

## ULTRA-SMALL SCALE-FREE GEOMETRIC NETWORKS

J. E. YUKICH,\* *Lehigh University*

### Abstract

We consider a family of long-range percolation models  $(G_p)_{p>0}$  on  $\mathbb{Z}^d$  that allow dependence between edges and have the following connectivity properties for  $p \in (1/d, \infty)$ : (i) the degree distribution of vertices in  $G_p$  has a power-law distribution; (ii) the graph distance between points  $\mathbf{x}$  and  $\mathbf{y}$  is bounded by a multiple of  $\log_{pd} \log_{pd} |\mathbf{x} - \mathbf{y}|$  with probability  $1 - o(1)$ ; and (iii) an adversary can delete a relatively small number of nodes from  $G_p(\mathbb{Z}^d \cap [0, n]^d)$ , resulting in two large, disconnected subgraphs.

*Keywords:* Scale-free graph; long-range percolation; chemical distance

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### 1. Introduction

The statistical properties of large networks have received considerable attention in the recent scientific literature [2], [14], [21], [25]. Of special interest are the power-law random networks in which the fraction of vertices of degree  $k$  is proportional to  $k^{-q}$  for some  $q > 0$ . Such networks lack an inherent scale and have been termed ‘scale free’. Scale-free graphs are ubiquitous in random network theory and have been proposed as a way to model the behavior of technological, social, and biological networks [1], [21].

Networks often have a geometric component to them where the vertices have positions in space and geographic proximity plays a role in deciding which vertices get connected. In this context, random geometric graphs are a natural alternative to the classical Erdős–Rényi random graph models. Random connection models [20] provide one way to describe networks with spatial content. In these models the event,  $E_{\mathbf{x},\mathbf{y}}$ , of a connection between points  $\mathbf{x}$  and  $\mathbf{y}$  has probability  $p_{\mathbf{x},\mathbf{y}} := \mathbb{P}[E_{\mathbf{x},\mathbf{y}}] = g(|\mathbf{x} - \mathbf{y}|)$ , where  $g: \mathbb{R}^+ \rightarrow [0, 1]$  is a connection function and  $|\mathbf{x}|$  denotes the Euclidean norm of  $\mathbf{x}$ . The standard long-range percolation model assumes independence of  $E_{\mathbf{x},\mathbf{y}}$  and  $E_{\mathbf{x},\mathbf{u}}$ ,  $\mathbf{y} \neq \mathbf{u}$ , which may not be the case in networked systems. Moreover, the degree distribution in this connection model generally does not follow a power law.

Allowing dependency between edges will in general result in technically more complicated models. In this note we show that a natural edge dependency gives rise to a family of long-range percolation models,  $(G_p)_{p>0}$ , which is technically tractable and which exhibits three connectivity properties for  $p \in (1/d, \infty)$ . First,  $G_p$  has a power-law distribution. Second,  $G_p$  is ultra-small, in the sense that the graph distance between lattice points  $\mathbf{x}$  and  $\mathbf{y}$  is bounded by a multiple of  $\log_{pd} \log_{pd} |\mathbf{x} - \mathbf{y}|$  with probability  $1 - o(1)$ , where  $o(1)$  denotes a quantity tending to 0 as  $|\mathbf{x} - \mathbf{y}| \rightarrow \infty$ . Ultra-small graph distances imply efficiency, are consistent

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\* Postal address: Department of Mathematics, Lehigh University, Bethlehem, PA 18015, USA.

Email address: joseph.yukich@lehigh.edu

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with the ‘small-world phenomenon’ [2], [14], [24], [25], and are relevant in the context of routing, searching, and transport of information. Third, an adversary can delete a relatively small number of nodes from  $G_p(\mathbb{Z}^d \cap [0, n]^d)$ , after which there are two disconnected subgraphs, each containing nearly one-half of the total number of network nodes.

**1.1. A general dependent random connection model**

Let  $\{U_z\}_{z \in \mathbb{Z}^d}$  be independent, identically distributed uniform[0, 1] random variables indexed by  $\mathbb{Z}^d$ . Let  $p > 0$  and  $\delta \in (0, 1]$ . For each  $z \in \mathbb{Z}^d$ , we take  $\delta U_z^{-p}$  to represent a weight at node  $z$  defining the radius of the ‘ball of influence’ at  $z$ . Consider the graph  $G_{p,\delta} := G_{p,\delta}(\mathbb{Z}^d)$  which puts an edge between nodes  $x, y \in \mathbb{Z}^d$  whenever each node is contained in the other’s ball of influence. Thus, this connection rule says that the edge  $(x, y)$  appears in  $G_{p,\delta}(\mathbb{Z}^d)$  whenever

$$|x - y| \leq \delta \min(U_x^{-p}, U_y^{-p}). \tag{1.1}$$

Let  $\delta = 1$ . By the independence of the  $U_z$ , we have  $p_{x,y} := P[E_{x,y}] = |x - y|^{-2/p}$ , showing that the probability of (there being) long edges in  $G_p := G_{p,1}$  increases with  $p$ . Edges in  $G_p$  have dependent probabilities: if  $|y| < |x|$  then the probability of the edge  $(0, y)$  given the edge  $(0, x)$  is  $|y|^{-1/p}$  instead of  $|y|^{-2/p}$ .

The family of random connection models  $G_{p,\delta}$  is disconnected for general  $p$  and  $\delta$ , but not for  $\delta = 1$ , since having  $U_z^{-p} \geq 1$  for all  $z \in \mathbb{Z}^d$  implies that adjacent lattice points are connected in  $G_p$ . The main results below show, for all  $p \in (1/d, \infty)$ , that the components of  $G_p$  are of arbitrarily large diameter with arbitrarily large probability. Moreover, in accordance with their Poisson Boolean model counterparts (cf. [20]), it is easy to check, for all  $\delta \in (0, 1]$  and large  $p$ , that the expected number of nodes in the component of  $G_{p,\delta}$  containing  $0$  is infinite, whereas, for  $p$  and  $\delta$  both small, the expected number of such nodes is finite. Our purpose here is to explore the connectivity properties of  $G_p, p \in (1/d, \infty)$ .

**1.2. Main results**

Let  $D_p(0)$  denote the degree of the origin in  $G_p(\mathbb{Z}^d)$ , let  $\omega_d$  denote the volume of the unit-radius ball in  $\mathbb{R}^d$ , and let  $\alpha := pd - 1$ . Our first result shows that if  $p \in (1/d, \infty)$  then the degree of a typical vertex follows a power law, i.e.  $G_p$  is scale free.

**Theorem 1.1.** ( $G_p(\mathbb{Z}^d)$  has a power-law degree distribution.) *For all  $d = 1, 2, \dots$  and all  $p \in (1/d, \infty)$ ,*

$$\lim_{t \rightarrow \infty} t^{1/\alpha} P[D_p(0) > t] = (pd\omega_d/\alpha)^{1/\alpha}.$$

For all  $x, y \in \mathbb{Z}^d$ ,  $d_p(x, y)$  denotes the  $G_p$  graph distance (‘chemical distance’) between  $x$  and  $y$ . Our next result says that  $G_p$  is ultra-small (cf. [12]), in that  $d_p(x, y)$  is bounded by  $4(2 + \log \log |x - y|)$  with probability  $1 - o(1)$ , where throughout, for all  $s > 0$ ,  $\log s$  is short for  $\log_{pd} s$ . We expect that the upper bound in this result can be improved but have not tried to obtain the sharpest bound.

**Theorem 1.2.** ( $G_p(\mathbb{Z}^d)$  has small graph distance.) *For all  $d = 1, 2, \dots$  and all  $p \in (1/d, \infty)$ ,*

$$\frac{d_p(0, x)}{2 + \log \log |x|} \leq 4$$

*with probability  $1 - o(1)$ , where  $o(1)$  tends to 0 as  $|x| \rightarrow \infty$ .*

The network failure of  $G_p(\mathbb{Z}^d)$  is easily quantified, as follows.

**Theorem 1.3.** (Network failure.) *For all  $d = 1, 2, \dots$  and all  $p \in (1/d, \infty)$ , an adversary can delete  $N$  nodes from  $G_p(\mathbb{Z}^d \cap [0, n]^d)$ , where  $E[N] = O(n^{d-1}[n^{1-1/p} \vee 1])$ , resulting in two disconnected subgraphs on vertex sets of cardinality at least  $n^d/2 - N$ .*

Theorem 1.3 implies, in particular, that if  $p \in (1/d, 1)$  then removing roughly  $O(n^{d-1})$  nodes may reduce  $G_p(\mathbb{Z}^d \cap [0, n]^d)$  to two large, disconnected subgraphs.

**Remarks.** 1. *Standard long-range percolation models.* Assume that  $p_{x,y} := P[E_{x,y}] = |x - y|^{-s+o(1)}$  as  $|x - y| \rightarrow \infty$ , for some constant  $s \in (0, \infty)$ ;  $E_{x,y}$  and  $E_{x,u}$  are independent for all  $x, y, u \in \mathbb{Z}^d$ . For  $s \in (0, d)$ , Benjamini *et al.* [4] showed that the graph distance  $d(\mathbf{0}, \mathbf{x})$  behaves like the constant  $\lceil s/(d - s) \rceil$  as  $|\mathbf{x}| \rightarrow \infty$ . Here,  $\lceil x \rceil$  denotes the greatest integer less than  $x$ . For  $s = d$ , Coppersmith *et al.* [13] showed that  $d(\mathbf{0}, \mathbf{x})$  scales as  $\log |\mathbf{x}|/\log \log |\mathbf{x}|$ , whereas, for  $s \in (d, 2d)$ , Biskup [7], [8] showed that  $d(\mathbf{0}, \mathbf{x})$  scales as  $(\log |\mathbf{x}|)^{\Delta+o(1)}$ , where  $\Delta := \Delta(s, d) := \log 2/\log(2d/s)$ . The case  $s = 2d$  is open and, for  $s \in (2d, \infty)$ ,  $d(\mathbf{0}, \mathbf{x})$  scales at least linearly in  $|\mathbf{x}|$ , as shown by Berger [5]. The different scalings for the standard long-range percolation model suggest that  $G_p$  also has different scalings for  $p \in (0, 1/d)$ , but we have not determined them. Kleinberg [19] proposed a lattice model where long-range contacts are added in a biased way, there being, however, a uniform bound on the number of such contacts.

2. *Geometric networks in  $\mathbb{R}^d$ .* We expect that Theorems 1.1–1.3 extend to analogously defined continuum models on Poisson point sets in  $\mathbb{R}^d$ . This would add to the following related results.

- (a) Let  $f: \mathbb{R}^d \rightarrow \mathbb{R}^+$  and let  $\mathcal{P}_f$  be a Poisson point process on  $\mathbb{R}^d$  with intensity  $f$ . The *geometric graph*, described in depth by Penrose [23], joins two nodes in  $\mathcal{P}_f$  whenever their Euclidean distance is less than a specified cutoff. Hermann *et al.* [18, Section II.B] showed that if  $\int_{\mathbb{R}^d} f^r(\mathbf{x}) \, d\mathbf{x} = \infty$  for all  $r > r_0$ , then the degree distribution is effectively a power law.
- (b) The *on-line nearest-neighbors graph* is defined on randomly ordered point sets in  $\mathbb{R}^d$ , and places an edge between each point and its nearest neighbor amongst the points preceding it in the ordering. Such graphs have scale-free properties over certain degree domains [6], [16].
- (c) Franceschetti and Meester [17] developed a scale-free continuum model but did not obtain iterated log bounds on interpoint graph distances.
- (d) The standard *Boolean connection model* puts an edge between  $\mathbf{x}$  and  $\mathbf{y}$  whenever the respective balls of influence overlap. In the context of (1.1),  $(\mathbf{x}, \mathbf{y})$  is an edge whenever  $|\mathbf{x} - \mathbf{y}| \leq \delta(U_{\mathbf{x}}^{-p} + U_{\mathbf{y}}^{-p})$ . These models are not in general scale free.

3. *Power exponents  $q \in (2, 3)$ .* Consider a random graph on  $n$  nodes  $v_1, v_2, \dots, v_n$  with weight (expected degree)  $w_i$  at node  $v_i$ . Nodes  $v_i$  and  $v_j$  are connected with probability  $\rho w_i w_j$ , where  $\rho = (\sum_{i=1}^n w_i)^{-1}$ . Chung and Lu [10], [11] provided conditions on the weights under which the degree distribution is proportional to  $k^{-q}$ ,  $q \in (2, 3)$ ,  $k \in \mathbb{Z}$ , the average distance between nodes is almost surely  $O(\log \log n)$ , and the diameter is  $O(\log n)$ . In unrelated work, Cohen and Havlin [12] argued that whenever the degree distribution of a random graph on  $n$  vertices is proportional to  $k^{-q}$ , where  $q \in (2, 3)$ ,  $k$  is restricted to  $(m, K)$ , and where  $m$  and  $K := K(n)$  are well-defined ‘cutoffs’, then the diameter behaves like  $\log \log n$ .

4. *Preferential attachment models.* These dynamic graphs evolve with time in such a way that a newly arriving vertex connects to an existing vertex with a probability proportional to the degree of the (latter) vertex. Thus, nodes of high degree tend to acquire more new links than do nodes of low degree. Albert and Barabási [1] showed that such models follow a power law, are not geometry dependent, and, as shown by Bollabás and Riordan [9], are not ultra-small in general.

5. *Degree dependence on  $p$ .* Theorem 1.1 tells us that  $P[D_p(\mathbf{0}) = k] \sim Ck^{-q}$ , where  $q := pd/(pd - 1)$ . Thus, as  $p$  increases on  $(1/d, \infty)$ , the exponent of the degree distribution,  $q$ , decreases to 1.

6. *Further connectivity results.* Theorems 1.1–1.3 describe the connectivity of  $G_p(\mathbb{Z}^d)$ . Further analysis of the connectivity of  $G_p(\mathbb{Z}^d)$ , such as thermodynamic and Gaussian limits for the number of three cycles (or other clustering coefficients) on  $G_p(\mathbb{Z}^d \cap [0, n]^d)$ , is simplified by appealing to the stabilization properties of  $G_p$  (see especially [22]).  $G_p(\mathbb{Z}^d)$  is *assortative* in that high-degree nodes tend to link to high-degree nodes and low-degree nodes tend to link to low-degree nodes.

7. *The case  $p \in (0, 1/d)$ .* If  $p \in (0, 1/d)$  then  $G_p$  has few long edges and the proofs of the scale-free and ultra-small properties break down. The scalar  $1/d$  thus represents the boundary between graphs that are ultra-small scale free and those which are not.

### 2. Proof of Theorem 1.1

Throughout, we adopt the following notation:  $B_r(\mathbf{x})$  denotes the Euclidean ball of radius  $r$  centered at  $\mathbf{x} \in \mathbb{R}^d$ ,  $L_r(\mathbf{x}) := B_r(\mathbf{x}) \cap \mathbb{Z}^d \setminus \{\mathbf{x}\}$  denotes the lattice points a distant at most  $r$  from  $\mathbf{x}$ , and  $C$  denotes a generic positive constant whose value may change from line to line. The underlying probability space is  $\Omega := [0, 1]^{\mathbb{Z}^d}$  and is equipped with the product probability measure  $P := \mu^{\mathbb{Z}^d}$ , where  $\mu$  is the uniform probability measure on  $[0, 1]$ . Conditional on  $U_{\mathbf{0}} = u$ ,  $D_p(\mathbf{0})$  is the number of points  $\mathbf{y}$  in  $L_{u^{-p}}(\mathbf{0})$  with weight,  $U_{\mathbf{y}}^{-p}$ , exceeding  $|\mathbf{y}|$ ; hence,  $U_{\mathbf{y}} \in [0, |\mathbf{y}|^{-1/p}]$ . Writing  $D(u^{-p})$  for the value of  $D_p(\mathbf{0})$  conditioned on  $\mathbf{0}$  having weight  $u^{-p}$ , we have

$$D(u^{-p}) = \sum_{\mathbf{y} \in L_{u^{-p}}(\mathbf{0})} 1_{\{U_{\mathbf{y}} \leq |\mathbf{y}|^{-1/p}\}}.$$

Thus, to prove Theorem 1.1 we condition on  $U_{\mathbf{0}}$  and show that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_0^1 P[D(u^{-p}) > t] du = \left(\frac{pd\omega_d}{\alpha}\right)^{1/\alpha}, \tag{2.1}$$

where, recall,  $\alpha := pd - 1$ . The next lemma will be useful in establishing (2.1). Let  $\beta := pd\omega_d/\alpha$ .

**Lemma 2.1.** *For all  $p \in (1/d, \infty)$ , we have*

$$E[D(u^{-p})] = \beta u^{-\alpha} + O(\max(1, u^{-pd+p+1})), \tag{2.2}$$

where the error on the right-hand side of (2.2) holds as  $u \rightarrow 0^+$ .

*Proof.* Note that  $E[D(u^{-p})]$  is approximated by

$$\int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} = d\omega_d \int_0^{u^{-p}} t^{d-1-1/p} dt = \beta u^{-\alpha}.$$

Let  $R := R(u)$  be the maximal collection of grid cubes (cubes centered at points in  $\mathbb{Z}^d$  with edge length 1) contained within  $B_{u^{-p}}(\mathbf{0})$ . The approximation error

$$\left| E[D(u^{-p})] - \int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|$$

is bounded by the sum of the following three errors:

$$\begin{aligned} E_1 &:= \left| E[D(u^{-p})] - \sum_{\mathbf{y} \in R(u) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p} \right|, \\ E_2 &:= \left| \sum_{\mathbf{y} \in R(u) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p} - \int_{R(u)} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|, \\ E_3 &:= \left| \int_{R(u)} |\mathbf{x}|^{-1/p} d\mathbf{x} - \int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|. \end{aligned}$$

Now,

$$E_1 = \sum_{\mathbf{y} \in (B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p}$$

and, so, is bounded by the product of

$$\text{card}\{(B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d\} \quad \text{and} \quad \sup\{|\mathbf{y}|^{-1/p} : \mathbf{y} \in (B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d\}.$$

Since the first factor is bounded by  $Cu^{-p(d-1)}$  and the second by  $Cu$ , it follows that  $E_1 \leq Cu^{-pd+p+1}$ . A similar method shows that  $E_3 \leq Cu^{-pd+p+1}$ .

We estimate  $E_2$  as follows. For all  $\mathbf{y} \in \mathbb{Z}^d$ , let  $Q_{\mathbf{y}}$  denote the grid cube with center  $\mathbf{y}$ . For all  $s = 1, 2, \dots$ , let  $M(s) := \text{card}\{\mathbf{y} \in \mathbb{Z}^d : |\mathbf{y}| \in [s, s+1)\}$ . Since there is a constant  $C > 0$  such that, for all  $\mathbf{x} \in Q_{\mathbf{y}}$  and all  $\mathbf{y} \in \mathbb{Z}^d$ ,

$$||\mathbf{y}|^{-1/p} - |\mathbf{x}|^{-1/p}| \leq C|\mathbf{y}|^{-1/p-1},$$

it follows that

$$E_2 \leq C \sum_{s=1}^{u^{-p}} s^{-1/p-1} M(s) \leq C \sum_{s=1}^{u^{-p}} s^{-1/p+d-2} \leq C \max(1, u^{-pd+p+1}),$$

since  $M(s) \leq Cs^{d-1}$ . Combining the bounds for  $E_1, E_2,$  and  $E_3$  yields Lemma 2.1.

Letting  $s := u^{-p}$  in (2.1), note that, to prove Theorem 1.1, it suffices to show that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_1^\infty \mathbb{P}[D(s) > t] \frac{1}{p} s^{-1/p-1} ds = \beta^{1/\alpha}. \tag{2.3}$$

We observe that (2.3) is plausible because Lemma 2.1 suggests that  $\mathbb{P}[D(s) > t]$  is close to 1 for  $t \ll \beta s^{\alpha/p}$  and close to 0 for  $t \gg \beta s^{\alpha/p}$ , indicating that the left-hand side of (2.3) behaves like

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{(t/\beta)^{p/\alpha}}^\infty \frac{1}{p} s^{-1/p-1} ds = \beta^{1/\alpha}.$$

To put this heuristic argument on a rigorous footing, we rewrite the integral in (2.3) as a sum of two integrals. The first integral is estimated via Bernstein’s inequality and the second is handled using Poisson approximation arguments. We do this as follows.

For all  $v > 0$ , let  $m(v) := \sup\{s : E[D(s)] \leq v\}$ . Recalling that  $\alpha := pd - 1$ , from Lemma 2.1 we obtain

$$E[D(s)] = \beta s^{\alpha/p} + O(\max(1, s^{d-1-1/p})) = \beta s^{\alpha/p}(1 + \max(O(s^{1/p-d}), O(s^{-1}))). \tag{2.4}$$

It follows, for large  $v$  and  $p \in (1/d, \infty)$ , that

$$m(v) = \left( \frac{v}{(1 + o(1))\beta} \right)^{p/\alpha},$$

where  $o(1)$  tends to 0 as  $v \rightarrow \infty$ . Given fixed  $t \geq \beta$  and  $\varepsilon \in (0, \frac{1}{2})$ , define the following two integration domains:

$$I_1 := [1, m(t - t^{1/2+\varepsilon})], \quad I_2 := [m(t - t^{1/2+\varepsilon}), \infty).$$

Rewrite the left-hand side of (2.3) as

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{I_1} P[D(s) > t] \frac{1}{p} s^{-1/p-1} ds + \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{I_2} P[D(s) > t] \frac{1}{p} s^{-1/p-1} ds =: S_1 + S_2,$$

provided that both limits exist.

To prove Theorem 1.1 it suffices to show that  $S_1 = 0$  and  $S_2 = \beta^{1/\alpha}$ . We first show that  $S_1 = 0$ . Bernstein’s inequality [15, p. 12] for sums of independent, bounded random variables yields, for all  $s \in I_1$ ,

$$P[D(s) > t] \leq \exp\left(\frac{-(t - E[D(s)])^2}{2 E[D(s)] + 4t/3}\right).$$

Using the bounds  $\inf_{s \in I_1} (t - E[D(s)]) \geq t^{1/2+\varepsilon}$  and  $\sup_{s \in I_1} E[D(s)] \leq t - t^{1/2+\varepsilon} < t$ , for all  $s \in I_1$  we thus obtain

$$P[D(s) > t] \leq \exp\left(\frac{-(t^{1/2+\varepsilon})^2}{10t/3}\right) = \exp\left(-\frac{3t^{2\varepsilon}}{10}\right).$$

It follows that

$$S_1 \leq \limsup_{t \rightarrow \infty} t^{1/\alpha} \exp\left(-\frac{3t^{2\varepsilon}}{10}\right) \int_1^\infty \frac{1}{p} s^{-1/p-1} ds = 0.$$

We next show that  $S_2 = \beta^{1/\alpha}$ . By approximating  $D(s)$  with a Poisson random variable we establish the following simplified expression for  $S_2$ . Here and elsewhere,  $Po(\lambda)$  denotes a Poisson random variable with mean  $\lambda$ .

**Lemma 2.2.** *For all  $p \in (1/d, \infty)$ , we have*

$$S_2 = \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^\infty P[Po(E[D(s)]) > t] \frac{1}{p} s^{-1/p-1} ds.$$

*Proof.* For all  $y \in \mathbb{Z}^d$ , let  $p_y := E[1_{\{U_y \leq |y|^{-1/p}\}}] = |y|^{-1/p}$ . Letting  $d_{TV}$  be the total variation distance, it follows from well-known Poisson approximation bounds (e.g. Equation (1.23)

of [3]) that

$$d_{TV}(D(s), \text{Po}(E[D(s)])) \leq \left( \sum_{y \in L_s(\mathbf{0})} p_y \right)^{-1} \sum_{y \in L_s(\mathbf{0})} p_y^2.$$

By an analysis similar to that in the proof of Lemma 2.1 and (2.4), for  $d > 2/p$  we obtain

$$\sum_{y \in L_s(\mathbf{0})} p_y^2 = \frac{pd\omega_d}{pd - 2} s^{d-2/p} (1 + o(1)),$$

whereas, for  $1/p < d \leq 2/p$ , we have

$$\sum_{y \in L_s(\mathbf{0})} p_y^2 = O(1).$$

It follows from Lemma 2.1 that, for  $d > 2/p$ , we obtain

$$\begin{aligned} d_{TV}(D(s), \text{Po}(E[D(s)])) &\leq (\beta s^{d-1/p} (1 + o(1)))^{-1} \beta \left( \frac{pd - 1}{pd - 2} \right) s^{d-2/p} (1 + o(1)) \\ &= O(s^{-1/p}), \end{aligned}$$

whereas, for  $1/p < d \leq 2/p$ , we have

$$d_{TV}(D(s), \text{Po}(E[D(s)])) = O(s^{-d+1/p}).$$

Letting

$$e(s, t) := P[D(s) > t] - P[\text{Po}(E[D(s)]) > t],$$

it follows that, uniformly in  $t \in (0, \infty)$ , we have  $|e(s, t)| = O(s^{-\xi})$ , where  $\xi = 1/p$  for  $d > 2/p$  and  $\xi = d - 1/p$  for  $1/p < d \leq 2/p$ . We now rewrite  $S_2$  as

$$S_2 = \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^{\infty} (P[\text{Po}(E[D(s)]) > t] + e(s, t)) \frac{1}{p} s^{-1/p-1} ds$$

and show that the term containing  $e(s, t)$  is negligible.

Recall that

$$m(t - t^{1/2+\varepsilon}) = \left( \frac{t - t^{1/2+\varepsilon}}{(1 + o(1))\beta} \right)^{p/\alpha},$$

where, here and in the remainder of this section,  $o(1)$  tends to 0 as  $t \rightarrow \infty$ . It follows that

$$\int_{m(t-t^{1/2+\varepsilon})}^{\infty} e(s, t) s^{-1/p-1} ds = O\left( \int_{m(t-t^{1/2+\varepsilon})}^{\infty} s^{-\xi-1/p-1} ds \right) = O(t^{-p/\alpha(\xi+1/p)})$$

and, therefore, that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^{\infty} e(s, t) s^{-1/p-1} ds = 0.$$

We thus obtain Lemma 2.2.

It is now straightforward to show that  $S_2 = \beta^{1/\alpha}$ . Letting  $z := \beta s^{d-1/p}/t$ , whence  $s = (tz/\beta)^{p/\alpha}$  and  $E[D(s)] = tz(1 + O((tz)^{-\rho}))$  with  $\rho := \rho(p, d) > 0$ , we obtain, via Lemma 2.2,

$$S_2 = \lim_{t \rightarrow \infty} \frac{\beta^{1/\alpha}}{\alpha} \int_{1+o(1)}^{\infty} P[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] z^{-1/\alpha-1} dz.$$

The integrability of the integrand on  $[1 + o(1), \infty)$  gives, for all  $\gamma > 0$ ,

$$S_2 = \lim_{t \rightarrow \infty} \frac{\beta^{1/\alpha}}{\alpha} \int_{1+\gamma}^{\infty} \mathbb{P}[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] z^{-1/\alpha-1} dz + \gamma \cdot O(1).$$

For all  $z \in [1 + \gamma, \infty)$ , we have  $\mathbb{P}[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] \rightarrow 1$  as  $t \rightarrow \infty$ . The dominated convergence theorem yields

$$S_2 = \frac{\beta^{1/\alpha}}{\alpha} \int_1^{\infty} z^{-1/\alpha-1} dz + \gamma O(1) = \beta^{1/\alpha} + \gamma \cdot O(1).$$

Now let  $\gamma \rightarrow 0$  to obtain  $S_2 = \beta^{1/\alpha}$ , as desired.

### 3. Proof of Theorem 1.2

We prove Theorem 1.2 by showing, for all  $\mathbf{x} \in \mathbb{Z}^d$ , the existence of an event  $E := E(\mathbf{x}) \subset \Omega$ , with  $\mathbb{P}[E] = 1 - o(1)$ , such that on  $E$  there is a path  $\pi$  consisting of  $N$  edges in  $G_p(\mathbb{Z}^d)$  joining  $\mathbf{0}$  to  $\mathbf{x}$ , where  $N \leq 4(2 + \log \log |\mathbf{x}|)$ . Here and in the sequel,  $o(1)$  denotes a quantity tending to 0 as  $|\mathbf{x}| \rightarrow \infty$ .

Constructing the path  $\pi$  would be easy if the balls of influence at  $\mathbf{0}$  and  $\mathbf{x}$  both had radius at least  $|\mathbf{x}|$ , for then  $\pi$  would consist merely of the single edge  $(\mathbf{0}, \mathbf{x})$ . In general, the balls of influence at  $\mathbf{0}$  and  $\mathbf{x}$  have much smaller radii and the path  $\pi$  thus needs to join a sequence of balls such that consecutive balls contain each other’s centers.

The heart of the proof will consist of constructing a sequence of nodes of cardinality roughly  $2 \log \log |\mathbf{x}|$  with these properties: the first node,  $\mathbf{0}'$ , is at distance at most  $\frac{1}{2} \log \log |\mathbf{x}|$  from  $\mathbf{0}$ ; the last node,  $\mathbf{x}'$ , is at distance at most  $\frac{1}{2} \log \log |\mathbf{x}|$  from  $\mathbf{x}$ ; and the edges defined by consecutive nodes are in  $G_p$ , i.e. the balls of influence at consecutive nodes contain each other’s centers. Since  $\mathbf{0}$  and  $\mathbf{0}'$  can be joined with a path of at most  $\log \log |\mathbf{x}|$  edges, and likewise for  $\mathbf{x}$  and  $\mathbf{x}'$ , we can obtain a path  $\pi$  consisting of roughly  $4 \log \log |\mathbf{x}|$  edges. The construction of this sequence of nodes depends critically on an intermediate node, denoted here by  $\mathbf{P}_0$ , that has an unusually large ball of influence. Before defining  $\mathbf{0}'$ ,  $\mathbf{P}_0$ , and  $\mathbf{x}'$ , we need some terminology.

For all  $\mathbf{x} \in \mathbb{R}^d$  and  $r > 0$ , let  $L_r^+(\mathbf{x})$  and  $L_r^-(\mathbf{x})$  denote the lattice points in the upper and lower hemispheres of radius  $r$  centered at  $\mathbf{x}$ . That is,  $L_r^+(\mathbf{x}) := B_r(\mathbf{x}) \cap (\mathbb{Z}^{d-1} \times \mathbb{Z}^+)$  and, similarly,  $L_r^-(\mathbf{x}) := B_r(\mathbf{x}) \cap (\mathbb{Z}^{d-1} \times \mathbb{Z}^-)$ . Here  $\mathbb{Z}^+ := \{1, 2, \dots\}$  and  $\mathbb{Z}^- := \{-1, -2, \dots\}$ .

#### 3.1. Definition of $\mathbf{0}'$ , $\mathbf{P}_0$ , and $\mathbf{x}'$

Throughout, we appeal to the following elementary fact. Recall that  $\log s$  is short for  $\log_{pd} s$ .

**Lemma 3.1.** *Let  $U_1, \dots, U_n$  be independent and identically uniformly distributed on  $[0, 1]$ . Then, for all  $n > pd$ , we have*

$$\min_{i \leq n} U_i \leq \frac{K \log n}{n}$$

with probability at least  $1 - n^{-K}$ .

In the sequel, we fix  $K$  to be large, with a value to be determined later.

3.1.1. *Definition of  $\mathbf{0}'$ .* Let  $E_0 := E_0(\mathbf{x})$  be the event that there is a node  $z \in L_{(1/2) \log \log |\mathbf{x}|}^-(\mathbf{0})$  such that

$$U_z \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^d}.$$



Clearly,  $E_0$  depends only on  $U_z, z \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{0})$ . In case more than one node in  $L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{0})$  satisfies the last bound, we choose  $z$  to be that node with smallest lexicographical order.

By Lemma 3.1,  $P[E_0] \geq 1 - C(\log\log|\mathbf{x}|)^{-dK}$ . Given  $E_0$  we let  $\theta' := z$ . Note that  $\theta'$  is random and that, since  $pd > 1$ , for all large  $|\mathbf{x}|$  we have

$$U_{\theta'}^{-p} \geq 2 \log\log|\mathbf{x}|. \tag{3.1}$$

Inequality (3.1) will be important in the sequel. For now note that, since  $G_p(\mathbb{Z}^d)$  connects adjacent lattice points, it follows that  $d_p(\mathbf{y}, \mathbf{x}) \leq 2|\mathbf{y} - \mathbf{x}|$  for all  $\mathbf{x}, \mathbf{y} \in \mathbb{Z}^d$ , i.e. that

$$d_p(\mathbf{0}, \theta') \leq \log\log|\mathbf{x}|. \tag{3.2}$$

3.1.2. *Definition of  $\mathbf{x}'$ .* Similarly, given  $\mathbf{x}$ , with probability at least  $1 - C(\log\log|\mathbf{x}|)^{-dK}$  there is an event  $E_x$  such that there is a node  $\mathbf{x}' \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{x})$  on  $E_x$  with weight

$$U_{\mathbf{x}'}^{-p} \geq 2 \log\log|\mathbf{x}|.$$

Clearly  $d_p(\mathbf{x}, \mathbf{x}') \leq \log\log|\mathbf{x}|$  and  $E_x$  depends only on  $U_z, z \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{x})$ .

3.1.3. *Definition of  $P_0$ .* Assume without loss of generality that the components of  $\mathbf{x}$  have even parity, meaning that  $\mathbf{x}/2 \in \mathbb{Z}^d$ . Consider the event,  $E_{\mathbf{x}/2}$ , that there is a node  $P_0 \in L_{|\mathbf{x}|/10}(\mathbf{x}/2)$  with

$$U_{P_0} \leq \frac{K \log(|\mathbf{x}|)^d}{|\mathbf{x}|^d}.$$

Lemma 3.1 implies that  $P[E_{\mathbf{x}/2}] \geq 1 - C(|\mathbf{x}|^{-dK})$ . We note that, for large  $|\mathbf{x}|$ ,

$$U_{P_0}^{-p} \geq 2|\mathbf{x}| \tag{3.3}$$

since  $pd > 1$ .

**3.2. Construction of the path  $\pi$  via  $\theta', P_0$ , and  $\mathbf{x}'$**

It will suffice to show that there is an event  $E := E(\mathbf{x})$ , with  $P[E(\mathbf{x})] = 1 - o(1)$ , such that on  $E$  there are two paths, each having at most  $2 + 2\lceil\log\log|\mathbf{x}|\rceil$  edges, with one path joining  $P_0$  to  $\mathbf{0}$  and the other joining  $P_0$  to  $\mathbf{x}$ . It will be enough to show the existence of a path between  $P_0$  and  $\mathbf{0}$ , for the method can be repeated verbatim to yield the path between  $P_0$  and  $\mathbf{x}$ . We first introduce some additional terminology.

We abbreviate our notation by letting  $b := pd$ . Note that  $b > 1$  by assumption. Fix  $\varepsilon \in (0, 1)$  and  $\mathbf{x} \in \mathbb{Z}^d$  with  $|\mathbf{x}|$  large. For all  $j = 1, 2, \dots$ , let

$$r_j := r_j(\mathbf{x}, \varepsilon) := |\mathbf{x}|^{b^{-j(1-\varepsilon)}}$$

and note that  $r_j \downarrow 1$  and  $1 < r_j < |\mathbf{x}|$  for all  $j = 1, 2, \dots$ . We record an elementary fact.

**Lemma 3.2.**  $r_{j+1} = r_j^{\beta(p,d,\varepsilon)}$ , where  $\beta(p, d, \varepsilon) := b^{-1+\varepsilon}$ .

For all  $j = 1, 2, \dots$ , consider the following disjoint ‘semi-annular’ regions of lattice points:

$$A_j := [(L_{r_j}^+(\theta') - L_{r_{j+1}}^+(\theta')) \setminus L_{|\mathbf{x}|/10}^+(\mathbf{x}/2)].$$

The construction of the path joining  $P_0$  to  $\mathbf{0}$  is facilitated by the following four lemmas. The first three lemmas show that, for all  $j, 1 \leq j \leq \lceil\log\log|\mathbf{x}|\rceil + 1$ , there are points  $P_j \in A_j$  such that  $(P_j, P_{j-1})$  and  $(P_{\lceil\log\log|\mathbf{x}|\rceil+1}, \theta')$  belong to  $G_p(\mathbb{Z}^d)$ . The fourth lemma shows that this happens on an event with probability  $1 - o(1)$ . By consecutively linking  $P_j, 0 \leq j \leq \lceil\log\log|\mathbf{x}|\rceil + 1$ , and  $\theta'$ , we construct a path joining  $P_0$  to  $\theta'$  with  $\lceil\log\log|\mathbf{x}|\rceil + 2$

edges. Since  $\mathbf{0}'$  is within  $\frac{1}{2} \log \log |\mathbf{x}|$  of  $\mathbf{0}$ , we need at most  $\lceil \log \log |\mathbf{x}| \rceil$  edges to join  $\mathbf{0}'$  to  $\mathbf{0}$  (recall (3.2)). This gives a path joining  $\mathbf{P}_0$  to  $\mathbf{0}$  with at most  $2\lceil \log \log |\mathbf{x}| \rceil + 2$  edges. Since  $2 + 2\lceil \log \log |\mathbf{x}| \rceil \leq 4 + 2 \log \log |\mathbf{x}|$ , we obtain Theorem 1.2, as desired. We now turn to our four lemmas.

**Lemma 3.3.** *There exists an event  $E_1$ , with  $P[E_1] = 1 - O(r_1^{-dK})$ , such that on  $E_1$  there is a node  $\mathbf{P}_1 \in A_1$  which is linked to  $\mathbf{P}_0$ , i.e. the edge  $(\mathbf{P}_0, \mathbf{P}_1)$  is in  $G_p(\mathbb{Z}^d)$ .*

*Proof.* The number of lattice points in  $A_1$  is  $\Theta(|\mathbf{x}|^{db^{-1+\varepsilon}})$ , i.e. there is a constant  $C > 0$  such that the number of lattice points is bounded from above by  $C|\mathbf{x}|^{db^{-1+\varepsilon}}$  and bounded from below by  $C^{-1}|\mathbf{x}|^{db^{-1+\varepsilon}}$ . Lemma 3.1 implies that there is an event  $E_1$ , depending only on  $\{U_z\}_{z \in A_1}$  and with

$$P[E_1] = 1 - O(|\mathbf{x}|^{-dKb^{-1+\varepsilon}}),$$

such that, for large  $|\mathbf{x}|$ ,  $E_1$  implies the existence of  $\mathbf{P}_1 \in A_1$  with

$$U_{\mathbf{P}_1} \leq \frac{K \log(|\mathbf{x}|^{db^{-1+\varepsilon}})}{|\mathbf{x}|^{db^{-1+\varepsilon}}}.$$

Again, if there is more than one node in  $A_1$  satisfying this inequality, we choose the one with smallest lexicographical order. Since  $b := pd$  it follows for large  $|\mathbf{x}|$  that  $\mathbf{P}_1$  has weight

$$U_{\mathbf{P}_1}^{-p} \geq \frac{|\mathbf{x}|^{b\varepsilon}}{(K \log(|\mathbf{x}|^{db^{-1+\varepsilon}}))^p} \geq 2|\mathbf{x}|. \tag{3.4}$$

We now show that  $\mathbf{P}_1$  is linked to  $\mathbf{P}_0$ . It suffices to show that

$$|\mathbf{P}_0 - \mathbf{P}_1| \leq \min(U_{\mathbf{P}_0}^{-p}, U_{\mathbf{P}_1}^{-p}).$$

However,  $|\mathbf{P}_0 - \mathbf{P}_1| \leq |\mathbf{P}_0| + |\mathbf{P}_1| \leq 2|\mathbf{x}|$ , so Lemma 3.3 follows from (3.3) and (3.4).

Given  $\mathbf{x}$ , let  $m := m(\mathbf{x})$  denote the largest integer such that  $r_m \geq \log \log |\mathbf{x}|$ ;  $m$  is well defined since  $r_j \downarrow 1$ . If  $t := \lceil 1/(1 - \varepsilon) \rceil \log \log |\mathbf{x}|$  then

$$|\mathbf{x}|^{b^{-t(1-\varepsilon)}} = |\mathbf{x}|^{1/\log |\mathbf{x}|} = b,$$

showing that  $m$  is bounded by  $t$ . The next lemma extends the arguments of Lemma 3.3 and builds a path of  $m$  edges from  $\mathbf{P}_0$  to a node  $\mathbf{P}_m \in A_m$ .

**Lemma 3.4.** *For all  $j$ ,  $1 \leq j \leq m$ , there is an event  $E_j$ , depending only on  $\{U_z\}_{z \in A_j}$ , such that*

- (i)  $P[E_j] = 1 - O(r_j^{-dK})$ , and
- (ii) on each  $E_j$  there is a node  $\mathbf{P}_j \in A_j$  such that, on  $E_{j-1} \cap E_j$ , the edge  $(\mathbf{P}_{j-1}, \mathbf{P}_j)$  is in  $G_p$ .

*Proof.* Since  $\text{card}\{A_j\} = \Theta(r_j^d)$ , Lemma 3.1 implies that for large  $|\mathbf{x}|$  there is an event  $E_j$ , with  $P[E_j] = 1 - O(r_j^{-dK})$ , which depends only on  $\{U_z\}_{z \in A_j}$  and which implies the existence of a  $\mathbf{P}_j \in A_j$  satisfying

$$U_{\mathbf{P}_j} \leq \frac{K \log(r_j^d)}{r_j^d} =: W_j.$$

It remains to show that

$$|\mathbf{P}_j - \mathbf{P}_{j-1}| \leq \min(U_{\mathbf{P}_j}^{-p}, U_{\mathbf{P}_{j-1}}^{-p}) \tag{3.5}$$

for all  $j, 1 \leq j \leq m$ . Lemma 3.3 shows (3.5) for  $j = 1$ . The maximal distance between points in  $A_j$  and  $A_{j-1}$  is at most twice  $r_{j-1}$ , i.e.  $|\mathbf{P}_j - \mathbf{P}_{j-1}| \leq 2r_{j-1}$ . It thus suffices to show that

$$2r_{j-1} \leq \min(W_j^{-p}, W_{j-1}^{-p}) = W_j^{-p}, \tag{3.6}$$

which holds since  $W_{j-1}^{-p} \geq W_j^{-p}$  for all  $j, 1 \leq j \leq m$ .

However, by Lemma 3.2,

$$W_j^{-p} = \frac{r_j^{pd}}{(Kd \log r_j)^p} = \frac{((r_{j-1})^{b^{-1+\epsilon}})^{pd}}{(Kdb^{-1+\epsilon} \log(r_{j-1}))^p}.$$

Thus, for all  $j, 1 \leq j \leq m$ ,

$$\frac{W_j^{-p}}{r_{j-1}} = \frac{(r_{j-1})^{b^\epsilon - 1}}{(Kdb^{-1+\epsilon} \log(r_{j-1}))^p} \geq \frac{(r_m)^{b^\epsilon - 1}}{(Kdb^{-1+\epsilon} \log(r_m))^p},$$

since the  $r_j$  are decreasing. By definition of  $r_m$  and since  $b^\epsilon - 1 > 0$ , the last ratio clearly exceeds 2 for large  $|\mathbf{x}|$ , showing (3.6) and completing the proof of Lemma 3.4.

The next lemma shows that we may link  $\mathbf{P}_m$  and  $\mathbf{0}'$  via a node  $\mathbf{P}_{m+1} \in A_{m+1}$ . Combined with Lemmas 3.2 and 3.3, this builds a path between  $\mathbf{P}_0$  and  $\mathbf{0}'$  which contains  $m + 2$  edges.

**Lemma 3.5.** *There is an event  $E_{m+1}$ , depending only on  $\{U_z\}_{z \in A_{m+1}}$ , such that  $P[E_{m+1}] = 1 - O(r_{m+1}^{-dK})$ , and on  $E_0 \cap E_m \cap E_{m+1}$  there is a point  $\mathbf{P}_{m+1} \in A_{m+1}$  such that the edges  $(\mathbf{P}_m, \mathbf{P}_{m+1})$  and  $(\mathbf{P}_{m+1}, \mathbf{0}')$  both belong to  $G_p(\mathbb{Z}^d)$ .*

*Proof.* First, by definition of  $m$  and by Lemma 3.2 we have

$$(\log \log |\mathbf{x}|)^\beta \leq r_m^\beta = r_{m+1} < \log \log |\mathbf{x}|.$$

By Lemma 3.1, for large  $|\mathbf{x}|$  there is an event  $E_{m+1}$ , with  $P[E_{m+1}] = 1 - O(r_{m+1}^{-dK})$ , which depends only on  $\{U_z\}_{z \in A_{m+1}}$  and which implies the existence of a point  $\mathbf{P}_{m+1} \in A_{m+1}$  with

$$U_{\mathbf{P}_{m+1}} \leq \frac{K \log(r_{m+1}^d)}{r_{m+1}^d} \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^{\beta d}} \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^{(pd)^\epsilon / p}}$$

since  $\beta d = (pd)^\epsilon / p$ . Since  $(pd)^\epsilon > 1$ , it follows that, for large  $|\mathbf{x}|$ , on  $E_{m+1}$  we have

$$U_{\mathbf{P}_{m+1}}^{-p} \geq 2 \log \log |\mathbf{x}|. \tag{3.7}$$

Following the arguments in the proof of Lemma 3.4 (with  $j$  equal to  $m + 1$  there), we find that, on  $E_m \cap E_{m+1}$ ,  $(\mathbf{P}_m, \mathbf{P}_{m+1})$  is an edge in  $G_p(\mathbb{Z}^d)$ . Furthermore, on  $E_0 \cap E_m \cap E_{m+1}$ , the edge  $(\mathbf{P}_{m+1}, \mathbf{0}')$  belongs to  $G_p(\mathbb{Z}^d)$  if and only if

$$|\mathbf{0}' - \mathbf{P}_{m+1}| \leq \min(U_{\mathbf{0}'}^{-p}, U_{\mathbf{P}_{m+1}}^{-p}). \tag{3.8}$$

However,

$$|\mathbf{0}' - \mathbf{P}_{m+1}| \leq |\mathbf{0}' - \mathbf{0}| + |\mathbf{0} - \mathbf{P}_{m+1}| \leq \log \log |\mathbf{x}| + r_{m+1} \leq 2 \log \log |\mathbf{x}|,$$

showing that (3.8) follows, using (3.7) and (3.1).

The last lemma completes the proof of Theorem 1.2.

**Lemma 3.6.** *For all  $x \in \mathbb{Z}^d$ , there is an event  $E(x)$ , with  $P[E(x)] = 1 - o(1)$ , such that on  $E(x)$  there exists a path joining  $P_0$  to  $\mathbf{0}$  with  $4 + 2 \log \log |x|$  edges.*

*Proof.* Let  $E(x) := E_0 \cap E_{x/2} \cap (\bigcap_{j=1}^{m+1} E_j)$ . On  $E(x)$  we have shown that there is a path,  $\pi$ , joining  $P_0$  to  $\mathbf{0}$  via the successive nodes  $P_1, P_2, \dots, P_m, P_{m+1}, \mathbf{0}', \mathbf{0}$ . The number of edges in  $\pi$  is bounded by  $m + 2 + \lceil \log \log |x| \rceil$ , where  $\lceil \log \log |x| \rceil$  denotes an upper bound on the number of edges between  $\mathbf{0}'$  and  $\mathbf{0}$ . Since  $\varepsilon$  is arbitrary in the definition of  $t$ , it follows that  $m \leq \lceil \log \log |x| \rceil$ . Thus,  $\text{card}\{\pi\} \leq 4 + 2 \log \log |x|$ .

Finally, we show that  $P[E(x)] = 1 - o(1)$ . For all  $j, 1 \leq j \leq m + 1$ ,  $E_j$  depends only on  $\{U_z\}_{z \in A_j}$  and, since the  $A_j$  are disjoint, the  $\{E_j\}_{1 \leq j \leq m+1}$  are independent. Clearly, since  $E_0$  depends on  $\{U_z\}_{z \in \mathbb{Z}^{d-1} \times \mathbb{Z}^-}$ , we have independence of  $E_0, E_1, E_2, \dots, E_{m+1}$ . Similarly,  $E_{x/2}, E_0, E_1, E_2, \dots, E_{m+1}$  are independent.

By independence, we have

$$P[E(x)] = P\left[\bigcap_{j=1}^{m+1} E_j\right] P[E_0] P[E_x] P[E_{x/2}] = (1 - o(1))^3 \prod_{j=1}^{m+1} P[E_j].$$

Now,  $m$  is bounded by  $C \log \log |x|$  and the definition of  $r_m$  shows, for large  $K$ , that  $m r_{m+1}^{-dK} \rightarrow 0$  as  $|x| \rightarrow \infty$ . Since  $1 - 2s \leq \exp(-s) \leq 1 - s/2$  for small, positive  $s$ , it follows that

$$\begin{aligned} \prod_{j=1}^{m+1} P[E_j] &= \prod_{j=1}^{m+1} (1 - O(r_j^{-dK})) \\ &\geq \exp\left(-C \sum_{j=1}^{m+1} r_j^{-dK}\right) \\ &\geq 1 - C \sum_{j=1}^{m+1} r_j^{-dK} \\ &\geq 1 - C \sum_{j=1}^{m+1} r_{m+1}^{-dK}. \end{aligned}$$

This yields  $P[E(x)] = 1 - o(1)$ , as desired, completing the proof of Lemma 3.6.

#### 4. Proof of Theorem 1.3

Assume without loss of generality that  $n$  has even parity. Partition  $[0, n]^d \cap \mathbb{Z}^d$  into  $Q_1 := [0, \frac{1}{2}n] \times [0, n]^{d-1} \cap \mathbb{Z}^d$  and  $Q_2 := (\frac{1}{2}n, n] \times [0, n]^{d-1} \cap \mathbb{Z}^d$ . For all  $k = 0, 1, 2, \dots, n/2$ , write  $Q_{1,k} := \{n/2 - k\} \times [0, n]^{d-1} \cap \mathbb{Z}^d$  and note that  $Q_1 = \bigcup_{k=0}^{n/2} Q_{1,k}$ .

The number of nodes in  $Q_1$  whose balls of influence have nonempty intersection with  $Q_2$  is

$$N := \sum_{k=0}^{n/2} \sum_{i \in Q_{1,k}} 1_{\{U_i^{-p} \geq k+1\}}.$$

If we remove these  $N$  nodes from  $Q_1$  then  $G_p(Q_1)$  and  $G_p(Q_2)$  are disconnected, i.e. the graphs have no edges joining them. Moreover, as the number of nodes in  $Q_{1,k}$  equals  $n^{d-1}$ , we obtain

$$E[N] = \sum_{k=0}^n n^{d-1} \mathbb{P}[U_0^{-p} \geq k+1] = n^{d-1} \sum_{k=0}^n (k+1)^{-1/p} \leq Cn^{d-1} [n^{1-1/p} \vee 1],$$

which is exactly the desired upper bound.

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