimple. This is in contrast with $Th(\mathbb{Z}, +, 0, Pr, <)$ which is known to be undecidable by the works of Jockusch, Bateman, and Woods.

§1. Introduction. It is well known that Presburger arithmetic $T_{+,<} = Th(\mathbb{Z}, +, 0, 1, <)$ is decidable and enjoys quantifier elimination after introducing predicates for divisibility by n for every natural number $n > 1$ (see e.g., [9, Corollary 3.1.21]). The same is true for $T_+ = Th(\mathbb{Z}, +, 0, 1)$. This is, of course, in contrast to the situation with the theory of Peano arithmetics or $Th(\mathbb{Z}, +, \cdot, 0, 1)$ which is not decidable.

If we are interested in classifying these theories in terms of stability theory, quantifier elimination gives us that T_+ is superstable of U -rank 1, while $T_{+,<}$ is dp-minimal (a subclass of dependent, or NIP, theories, see e.g., [5, 10, 15]).

Over the years there has been quite extensive research on structures with universe \mathbb{Z} or \mathbb{N} and some extra structure, usually definable from Peano. A very good survey regarding questions of decidability is [2] and a list of such structures defining addition and multiplication is available in [8].

Less research was done on classifying these structures stability-theoretically. For instance, in [12, Theorem 25] and also in [11] it is proved that $Th(\mathbb{Z}, +, 0, P_q)$ is superstable of U -rank ω , where P_q is the set of powers of q .

In this paper we are interested in adding a predicate Pr for the primes and their negations and we consider $T_{+,Pr} = Th(\mathbb{Z}, +, 0, 1, Pr)$ and $T_{+,Pr,<} = Th(\mathbb{Z}, +, 0, 1, Pr, <)$. The language $\{+, 0, 1, Pr\}$ allows us to express famous number-theoretic conjectures such as the twin prime conjecture (for every n , there are at least n pairs of primes/negation of primes of distance 2), and a version of Goldbach's conjecture (all even integers can be expressed as a difference or a sum of primes). Adding the order allows us to express Goldbach's conjecture in full.

Up to now, the only known results about the theory are under a strong number-theoretic conjecture known as Dickson conjecture (D) (see below), which is also the assumption in the works of Jockusch, Bateman, and Woods. In [1, 19], they proved

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that assuming Dickson conjecture, $Th(\mathbb{N}, +, 0, Pr)$ is undecidable and even defines multiplication. It follows immediately that $T_{+,Pr,<}$ is undecidable and as complicated as possible in the sense of stability theory. This also explains why we need Pr to include also the negation of primes: by relatives of the Goldbach Conjecture (which are proved, see e.g., [17]), every positive integer greater than N is a sum of K primes for some fixed K, N , and hence the positive integers themselves are also definable from the positive primes.

CONJECTURE 1.1 (D) (Dickson, 1904 [6]). *Let $k \geq 1$ and $\bar{f} = \langle f_i \mid i < k \rangle$ where $f_i(x) = a_i x + b_i$ with a_i, b_i non-negative integers, $a_i \geq 1$ for all $i < k$. Assume that the following condition holds:*

$\star_{\bar{f}}$ *There does not exist any integer $n > 1$ dividing all the products $\prod_{i < k} f_i(s)$ for every (non-negative) integer s .*

Then there exist infinitely many natural numbers m such that $f_i(m)$ is prime for all $i < k$.

Note that in fact the condition $\star_{\bar{f}}$ follows easily from the conclusion that there are infinitely many m 's with $f_i(m)$ prime for all $i < k$. See also Remark 2.6.

Dickson's conjecture is the linear case of Schinzel's Hypothesis, see [13, pg. 292] for a discussion.

Our main result is the following.

THEOREM 1.2. *Assuming (D), the theory $T_{+,Pr}$ is decidable, unstable and supersimple of U -rank 1.*

In essence (D) implies that the set of primes is generic up to congruence conditions (while it is not generic in the sense of [3]), and this allows us to get quantifier elimination in a suitable language. Forking then turns out to be trivial: forking formulas are algebraic (Theorem 3.2).

To show that $T_{+,Pr}$ is unstable we show that it has the independence property (see Proposition 3.6). This turns out to follow from the proof of the Green-Tao theorem about arithmetic progressions in the primes [7] (i.e., without using (D)), as was told to us in a private communication by Tamar Ziegler (but we also show that this follows from (D)).

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§2. Quantifier elimination. In this section we will prove quantifier elimination in $T_{+,Pr}$ assuming (D) in a suitable language.

Let us first note some useful facts about (D).

REMARK 2.1. Suppose that $\langle f_i \mid i < k \rangle$ is as in (D) and $f_i(x) = a_i x + b_i$. Let

$$N = \max(\{a_i \mid i < k\} \cup \{k\}) + 1.$$

Then $\star_{\bar{f}}$ holds iff for every prime $p < N$, p does not divide $\prod_{i < k} f_i(s)$ for all $s \in \mathbb{Z}$ where

PROOF. If $\star_{\bar{f}}$ fails, then there is some prime p such that p divides $\prod_{i < k} f_i(s)$ for all s . Let $P(X) \in \mathbb{Z}[X]$ be the polynomial $\prod_{i < k} f_i(X)$. Let $P_p = P \pmod{p} \in \mathbb{F}_p[X]$ (where \mathbb{F}_p is the prime field of size p). It follows that $P_p(a) = 0$ for all $a \in \mathbb{F}_p$. So either $P_p = 0$ or $k \geq \deg(P_p) \geq p$, hence $p \leq k$ or $\prod_{i < k} a_i \equiv 0 \pmod{p}$ (as the leading coefficient) which means that for some $i < k$, $a_i \geq p$, so $p < N$ and we are done. \dashv

LEMMA 2.2. *Assume (D). Then (D) holds also when we allow b_i to be negative.*

PROOF. Suppose that $\langle f_i \mid i < k \rangle$ is a sequence of linear maps $f_i(x) = a_i x + b_i$ where $a_i \geq 1$ and $b_i \in \mathbb{Z}$, and assume that $\star_{\bar{f}}$ holds. Let N be as in Remark 2.1. Let $K = N!$ (enough to take the product of the primes below N). Suppose that $l \in \mathbb{N}$ is such that $lK + b_i > 0$ for all $i < k$. Let $f'_i(x) = a_i x + a_i lK + b_i$. Then $a_i \geq 1$, $b'_i = a_i lK + b_i > 0$, so let us show that $\star_{\bar{f}'}$ holds (where $\bar{f}' = \langle f'_i \mid i < k \rangle$). Note that when we compute N in Remark 2.1, we only use k and a_i which haven't changed, so by that remark, it is enough to check that for no prime $p < N$, $\prod_{i < k} f'_i(s) \equiv 0 \pmod{p}$ for all s . But for such p 's, $f'_i(s) = f_i(s) + a_i lK \equiv f_i(s) \pmod{p}$, so $\prod_{i < k} f'_i(s) \equiv \prod_{i < k} f_i(s) \not\equiv 0 \pmod{p}$.

By (D), there are infinitely many integers m such that $f'_i(m)$ is prime for all $i < k$. But $f'_i(m) = a_i m + a_i lK + b_i = a_i(m + lK) + b_i$. Hence substituting $m + lK$ for m we get what we wanted. \dashv

LEMMA 2.3. *Assume (D). Suppose that $k, k' \in \mathbb{N}$ and $\langle a_i, b_i \mid i < k \rangle, \langle c_j, d_j \mid j < k' \rangle$ are two tuples of integers with $a_i, c_j \geq 1$ for all $i < k, j < k'$.*

For $i < k$, let $f_i(x) = a_i x + b_i$ and for $j < k'$, let $g_j(x) = c_j x + d_j$.

Suppose that $\star_{\bar{f}}$ holds and that $(a_i, b_i) \neq (c_j, d_j)$.

Then there are infinitely many natural numbers m for which for all $i < k$ and $j < k'$, $f_i(m)$ is prime and $g_j(m)$ is composite.

Before giving the proof, we note that this lemma generalizes Lemma 1 from [1], which was key in the proof there of the undecidability of $T_{+,Pr,<}$.

COROLLARY 2.4 ([1, Lemma 1]). *(Assuming (D)) Let b_0, \dots, b_{n-1} be an increasing sequence of natural numbers, and assume that there is no prime p such that $\{b_i \pmod{p} \mid i < n\} = p$. Then there are infinitely many natural numbers x such that $x + b_0, \dots, x + b_{n-1}$ are consecutive primes.*

PROOF OF COROLLARY. This is immediate from Lemma 2.3 by taking $f_i(x) = x + b_i$ and $g_j(x) = x + c_j$ where c_j run over all numbers between the b_j 's. \dashv

PROOF OF LEMMA. By induction on k' . For $k' = 0$ there is nothing to prove by (D) and Lemma 2.2.

Suppose the lemma is true for k' and prove it for $k' + 1$. It is enough to prove that for any n , there is some $m > n$ such that $f_i(m)$ is prime for all $i < k$ and $g_j(m)$ is not prime for all $j \leq k'$.

Fix n . We may assume by enlarging it that for no $m > n$ is it the case that $f_i(m) = g_j(m)$ for $i < k, j \leq k'$.

Let $m > n$ be so that $f_i(m)$ is prime for all $i < k$ and $g_j(m)$ is composite for all $j < k'$. If it happens that $g_{k'}(m)$ is composite, then we are done, so suppose that $q = g_{k'}(m)$ is prime. Let $f'_i(x) = a_i(qx + m) + b_i$ and $g'_j(x) = c_j(qx + m) + d_j$ for $i < k$ and $j < k' + 1$. Then $g'_{k'}(x) = c_{k'}qx + q$ is composite for all $x \geq 1$

(so that $c_jx + 1 \geq 2$). Hence it is enough to find m' large enough so that $f'_i(m')$ is prime for all $i < k$ and $g'_j(m')$ is composite for all $j < k'$.

By the induction hypothesis, it is enough to check that $\star_{\vec{f}}$ holds (because $(a_iq, a_im + b_i) \neq (c_jq, c_jm + d_j)$). Suppose that $p > 1$ is a prime which divides $\prod_{i < k} f'_i(s)$ for all s . Hence $\prod_{i < k} f'_i(s) \equiv 0 \pmod{p}$, and if $p \neq q$, it follows (as q is invertible modulo p) that $\prod_{i < k} f_i(s) \equiv 0 \pmod{p}$ for all s — a contradiction. If $p = q$, then $f'_i(x) \equiv a_im + b_i \equiv f_i(m) \pmod{q}$ for all x , hence for some $i < k$, $f_i(m) = q = g_{k'}(m)$, contradicting our choice of m . \dashv

Expand the language $L = \{+, Pr, 0, 1\}$ to include the Presburger predicates P_n for $2 \leq n < \omega$ interpreted as $P_n(x) \Leftrightarrow x \equiv 0 \pmod{n}$, and also the predicates Pr_n for $2 \leq n < \omega$ interpreted as $Pr_n(x) \Leftrightarrow P_n(x) \wedge Pr(x/n)$. We need the latter predicate in order to eliminate the quantifiers from $\varphi(x) = \exists y (ny = x \wedge Pr(y))$. We also add negation (as a unary function). We need negation because of formulas of the form $\varphi(x, y) = Pr(x - y) = \exists w (w + y = x \wedge Pr(w))$.

Let L^* be the resulting language $\{+, -, 1, 0, Pr, Pr_n, P_n \mid 2 \leq n < \omega\}$, and let $T^*_{+,Pr}$ be the complete theory of M^* — the structure with universe \mathbb{Z} in L^* . Note that all the new predicates are definable from L .

REMARK 2.5. The condition $\star_{\vec{f}}$ of Dickson’s conjecture is first-order expressible in L^* . This means that for every tuple $\langle a_i \mid i < k \rangle$ of positive integers, there is a formula $\varphi_{\vec{a}}(y_0, \dots, y_{k-1})$ such that for any choice of $b_i \in \mathbb{Z}$ for $i < k$, $M^* \models \varphi_{\vec{a}}(\vec{b})$ iff $\star_{\vec{f}}$ holds where $f_i(x) = a_ix + b_i$ for $i < k$. It has the form $\bigwedge_{p < N \text{ prime}} \bigvee_{r < p} \bigwedge_{i < k} \neg P_p(a_ir + y_i)$ for some $N \in \mathbb{N}$.

PROOF. Recall Remark 2.1 and the choice of N from there (which depends only on $\langle a_i \mid i < k \rangle$ and k). Let $\varphi_{\vec{a}}(\vec{y})$ be as described in the remark: for every prime $p < N$, for some $0 \leq x < p$, for all $i < k$, $\neg P_p(a_ix + y_i)$. Note that $\varphi_{\vec{a}}$ is quantifier-free in L^* (as it contains 1). \dashv

REMARK 2.6. Given $\vec{f} = \langle f_i \mid i < k \rangle$ a tuple of linear maps as above, if there are more than $2k$ integers m such that $f_i(m)$ is prime or a negation of a prime, then $\star_{\vec{f}}$ holds.

PROOF. Indeed, otherwise there is some prime p which witnesses this, but then for some i and three different m ’s, $|f_i(m)| = p$ — a contradiction. \dashv

LEMMA 2.7. $T^*_{+,Pr}$ eliminates quantifiers in L^* provided (D).

PROOF. We start with the following observation.

◇ By Remark 2.5 and Lemma 2.3, our assumption that Dickson’s conjecture holds translates into a scheme of first-order statements:

For every n and every choice of positive integers $\langle a_i \mid i < k \rangle$ and $\langle a'_j \mid j < k' \rangle$ and for all $\langle b_i \mid i < k \rangle$ and $\langle b'_j \mid j < k' \rangle$, if $\varphi_{\vec{a}}(\vec{b})$ holds and for all $i < k, j < k'$, $(a_i, b_i) \neq (a'_j, b'_j)$ then there are at least n elements x with

$$\bigwedge_{i < k} Pr(a_ix + b_i) \wedge \bigwedge_{j < k'} \neg Pr(a'_jx + b'_j).$$

Conversely, by Remark 2.6, if there are more than $2k$ such elements x , then $\varphi_{\bar{a}}(\bar{b})$ holds. In particular, $\varphi_{\bar{a}}(\bar{b}) \wedge \bigwedge_{i,j} (a_i, b_i) \neq (a'_j, b'_j)$ holds iff there are more than $2k$ elements x with

$$\bigwedge_{i < k} Pr(a_i x + b_i) \wedge \bigwedge_{i < k'} \neg Pr(a'_i x + b'_i).$$

(Recall that Pr contains the primes and their negations.)

In order to prove quantifier elimination we will use a back-and-forth criteria. Namely, suppose that $\mathfrak{C} \models T_{+,Pr}^*$ is a monster model (very large, saturated model) and that $h : A \rightarrow B$ is an isomorphism of small substructures A, B . Given $a \in \mathfrak{C} \setminus A$ we want to extend h so that its domain contains a .

We may assume, by our choice of language (which includes Pr_n and $-$), that both A and B are groups such that if $c \in A$ and $\mathfrak{C} \models Pr_n(a)$ then $c/n \in A$ and similarly for B . Why? For such a c , elements of the group generated by adding c/n to A have the form $m(c/n) + b$ for $m \in \mathbb{Z}$ and $b \in A$. We have to show that the map taking c/n to $h(c)/n$ and extends h is an isomorphism. For instance, we have to show that if $\mathfrak{C} \models Pr(m(c/n) + b)$ then $\mathfrak{C} \models Pr(m(h(c)/n) + h(b))$. But $\mathfrak{C} \models Pr(m(c/n) + b)$ iff $\mathfrak{C} \models Pr_n(mc + nb)$. Similarly we deal with Pr_k and P_k .

Let $p^a(x) = \text{tp}^{\text{qf}}(a/A)$, and let $q^a(x) = h(p^a)$. Let $p_{\equiv}^a = p^a \upharpoonright L_{\equiv}^*$ and $p_{Pr}^a = p^a \upharpoonright L_{Pr}^*$, where $L_{\equiv}^* = L^* \setminus \{Pr, Pr_n \mid 2 \leq n < \omega\}$ and $L_{Pr}^* = L^* \setminus \{P_n \mid 2 \leq n < \omega\}$, so that $p^a = p_{\equiv}^a \cup p_{Pr}^a$, and we have to realize q^a .

CLAIM 2.8. *It is enough to prove that we can realize $q_{Pr}^a = h(p_{Pr}^a)$ for all a as above.*

PROOF. Easily, as we included 1 in the language, q_{\equiv}^a is isolated by $\{x \neq c \mid c \in B\}$ and equations of the form $x \equiv k \pmod{n}$ for $k < n$, and for every $n < \omega$ there is exactly one $k < n$ with such an equation appearing in q^a . Also, every finite set of such equations is implied by one such equation (e.g., if the equations are $\{x \equiv k_i \pmod{n_i} \mid i < s\}$ then take $x \equiv k \pmod{\prod_{i < s} n_i}$ where k is such that this equation is in q^a). Hence it is enough to show that $x \equiv k \pmod{n} \cup q_{Pr}^a(x)$ is consistent (q_{Pr}^a already contains $\{x \neq c \mid c \in B\}$). As $a \equiv k \pmod{n}$, $b = (a - k)/n \in \mathfrak{C}$. Let $p^b = \text{tp}^{\text{qf}}(b/A)$ so by our assumption there is some $d \in \mathfrak{C}$ such that $d \models h(p^b)_{Pr}$. Then $nd + k \models q_{Pr}^a(x)$ and of course satisfies the equation $x \equiv k \pmod{n}$. ⊢

Let $p_{Pr_0}^a = p^a \upharpoonright L_{Pr_0}$ where $L_{Pr_0} = L_{Pr} \setminus \{Pr_n \mid 2 \leq n < \omega\}$.

CLAIM 2.9. *It is enough to prove that we can realize $q_{Pr_0}^a = h(p_{Pr_0}^a)$ for a as above.*

PROOF. This is similar to Claim 2.8. It is enough to show that $q_{Pr_0}^a(x) \cup \Sigma(x)$ is consistent where Σ is a finite set of formulas from $q_{Pr}^a \setminus q_{Pr_0}^a$. So Σ consists of formulas of the form $Pr_n(mx + c)$ or its negation for $m \in \mathbb{Z}$, $1 < n \in \mathbb{N}$ and $c \in B$. Without loss of generality, by replacing the n 's with their product N and $Pr_n(mx + c)$ by $Pr_N((N/n)(mx + c))$, we may assume that all the n 's appearing in Σ are equal to $n > 1$. Let $b = (a - k)/n$ where $a \equiv k \pmod{n}$ and $k < n$. Let $p^b = \text{tp}^{\text{qf}}(b/A)$. By our assumption there is some $d \in \mathfrak{C}$ such that $d \models h(p^b)_{Pr_0}$. Let us check that $nd + k \models q_{Pr_0}^a(x) \cup \Sigma(x)$.

First, if $\varphi(x, c) \in q_{Pr_0}^a(x)$ (c a tuple from B) then $\mathfrak{C} \models \varphi(a, h^{-1}(c))$ so that $\mathfrak{C} \models \varphi(nb + k, h^{-1}(c))$ so $d \models \varphi(nd + k, c)$ so $nd + k \models \varphi(x, c)$.

Now, suppose that $Pr_n(mx + c) \in \Sigma$.

Then $\mathfrak{C} \models Pr_n(ma + h^{-1}(c))$, so $\mathfrak{C} \models Pr_n(m(nb + k) + h^{-1}(c))$. Hence $m(nb + k) + h^{-1}(c)$ is divisible by n which means that $mk + h^{-1}(c)$ is divisible by n , and as h is an isomorphism (and the language includes 1), so is $mk + c$, hence $m(nd + k) + c$ is also divisible by n . Moreover the quotient $e = [mk + h^{-1}(c)]/n \in A$ maps to $e' = [mk + c]/n \in B$. As $\mathfrak{C} \models Pr(mb + e)$, it follows that $\mathfrak{C} \models Pr(md + e')$, so that $\mathfrak{C} \models Pr_n(m(nd + k) + c)$. The same logic works if $\neg Pr_n(mx + c) \in \Sigma$. ⊣

Divide into cases.

CASE 1: There are infinitely many solutions to $p_{Pr_0}^a$.

Given any finite set $\Sigma \subseteq q_{Pr_0}^a$, it has the form

$$\{Pr(m_i x + c_i) \mid i < k\} \cup \{\neg Pr(m'_j x + c'_j) \mid j < k'\}$$

where $m_i, m'_j \in \mathbb{Z}$ and $c_i, c'_j \in B$ (it also includes formulas of the form $x \neq c$). As¹ $\mathfrak{C} \models \forall x Pr(x) \leftrightarrow Pr(-x)$, we may assume that $m_i, m'_j \geq 1$. Also, it is of course impossible that $(m_i, c_i) = (m'_j, c'_j)$. By \diamond , it is enough to check that $\mathfrak{C} \models \varphi_{\bar{m}}(\bar{c})$ where $\bar{m} = \langle m_i \mid i < k \rangle$ and $\bar{c} = \langle c_i \mid i < k \rangle$ and $\varphi_{\bar{m}}$ is from Remark 2.5. As $\varphi_{\bar{m}}$ is quantifier-free, and as $\mathfrak{C} \models \varphi_{\bar{m}}(h^{-1}(\bar{c}))$ (because $h^{-1}(\Sigma)$ has infinitely many solutions and by \diamond), we are done.

CASE 2: There are only finitely many solutions to p_{Pr_0} .

By \diamond , and as $\mathfrak{C} \models \forall x Pr(x) \leftrightarrow Pr(-x)$, there are some $m_i \geq 1, e_i \in A$ such that $\{Pr(m_i x + e_i) \mid i < k\}$ already has finitely many solutions. Hence $\varphi_{\bar{m}}(\bar{e})$ fails. Let N be from Remark 2.5. We get that for some $p < N$, there is some $i < k$ such that $P_p(m_i a + e_i)$. But as $Pr(m_i a + e_i)$, it must be that $\pm p = m_i a + e_i$. As $\pm p, e_i \in A$, and as A is closed under dividing by m_i , it follows that $a \in A$ and this cannot happen by assumption. ⊣

§3. Decidability and classification. We start with the decidability result that is now almost immediate.

COROLLARY 3.1. *The theory $T_{+,Pr}^*$ is decidable and hence so is $T_{+,Pr}$ provided that Dickson's conjecture holds.*

PROOF. Observing the proof of Lemma 2.7, we see that we can recursively enumerate the axioms that we used. Let us denote this set by Σ . Let Σ' be the complete quantifier-free theory of \mathbb{Z} in L^* . Then Σ' is recursive and contained in $T_{+,Pr}^*$.

Now $\Sigma \cup \Sigma'$ is consistent and complete (every sentence is equivalent to a quantifier free sentence which is decided by Σ'). Hence it is decidable. ⊣

Now we turn to classification in the sense of [14], where one is interested in classifying first-order theories by finding “dividing lines” between them, inducing classes with interesting properties both inside and outside. The most studied such class is that of stable theories, which is a very well-behaved and well-understood class. Containing it is the class of simple theories, and among them the “simplest” simple theories are supersimple of U -rank 1. For the definitions of simple and supersimple theories as well as of forking and dividing, we refer the reader to e.g., [18, Chapter 7, Definition 8.6.3].

¹Here we use the fact that Pr contains both the primes and their negations.

THEOREM 3.2. *Assuming (D), $T_{+,Pr}^*$ (and $T_{+,Pr}$) is supersimple of U -rank 1: working in the monster model \mathfrak{C} , if $\varphi(x, a)$ forks over \emptyset where x is a singleton and a is some tuple from \mathfrak{C} then φ is algebraic (i.e., $\varphi \vdash \bigvee_{i < \kappa} x = c_i$).*

PROOF. The proof is similar to that of Lemma 2.7.

Let N be an ω -saturated model. Suppose that φ forks over \emptyset but is not algebraic. Extend φ to a type $p(x) \in S(N)$ which is nonalgebraic over N . So p forks over \emptyset , and hence it divides over \emptyset by saturation. By quantifier elimination we may assume that p is quantifier free.

Recalling the notation from the proof of Lemma 2.7, we have the following claim.

CLAIM 3.3. *It is enough to prove that for every type $q(x) \in S(N)$, if $q_{Pr} = q \upharpoonright L_{Pr}^*$ divides over \emptyset , then q_{Pr} is algebraic.*

PROOF. We want to show that p is algebraic, thus getting a contradiction. Let $\langle N_i \mid i < \omega \rangle$ be an indiscernible sequence starting with $N_0 = N$ in \mathfrak{C} , which witnesses that p divides.

By indiscernibility, all the congruent conditions in $p(x, N_i)$ (i.e., equations such as $mx + c \equiv d$) are implied by the congruent conditions in $p \upharpoonright \emptyset$. It follows that $\bigcup \{p_{Pr}(x, N_i) \mid i < \omega\} \cup \Sigma$ is inconsistent for some finite $\Sigma \subseteq p$, which is isolated by a formula of the form $x \equiv k \pmod{n}$ for some $k < n$.

Let $c \models p$. Then $c \equiv k \pmod{n}$ for some $k < n$, and let $d = (c - k)/n$. Then $[\text{tp}(d/N)]_{Pr}$ divides over \emptyset as witnessed by the same sequence $\langle N_i \mid i < \omega \rangle$ (let $r = \text{tp}(d/N)$, then if $d' \models \bigcup \{r_{Pr}(x, N_i) \mid i < \omega\}$ then $nd' + k \models \Sigma \cup \bigcup \{p_{Pr}(x, N_i) \mid i < \omega\}$). Hence, $[\text{tp}(d/N)]_{Pr}$ is algebraic, i.e., $d \in N$, but then so is c . \dashv

CLAIM 3.4. *It is enough to prove that for every type $q(x) \in S(N)$, if $q_{Pr_0} = q \upharpoonright L_{Pr_0}^*$ divides over \emptyset , then q_{Pr_0} is algebraic.*

PROOF. This is similar to the proof of Claim 3.3.

By Claim 3.3, it is enough to prove that for any $q(x) \in S(N)$, if q_{Pr} divides over \emptyset then q_{Pr} is algebraic. Suppose that q_{Pr} divides over \emptyset and let $\langle N_i \mid i < \omega \rangle$ be as in the proof of Claim 3.3. There is some finite set of formulas $\Sigma(x, N) \subseteq q_{Pr} \setminus q_{Pr_0}$ such that $\bigcup \{q_{Pr_0}(x, N_i) \cup \Sigma(x, N_i) \mid i < \omega\}$ is inconsistent. As in the proof of Lemma 2.7, we may assume that for some $n \in \mathbb{N}$, Σ consists of formulas of the form $Pr_n(mx + c)$ for $c \in N$ and $m \in \mathbb{Z}$. Let $d \models q$, and assume that $d \equiv k \pmod{n}$ for $k < n$. Then for some $e \in \mathfrak{C}$, $d = ne + k$, and $[\text{tp}(e/N)]_{Pr_0}$ divides over \emptyset (let $r = \text{tp}(e/N)$, then if $e' \models \bigcup \{r_{Pr_0}(x, N_i) \mid i < \omega\}$ then $ne' + k \models \bigcup \{q_{Pr_0}(x, N_i) \cup \Sigma(x, N_i) \mid i < \omega\}$, as in the proof of Lemma 2.7). Hence this type is algebraic and hence so is q . \dashv

CLAIM 3.5. *It is enough to prove that if $\Sigma(x, \bar{c})$ is a finite set of formulas of the form $Pr(mx + c)$ or $\neg Pr(mx + c)$ for $m \in \mathbb{Z}$ and $c \in N$, which has infinitely many solutions, then for any indiscernible sequence $\langle \bar{c}_i \mid i < \omega \rangle$ starting with \bar{c} , $\{\Sigma(x, \bar{c}_i) \mid i < \omega\}$ has infinitely many solutions.*

PROOF. Use Claim 3.4. We have to prove that if q_{Pr_0} divides over \emptyset then it is algebraic. Suppose it is not, and let $\Sigma(x, \bar{c}) \subseteq q_{Pr_0}$ be a finite set of formulas of the form $Pr(mx + c)$ or $\neg Pr(mx + c)$ for $m \in \mathbb{Z}$ and $c \in N$, and let $S \subseteq N$ be finite such that $\Delta(x, \bar{c}, \bar{d}) = \Sigma(x, \bar{c}) \cup \{x \neq d \mid d \in S\}$ divides over \emptyset . Let $\{(\bar{c}_i, \bar{d}_i) \mid i < \omega\}$ be an indiscernible sequence witnessing dividing. But then $\bigcup \{\Sigma(\bar{x}, \bar{c}_i) \mid i < \omega\}$ has

infinitely many solutions by assumption, so by saturation (of \mathfrak{C}) there is a solution which is distinct from $\bigcup \left\{ \bar{d}_i \mid i < \omega \right\}$, contradicting dividing. \dashv

Let $\Sigma(x)$ be as in Claim 3.5.

Then $\Sigma(x, \bar{c}, \bar{c}') = \{Pr(m_i x + c_i) \mid i < k\} \cup \{\neg Pr(m'_j x + c'_j) \mid j < k'\}$, for $m_i, m'_j \in \mathbb{Z}$ and $c_i, c'_j \in N$. Now take an indiscernible sequence $\langle \bar{c}_\alpha, \bar{c}'_\alpha \mid \alpha < \omega \rangle$ starting with $\langle c_i \mid i < k \rangle \frown \langle c'_j \mid j < k' \rangle$. Consider a finite union of the form $\bigcup \{ \Sigma(x, \bar{c}_\alpha, \bar{c}'_\alpha) \mid \alpha < l \}$. Then by indiscernibility it cannot be that $(m_i, c_{i,\alpha}) = (m'_j, c'_{j,\beta})$ for some $\alpha, \beta < l$, $i < k$, and $j < k'$. Hence by (D), it is enough to show that $\star_{\bar{f}}$ holds for $\bar{f} = \langle f_{i,\alpha} \mid i < k, \alpha < l \rangle$ where $f_{i,\alpha}(x) = m_i x + c_{i,\alpha}$, and by Remark 2.5 we have to show that $\varphi_{\bar{m}}(\langle \bar{c}_\alpha \mid \alpha < l \rangle)$ holds.

Let $N \in \mathbb{N}$ be from Remark 2.5 (it depends only on \bar{m} , k and l). We have to check that if $r < N$ is a prime, for some $0 \leq t < r$, for all $i < k$ and $\alpha < l$, $m_i t + c_{i,\alpha} \not\equiv 0 \pmod{r}$. If this does not happen for r , then, as (by indiscernibility) $c_{i,\alpha} \equiv c_i \pmod{r}$, we get that for all $0 \leq t < r$, for some $i < k$, $m_i t + c_i \equiv 0 \pmod{r}$. But this means that Σ cannot have infinitely many solutions by Remark 2.6 — contradiction. \dashv

We move to NIP. We will show that $T_{+,Pr}$ has the independence property IP (and thus the theory is not NIP), and even the n -independence property. This shows in particular that $T_{+,Pr}$ is unstable. We will recall the definition in the proof of Theorem 3.7, but the interested reader may find more in [16] (about NIP) and [4] (on n -dependence).

We will use the following proposition.

PROPOSITION 3.6. *For all $n < \omega$ and $s \subseteq n$ there is an arithmetic progression $\langle at + b \mid t < n \rangle$ of natural numbers such that $at + b$ is prime iff $t \in s$.*

PROOF. As we said in the introduction, according to a private communication with Tamar Ziegler, this follows from the proof of the Green-Tao theorem about arithmetic progression of primes [7].

We give a very detail-free explanation of why this should be true. Heuristically, the primes below N behave like a random set of density $1/\log N$, so the number of $x, d \leq N$ such that $x + d, x + 2d, \dots, x + kd$ are all primes is $N^2 / (\log N)^k$. If we skip the i 'th element in the sequence (i.e., we do not ask it to be prime), then the number is $N^2 / (\log N)^{k-1}$. Hence, we may remove all the prime arithmetic progressions and still find some sequence where the i 'th element is not prime.

We will however give a proof that relies on (D). Fix n and s . Let $b = n! + 1$. Use Lemma 2.3, with the linear maps $x + b, 2x + b, \dots, nx + b$. By Remark 2.1, it is enough to check that for all primes $p \leq n$, for some $t < p$, $kt + b \not\equiv 0 \pmod{p}$ for all $1 \leq k \leq n$. But $b \equiv 1 \pmod{p}$ so this holds for $t = 0$. \dashv

THEOREM 3.7 (Without assuming Dickson's conjecture). *$T_{+,Pr}$ has the independence property and even the n -independence property. Hence so does $T_{+,Pr}^*$.*

PROOF. We use only Proposition 3.6. To prove that T is n -independent, we have to find a formula $\varphi(x, y_1, \dots, y_n)$ such that for all $k < \omega$, there are tuples $a_{i,j}$ for $i < n, j < k$ inside some model $M \models T$ such that for every subset $s \subseteq k^n$, there is some tuple $b_s \in M$ with $M \models \varphi(b_s, a_{0,j_0}, \dots, a_{n-1,j_{n-1}})$ iff $(j_0, \dots, j_{n-1}) \in s$. This of course implies the independent property.

The formula we take is $\varphi(x, y_1, \dots, y_n) = Pr(x + y_1 + \dots + y_n)$, and we work in \mathbb{Z} .

Given k , by Proposition 3.6 there is an arithmetic progression of length $k^n \cdot 2^{(k^n)}$, which we write as $\langle \bar{c}_s \mid s \subseteq k^n \rangle$ where $\bar{c}_s = \langle c_{s,l} \mid l < k^n \rangle$, such that for each subset $s \subseteq k^n$ and $l < k^n$, $Pr(c_{s,l})$ iff $(j_0, \dots, j_{n-1}) \in s$ where $j_i < k$ are (unique) such that $l = \sum_{i < n} j_i k^i$.

Suppose this progression has difference $d > 0$. Now we choose $a_{i,j}$ for $i < n, j < k$ and b_s for $s \subseteq k^n$ as follows.

Let $a_{0,j} = j \cdot d$ for $j < k$ and in general, for $i < n, a_{i,j} = j \cdot d \cdot k^i$. Let $b_s = c_{s,0}$.

Now note that

$$c_{s,0} + \sum_{i < n} (j_i d) k^i = c_{s, \sum_{i < n} j_i \cdot k^i}.$$

And so we are done. ⊖

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