

# Perfect Graphs of Fixed Density: Counting and Homogeneous Sets

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For  $c \in (0, 1)$  let  $\mathcal{P}_n(c)$  denote the set of  $n$ -vertex perfect graphs with density  $c$  and let  $\mathcal{C}_n(c)$  denote the set of  $n$ -vertex graphs without induced  $C_5$  and with density  $c$ .

We show that

$$\lim_{n \rightarrow \infty} \log_2 |\mathcal{P}_n(c)| / \binom{n}{2} = \lim_{n \rightarrow \infty} \log_2 |\mathcal{C}_n(c)| / \binom{n}{2} = h(c)$$

with  $h(c) = \frac{1}{2}$  if  $\frac{1}{4} \leq c \leq \frac{3}{4}$  and  $h(c) = \frac{1}{2}H(|2c - 1|)$  otherwise, where  $H$  is the binary entropy function.

Further, we use this result to deduce that almost all graphs in  $\mathcal{C}_n(c)$  have homogeneous sets of linear size. This answers a question raised by Loebl and co-workers.

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## 1. Introduction and results

In this paper we investigate classes of graphs that are defined by forbidding certain substructures. Let  $\mathcal{H}$  be such a class. We focus on two related goals: to approximate the cardinality of  $\mathcal{H}$  and to determine the structure of a typical graph in  $\mathcal{H}$ . In particular, we add the additional constraint that all graphs in  $\mathcal{H}$  must have the same density  $c$  and would like to know how the answer to these questions depends on the parameter  $c$ .

The quantity  $|\mathcal{H}_n|$ , where  $\mathcal{H}_n := \{G \in \mathcal{H} : V(G) = [n]\}$ , is also called the *speed* of  $\mathcal{H}$ . Often exact formulas or good estimates for  $|\mathcal{H}_n|$  are out of reach. In these cases, however, one might

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still ask for the asymptotic behaviour of the speed of  $\mathcal{H}$ . One prominent result in this direction was obtained by Erdős, Frankl and Rödl [9] who considered properties  $\mathcal{F}orb(F)$  defined by a single forbidden (weak) subgraph  $F$ . They proved that, for each graph  $F$  with  $\chi(F) \geq 3$ , the class  $\mathcal{F}orb_n(F)$  of  $n$ -vertex graphs that do not contain  $F$  as a subgraph satisfies  $|\mathcal{F}orb_n(F)| = 2^{\text{ex}(F,n)+o(n^2)}$ , where  $\text{ex}(F,n) := (\chi(F) - 2) \binom{n}{2} / (\chi(F) - 1)$ . In other words, if  $\chi(F) \geq 3$  then the speed of  $\mathcal{F}orb(F)$  asymptotically only depends on the chromatic number of  $F$ .

In this paper we are interested in features of the picture at a more fine-grained scale. More precisely, we fix a density  $0 < c < 1$  and are interested in the number  $|\mathcal{H}_n(c)|$  of graphs on  $n$  vertices with property  $\mathcal{H}$  and density  $c$ . Let  $\mathcal{F}orb_n(F, c) = \mathcal{F}orb_n(F) \cap \mathcal{G}_n(c)$ , where  $\mathcal{G}_n(c)$  is the set of all graphs on vertex set  $[n]$  with  $\lfloor c \binom{n}{2} \rfloor$  edges. Here  $\lfloor x \rfloor$  denotes the nearest integer to  $x$ . For the sake of readability, we will always assume in the following that  $c \binom{n}{2}$  is an integer, since rounding issues would not affect our asymptotic considerations.

Straightforward modifications of the proof of the theorem of Erdős, Frankl and Rödl [9] yield the following bounds for  $|\mathcal{F}orb_n(F, c)|$  (we will sketch this argument in Section 2.3). Let  $F$  be a graph with  $\chi(F) = r$ . For all  $c \in (0, \frac{r-2}{r-1})$ , we have

$$\lim_{n \rightarrow \infty} \frac{\log |\mathcal{F}orb_n(F, c)|}{\binom{n}{2}} = \frac{r-2}{r-1} H\left(\frac{r-1}{r-2}c\right), \tag{1.1}$$

where  $H(x)$  is the *binary entropy* function, that is, for  $x \in (0, 1)$  we set  $H(x) := -x \log x - (1 - x) \log(1 - x)$ . Here we denote by  $\log$  the logarithm to base 2. Notice that

$$\lim_{n \rightarrow \infty} \log |\mathcal{F}orb_n(F, c)| / \binom{n}{2} = 0$$

for  $c > \frac{r-2}{r-1}$ , by the theorem of Erdős and Stone [11].

The analogous problem for a graph class  $\mathcal{F}orb^*(F)$ , characterized by a forbidden *induced* subgraph  $F$ , is more challenging and was first considered by Prömel and Steger [20]. They specified a graph parameter, the so-called colouring number  $\chi^*(F)$  of  $F$ , that serves as a suitable replacement of the chromatic number in the theorem of Erdős, Frankl and Rödl. More precisely, they showed that  $|\mathcal{F}orb_n^*(F)| = 2^{\text{ex}^*(F,n)+o(n^2)}$  with  $\text{ex}^*(F,n) := (\chi^*(F) - 2) \binom{n}{2} / (\chi^*(F) - 1)$ , where  $\chi^*(F)$  is defined as follows. A *generalized  $r$ -colouring* of  $F$  with  $r' \in [0, r]$  cliques is a partition of  $V(F)$  into  $r'$  cliques and  $r - r'$  independent sets. The *colouring number*  $\chi^*(F)$  is the largest integer  $r + 1$  such that there is an  $r' \in [r]$  for which  $F$  has no generalized  $r$ -colouring with  $r'$  cliques. For example, we have  $\chi^*(C_5) = 3$  and  $\chi^*(C_7) = 4$ .

This naturally extends to hereditary graph properties, *i.e.*, classes of graphs  $\mathcal{H}$  which are closed under isomorphism and taking induced subgraphs (and may therefore be characterized by possibly infinitely many forbidden induced subgraphs). Let  $\mathcal{F}(r, r')$  denote the family of all graphs that admit a generalized  $r$ -colouring with  $r'$  cliques. Then the colouring number of  $\mathcal{H}$  is

$$\chi^*(\mathcal{P}) := \max\{r + 1 : \mathcal{F}(r, r') \subseteq \mathcal{H} \text{ for some } r' \in [0, r]\},$$

and we set

$$\text{ex}^*(\mathcal{H}, n) := (\chi^*(\mathcal{H}) - 2) \binom{n}{2} / (\chi^*(\mathcal{H}) - 1).$$

Observe that this definition implies  $\chi^*(\text{Forb}^*(F)) = \chi^*(F)$ . And indeed Alekseev [1] and Bollobás and Thomason [6] generalized the result of Prömel and Steger to arbitrary hereditary graph properties  $\mathcal{H}$  and showed that  $|\mathcal{P}_n| = 2^{\text{ex}^*(\mathcal{H},n)+o(n^2)}$ .

More precise estimates for the speed were given for monotone properties  $\mathcal{H}$  (properties that are closed under isomorphisms and taking subgraphs) by Balogh, Bollobás and Simonovits [5], who showed that  $2^{\text{ex}^*(\mathcal{H},n)} \leq |\mathcal{H}_n| \leq 2^{\text{ex}^*(\mathcal{H},n)+cn \log n}$  for some constant  $c$ , and for hereditary properties  $\mathcal{H}$  by Alon, Balogh, Bollobás and Morris [2], who proved  $2^{\text{ex}^*(\mathcal{H},n)} \leq |\mathcal{H}_n| \leq 2^{\text{ex}^*(\mathcal{H},n)+n^{2-\varepsilon}}$  for some  $\varepsilon = \varepsilon(\mathcal{H}) > 0$  and  $n$  sufficiently large. Prömel and Steger [18, 19] gave even more precise results for the speed of  $\text{Forb}_n^*(C_4)$  and  $\text{Forb}_n^*(C_5)$ , which they determined up to a factor of  $2^{O(n)}$ . In fact, they showed in [19] that almost all graphs in  $\text{Forb}_n^*(C_5)$  are generalized split graphs, that is, graphs of a rather simple structure which are defined as follows. We say that a graph  $G = (V, E)$  admits a *generalized clique partition* if there is a partition  $V = V_1 \dot{\cup} \dots \dot{\cup} V_k$  of its vertex set such that  $G[V_i]$  is a clique and for  $i > j > 1$  we have  $e(V_i, V_j) = e(V_j, V_i) = 0$ . A graph  $G$  is a *generalized split graph* if  $G$  or its complement admit a generalized clique partition.

It is illustrative to compare this result to the celebrated strong perfect graph theorem [8]. A graph  $G$  is *perfect* if  $\chi(G')$  equals the clique number  $\omega(G')$  for all induced subgraphs  $G'$  of  $G$ . The strong perfect graph theorem asserts that *all* graphs without induced copies of odd cycles  $C_{2i+1}$ ,  $i > 1$  and without induced copies of their complements  $\overline{C}_{2i+1}$  are perfect. Using this characterization, it is easy to see that generalized split graphs are perfect. Consequently the result of Prömel and Steger implies that *almost all* graphs without induced  $C_5$  are perfect (observe that  $C_5$  is self-complementary).

In this paper, we consider induced  $C_5$ -free graphs of density  $c$  and provide bounds for their number. In the spirit of the result by Prömel and Steger, we also relate this quantity to the number of  $n$ -vertex perfect graphs and generalized split graphs with density  $c$ .

**Definition 1.1.** We define the following graph classes:

$$\begin{aligned} \mathcal{C}(n, c) &:= \text{Forb}_n^*(C_5, c) := \text{Forb}_n^*(C_5) \cap \mathcal{G}_n(c), \\ \mathcal{P}(n, c) &:= \{G \in \mathcal{G}_n(c) : G \text{ is perfect}\}, \\ \mathcal{S}(n, c) &:= \{G \in \mathcal{G}_n(c) : G \text{ is a generalized split graph}\}. \end{aligned}$$

Observe that for all  $n$  and  $c \in [0, 1]$  we have  $\mathcal{S}(n, c) \subseteq \mathcal{P}(n, c) \subseteq \mathcal{C}(n, c)$ . Our first main result now bounds the multiplicative error term between  $|\mathcal{S}(n, c)|$  and  $|\mathcal{C}(n, c)|$ . In order to state this we define the following function. Let

$$h(c) := \begin{cases} H(2c)/2 & \text{if } 0 < c < \frac{1}{4}, \\ 1/2 & \text{if } \frac{1}{4} \leq c \leq \frac{3}{4}, \\ H(2c - 1)/2 & \text{otherwise.} \end{cases} \tag{1.2}$$

Note that the classes of all generalized split graphs, all perfect graphs, and all graphs without induced  $C_5$  are closed under taking complements. Hence, for example,  $|\mathcal{C}(n, c)| = |\mathcal{C}(n, 1 - c)|$  for all  $c \in (0, 1)$  and  $h$  is in fact symmetric in  $(0, 1)$ . Further, note that  $H(|2c - 1|)/2 = h(c)$  for  $c < 1/4$  or  $c > 3/4$ .

**Theorem 1.2.** *For all  $c \in (0, 1)$  we have*

$$\lim_{n \rightarrow \infty} \frac{\log_2 |\mathcal{C}(n, c)|}{\binom{n}{2}} = \lim_{n \rightarrow \infty} \frac{\log_2 |\mathcal{P}(n, c)|}{\binom{n}{2}} = \lim_{n \rightarrow \infty} \frac{\log_2 |\mathcal{S}(n, c)|}{\binom{n}{2}} = h(c).$$

The proof of this theorem uses Szemerédi’s regularity lemma and is given in Section 2.

We remark that Bollobás and Thomason [7] studied related questions of a more general type (see also the references in [7] for earlier results in this direction). They were interested in the probability  $\mathbb{P}_{\mathcal{H}} := \mathbb{P}[\mathcal{G}(n, p) \in \mathcal{H}]$  of an arbitrary hereditary property  $\mathcal{H}$  in the probability space  $\mathcal{G}(n, p)$ , and showed that for any  $\mathcal{H}$  there are very simple properties  $\mathcal{H}^*$  which closely approximate  $\mathcal{H}$  in the probability space  $\mathcal{G}(n, p)$ . In this context, our Theorem 1.2 estimates the probability of  $\mathcal{H} = \text{Forb}_n^*(C_5)$  in the probability space  $\mathcal{G}(n, m)$  with  $m = c \binom{n}{2}$ , and states that  $\mathcal{H} = \text{Forb}_n^*(C_5)$  is approximated by the property  $\mathcal{H}^*$  of being a generalized split graph in  $\mathcal{G}(n, m)$ . The actual value of the probability  $\mathbb{P}_{\mathcal{H}}$  was estimated by Marchant and Thomason in [17] for several properties  $\mathcal{H}$ , such as  $\mathcal{H} = \text{Forb}_n^*(C_5)$  (see [17, 22]). The probabilities  $\mathbb{P}[\mathcal{G}(n, p) \in \mathcal{H}]$  and  $\mathbb{P}[\mathcal{G}(n, m = p \binom{n}{2}) \in \mathcal{H}]$  are related (but not identical), and we discuss their relation in Section 4.

Let us now move from the question of approximating cardinalities to determining the structure of a typical element in  $\text{Forb}_n^*(C_5)$ . A well-known conjecture by Erdős and Hajnal [10] states that any family of graphs that does not contain a certain fixed graph  $F$  as an induced subgraph must contain a homogeneous set, *i.e.*, a clique or a stable set, which is of size at least some positive power of the number of vertices.

The conjecture is known to be true for certain graphs  $F$ , but is open, among others, for  $F = C_5$  (see [12]). However, Loeb, Reed, Scott, Thomason and Thomassé [16] recently showed that for any graph  $F$ , *almost all graphs* in  $\text{Forb}_n^*(F)$  have a homogeneous set of size at least some positive power of  $n$ . Moreover, they ask for which graphs  $F$  it is true that almost all graphs in  $\text{Forb}_n^*(F)$  do indeed have a *linearly* sized homogeneous set.

It may seem at first sight that our estimates derived in Theorem 1.2, carrying an  $o(n^2)$  term in the exponent, are too rough to tell us much about the structure of almost all graphs in  $\text{Forb}_n^*(C_5)$  or  $\text{Forb}_n^*(C_5, c)$ . However, we can combine them with the ideas of [16] to answer the question of Loeb, Reed, Scott, Thomason and Thomassé in the affirmative for the case  $F = C_5$ . In fact, we can prove this assertion even in the case where we again restrict the class to graphs with a given density.

**Theorem 1.3.** *For  $\eta > 0$ , let  $\text{Forb}_{n,\eta}^*(F, c)$  denote the set of graphs  $G \in \text{Forb}_n^*(F, c)$  with  $\text{hom}(G) := \max\{\alpha(G), \omega(G)\} < \eta n$ . Then, for every  $0 < c < 1$  there exists  $\eta > 0$  such that*

$$\frac{|\text{Forb}_{n,\eta}^*(C_5, c)|}{|\text{Forb}_n^*(C_5, c)|} \rightarrow 0 \quad (n \rightarrow \infty).$$

We provide the proof of this theorem in Section 3.

Statements similar to those in Theorem 1.2 and 1.3, for forbidden graphs  $F$  other than  $C_5$ , seem to require more work.

## 2. The proof of Theorem 1.2

In this section we prove Theorem 1.2. In Section 2.1 we start with the lower bound, by estimating the number of generalized split graphs with a given density. For the upper bound we need some preparations. We shall apply Szemerédi’s regularity lemma, which is introduced in Section 2.2. In Section 2.3 we illustrate how this lemma can be used for counting graphs without a fixed (not necessarily induced) subgraph. In Section 2.4 we explain how to modify these ideas in order to deal with forbidden induced subgraphs. In Section 2.5, finally, we prove the upper bound of Theorem 1.2.

### 2.1. The lower bound of Theorem 1.2

In this section we estimate the number of generalized split graphs with density  $c$  and prove the following lemma, which constitutes the lower bound of Theorem 1.2.

**Lemma 2.1.** *For all  $c, \gamma \in (0, 1)$  there is an  $n_0$  such that, for all  $n \geq n_0$ , we have*

$$|\mathcal{S}(n, c)| \geq 2^{h(c)\binom{n}{2} - \gamma \binom{n}{2}}.$$

We will use the following bound for binomial coefficients (see, e.g., [14]). For every  $\gamma > 0$  there exists  $n_0$  such that, for every integer  $m \geq n_0$  and for every real  $c \in (0, 1)$ , we have

$$2^{mH(c) - \gamma m} \leq \binom{m}{cm} \leq 2^{mH(c)}. \tag{2.1}$$

We call the term  $-\gamma m$  in the first exponent the *error term of equation (2.1)*.

**Proof of Lemma 2.1.** We prove this lower bound by constructing an adequate number of generalized split graphs. Choose  $n_0$  sufficiently large such that (2.1) holds for  $m = \frac{1}{2}\binom{n}{2}$  and error term  $\gamma m$ . Observe that it suffices to prove the lemma for  $c \leq \frac{1}{2}$ , since the complement of a split graph with density  $c$  is a split graph with density  $(1 - c)$ .

We distinguish two cases. First, assume  $c \leq \frac{1}{4}$ . To obtain a lower bound for  $|\mathcal{S}(n, c)|$  in this case, we simply count bipartite graphs with density  $c$  and with colour classes of size  $n/2$ . There are at least

$$\binom{\frac{n^2}{4}}{c\binom{n}{2}} \geq \binom{\frac{1}{2}\binom{n}{2}}{c\binom{n}{2}} \geq 2^{\frac{1}{2}\binom{n}{2}H(2c) - \gamma \binom{n}{2}}$$

such graphs.

Now assume that  $\frac{1}{4} < c \leq \frac{1}{2}$ . In this case we construct suitable  $k$ -partite graphs. For this purpose choose  $k$  such that

$$x := c\binom{n}{2} - (k - 2)\left(\frac{n}{2} - k + 2\right) - \binom{k - 2}{2} \in \left[\frac{n^2}{8} - n, \frac{n^2}{8} + n\right]. \tag{2.2}$$

Indeed such a  $k$  exists since  $x$  monotonically decreases from  $c\binom{n}{2}$  for  $k = 2$  to  $c\binom{n}{2} - (n^2 - 2n - 8)/8 \leq (n^2 + 8)/8$  for  $k = n/2$  in steps of size at most  $n/2$ .

Now, construct  $k$  (independent) vertex sets  $V_1, \dots, V_k$  with  $|V_1| = \frac{n}{2}$ ,  $|V_2| = \frac{n}{2} - k + 2$  and  $|V_i| = 1$  for  $i = \{3, \dots, k\}$  and insert all edges between  $V_i$  and  $V_j$  with  $i, j \in [k] \setminus \{1\}$ ,  $i \neq j$ . Call

the resulting graph  $G_0$ . By (2.2) we obtain a generalized split graph with density  $c$  from  $G_0$ , if we insert  $x$  edges between  $V_1$  and  $V_2 \cup \dots \cup V_k$ . Since this can be done in at least

$$\binom{\frac{n^2}{4}}{\frac{n^2}{8} - n} \geq 2^{\frac{1}{2} \binom{n}{2} - \gamma \binom{n}{2}}$$

ways, we obtain at least  $2^{\frac{1}{2} \binom{n}{2} - \gamma \binom{n}{2}}$  generalized split graphs with exactly  $c \binom{n}{2}$  edges. □

**2.2. Regularity**

In order to prove the upper bound from Theorem 1.2, *i.e.*,

$$|\mathcal{C}(n, c)| \leq 2^{h(c) \binom{n}{2} + \gamma \binom{n}{2}},$$

we will analyse the structure of graphs in  $\mathcal{C}(n, c)$  by applying a variant of the regularity lemma suitable for our purposes.

Let  $G = (V, E)$  be a graph. For disjoint non-empty vertex sets  $A, B \subseteq V$ , the *density*  $d(A, B) := e(A, B)/(|A||B|)$  of the pair  $(A, B)$  is the number of edges that run between  $A$  and  $B$  divided by the number of possible edges between  $A$  and  $B$ . In the following let  $\varepsilon, d \in [0, 1]$ . The pair  $(A, B)$  is  $\varepsilon$ -regular if, for all  $A' \subseteq A$  and  $B' \subseteq B$  with  $|A'| \geq \varepsilon|A|$  and  $|B'| \geq \varepsilon|B|$ , it is true that  $|d(A, B) - d(A', B')| \leq \varepsilon$ . An  $\varepsilon$ -regular pair  $(A, B)$  is called  $(\varepsilon, d)$ -regular if it has density at least  $d$ .

A partition  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  of  $V$  is an *equipartition* if  $|V_i| = |V_j|$  for all  $i, j \in [k]$ . An  $(\varepsilon, d)$ -regular partition of  $G$  with *reduced graph*  $R = (V_R, E_R)$  is an equipartition  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  of  $V$  with  $|V_0| \leq \varepsilon|V|$ , and  $V_R = [k]$  such that  $(V_i, V_j)$  is an  $(\varepsilon, d)$ -regular pair in  $G$  if and only if  $\{i, j\} \in E_R$ . In this case we also call  $R$  an  $(\varepsilon, d)$ -reduced graph of  $G$ . An  $(\varepsilon, 0)$ -regular partition  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  which has at most  $\varepsilon \binom{k}{2}$  pairs that are not  $\varepsilon$ -regular is also called an  $\varepsilon$ -regular partition. The partition classes  $V_i$  with  $i \in [k]$  are called *clusters* of  $G$  and  $V_0$  is the *exceptional set*.

With this terminology at hand we can state the celebrated regularity lemma of Szemerédi.

**Lemma 2.2 (regularity lemma [21]).** *For all  $\varepsilon > 0$  and  $k_0$  there is a  $k_1$  such that every graph  $G = (V, E)$  on  $n \geq k_1$  vertices has an  $\varepsilon$ -regular partition  $V = V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  with  $k_0 \leq k \leq k_1$ .* □

The strength of this lemma becomes apparent when it is complemented with corresponding embedding lemmas, such as the following (see, *e.g.*, [15]). A *homomorphism* from a graph  $H = (V_H, E_H)$  to a graph  $R = (V_R, E_R)$  is an edge-preserving mapping from  $V_H$  to  $V_R$ .

**Lemma 2.3 (embedding lemma).** *For every  $d > 0$  and every integer  $k$  there exists  $\varepsilon > 0$  with the following property. Let  $H$  be a graph on  $k$  vertices  $v_1, \dots, v_k$ . Let  $G$  be a graph. Let  $V_1, \dots, V_k$  be clusters of an  $(\varepsilon, d)$ -regular partition of  $G$  with reduced graph  $R = ([k], E_R)$ . If there is a homomorphism from  $H$  to  $R$ , then  $H$  is a subgraph of  $G$ .* □

**2.3. Regular partitions and counting**

As a warm-up (and for the sake of completeness) we consider the problem of counting graphs of a fixed density without a given (not necessarily induced) subgraph  $F$ , and prove (1.1). For

this purpose we mimic the proof given by Erdős, Frankl and Rödl in [9] for the corresponding problem without fixed density.

**Proof of (1.1).** Let  $F$  be a graph with  $\chi(F) = r$  and  $c \in (0, \frac{r-2}{r-1})$ . Let  $\gamma > 0$  be given. For large enough  $n$  the lower bound

$$|\mathcal{F}orb_n(F, c)| \geq 2^{\frac{r-2}{r-1}H(\frac{r-1}{r-2}c)\binom{n}{2}-\gamma\binom{n}{2}}$$

can easily be obtained by counting subgraphs with  $c\binom{n}{2}$  edges of the complete  $(r - 1)$ -partite graph with  $n/(r - 1)$  vertices in each part, and by applying (2.1).

It remains to show the upper bound

$$|\mathcal{F}orb_n(F, c)| \leq 2^{\frac{r-2}{r-1}H(\frac{r-1}{r-2}c)\binom{n}{2}+\gamma\binom{n}{2}}$$

for  $n$  sufficiently large. We choose  $d$  such that

$$0 < d \leq \min \left\{ \frac{1}{16}, \frac{1}{4}\gamma^2, \left( -2\frac{r-1}{r-2} \log \left( 1 - c\frac{r-1}{r-2} \right) \right)^{-2} \right\}.$$

Since the entropy function is concave, we have for each  $\hat{d} \in [0, d]$  that

$$\begin{aligned} H\left((c - 2\hat{d})\frac{r-1}{r-2}\right) &\leq H\left(c\frac{r-1}{r-2}\right) - 2\hat{d}\frac{r-1}{r-2}H'\left(c\frac{r-1}{r-2}\right) \\ &\leq H\left(c\frac{r-1}{r-2}\right) - 2\hat{d}\frac{r-1}{r-2} \log\left(1 - c\frac{r-1}{r-2}\right) \leq H\left(c\frac{r-1}{r-2}\right) + \sqrt{\hat{d}}, \end{aligned} \tag{2.3}$$

which we shall use later. Next, let  $\varepsilon$  be the constant returned from Lemma 2.3 for input  $d$  and with  $k$  replaced by  $r$ . Set  $k_0 = \lceil 10/d \rceil$  and let  $k_1$  be the constant returned by Lemma 2.2 for input  $k_0$  and  $\varepsilon$ . Further let  $n \geq k_1$ .

Now we use the regularity lemma, Lemma 2.2, with parameters  $\varepsilon, k_0$  for each graph  $G$  in  $\mathcal{F}orb_n(F, c)$ . For each such application the regularity lemma produces an  $\varepsilon$ -regular partition with at most  $k_1$  clusters, for which we can construct the corresponding  $(\varepsilon, d)$ -reduced graph  $R$ . Since  $k_1$  is finite there is only a finite number of different reduced graphs  $R$  resulting from these applications of the regularity lemma. Hence we can partition  $\mathcal{F}orb_n(F, c)$  into a finite number of classes  $\mathcal{R}(R, \varepsilon, d, n, F, c)$  of graphs with  $(\varepsilon, d)$ -reduced graph  $R$ . Accordingly, it suffices to show that for each  $R$  we have

$$|\mathcal{R}(R, \varepsilon, d, n, F, c)| \leq 2^{\frac{r-2}{r-1}H(\frac{r-1}{r-2}c)\binom{n}{2}+\gamma\binom{n}{2}}. \tag{2.4}$$

Let  $R = (V_R, E_R)$  be any graph such that  $\mathcal{R}(R, \varepsilon, d, n, F, c)$  is non-empty, let  $k = |V_R|$ , let  $G \in \mathcal{R}(R, \varepsilon, d, n, F, c)$  and let  $P$  be an  $\varepsilon$ -regular partition of  $G$  corresponding to  $R$ . By the choice of  $k_0$  at most  $k\binom{n/k}{2} \leq \frac{d}{2}\binom{n}{2}$  edges of  $G$  are inside clusters of  $P$ , at most  $d\binom{n}{2}$  edges of  $G$  are in regular pairs of  $P$  with density less than  $d$ , and at most  $2\varepsilon n^2 \leq \frac{d}{2}\binom{n}{2}$  edges of  $G$  are in irregular pairs of  $P$  or have a vertex in the exceptional set. We conclude that at least  $(c - 2d)\binom{n}{2}$  edges of  $G$  lie in  $(\varepsilon, d)$ -regular pairs of  $P$ . In addition, by the choice of  $\varepsilon$  and since  $F$  has chromatic number  $r$ , Lemma 2.3 implies that  $K_r \not\subseteq R$ . It follows from Turán’s theorem that  $|E_R| \leq \frac{r-2}{r-1}\binom{k}{2}$ . Summarizing, we can bound the number of graphs in  $\mathcal{R}(R, \varepsilon, d, n, F, c)$  by bounding the number of ways to distribute at least  $(c - 2d)\binom{n}{2}$  edges to at most  $\frac{r-2}{r-1}\binom{k}{2}$  regular pairs (corresponding to edges of  $R$ ) with clusters of size at most  $n/k$ , and distributing at most  $2d\binom{n}{2}$  edges arbitrarily. By

the choice of  $n$ , the first of these two factors can be bounded by

$$\begin{aligned} \max_{0 \leq \hat{d} \leq d} \binom{\frac{r-2}{r-1} \binom{n}{2}}{(c-2\hat{d}) \binom{n}{2}} &\stackrel{(2.1)}{\leq} \max_{0 \leq \hat{d} \leq d} 2^{\frac{r-2}{r-1} \binom{n}{2} H(\frac{r-1}{r-2}(c-2\hat{d}))} \\ &\stackrel{(2.3)}{\leq} \max_{0 \leq \hat{d} \leq d} 2^{\frac{r-2}{r-1} \binom{n}{2} H(c \frac{r-1}{r-2}) + \sqrt{\hat{d}} \binom{n}{2}} \\ &\leq 2^{\frac{r-2}{r-1} \binom{n}{2} H(c \frac{r-1}{r-2}) + \sqrt{d} \binom{n}{2}}, \end{aligned}$$

and the second by  $\leq 2^{2d \binom{n}{2}}$ . Since  $2d + \sqrt{d} \leq \gamma$  this implies (2.4) as desired. □

**2.4. Embedding induced subgraphs**

In the last section we showed how the regularity lemma and a corresponding embedding lemma can be used to count graphs with forbidden subgraphs. In this section we provide the tools that will allow us to adapt this strategy to the setting of forbidden induced subgraphs.

We remark that the concepts and ideas presented in this section are not new. They have been used for various similar applications, *e.g.*, by Bollobás and Thomason [7] and Loeb, Reed, Scott, Thomason and Thomassé [16], as well as for different applications such as property testing, *e.g.*, by Alon, Fischer, Krivelevich and Szegedy [3] and Alon and Shapira [4].

We start with an embedding lemma for induced subgraphs, which allows us to find an induced copy of a graph  $F$  in a graph  $G$  with reduced graph  $R$  if  $F$  is an induced subgraph of  $R$  (see, *e.g.*, [3]).

**Lemma 2.4 (injective embedding lemma for induced subgraphs).** *For every  $d > 0$  and every integer  $k$  there exists  $\varepsilon > 0$  such that for all  $f \leq k$  the following holds. Let  $V_1, \dots, V_f$  be clusters of an  $\varepsilon$ -regular partition of a graph  $G$  such that for all  $1 \leq i < j \leq f$  the pair  $(V_i, V_j)$  is  $\varepsilon$ -regular. Let  $F = (V_F, E_F)$  be a graph on  $f$  vertices and let  $g : V_F \rightarrow [f]$  be an injective mapping from  $F$  to the clusters of  $G$  such that for all  $1 \leq i < j \leq f$  we have  $d(V_i, V_j) \geq d$  if  $\{g^{-1}(i), g^{-1}(j)\} \in E(F)$  and  $d(V_i, V_j) \leq 1 - d$  if  $\{g^{-1}(i), g^{-1}(j)\} \notin E(F)$ . Then  $G$  contains an induced copy of  $F$ . □*

In contrast to Lemma 2.3, this lemma allows us to embed only one vertex per cluster of  $G$ . Our goal in the following will be to describe an embedding lemma for induced subgraphs which allows us to embed more than one vertex per cluster. Observe first that for this purpose we must have some control over the existence of edges respectively non-edges *inside* clusters of a regular partition of  $G$ . This can be achieved by applying the following lemma to each of these clusters. It is not difficult to infer this lemma from the regularity lemma (Lemma 2.2) by applying Turán’s theorem and Ramsey’s theorem (see, *e.g.*, [3]).

We use the following definition. A  $(\mu, \varepsilon, k)$ -subpartition of a graph  $G = (V, E)$  is a family of pairwise disjoint vertex sets  $W_1, \dots, W_k \subseteq V$  with  $|W_i| \geq \mu|V|$  for all  $i \in [k]$  such that each pair  $(W_i, W_j)$  with  $\{i, j\} \in \binom{[k]}{2}$  is  $\varepsilon$ -regular. A  $(\mu, \varepsilon, k)$ -subpartition  $W_1, \dots, W_k$  of  $G$  is *dense* if  $d(W_i, W_j) \geq \frac{1}{2}$  for all  $\{i, j\} \in \binom{[k]}{2}$ , and *sparse* if  $d(W_i, W_j) < \frac{1}{2}$  for all  $\{i, j\} \in \binom{[k]}{2}$ .

**Lemma 2.5.** *For every  $k$  and  $\varepsilon$  there exists  $\mu > 0$  such that every graph  $G = (V, E)$  with  $n \geq \mu^{-1}$  vertices either has a sparse or a dense  $(\mu, \varepsilon, k)$ -subpartition. □*



The idea for the embedding lemma for induced subgraphs  $F$  of  $G$  is now as follows. We first find a regular partition of  $G$ . By Lemma 2.4, if a regular pair  $(V_i, V_j)$  in this partition is very dense then we can embed edges of  $F$  into  $(V_i, V_j)$ , if it is very sparse then we can embed non-edges of  $F$ , and if its density is neither very small nor very big then we can embed both edges and non-edges of  $F$ . Moreover, Lemma 2.5 asserts that each cluster either has a sparse or a dense subpartition. In the first case we can embed non-edges inside this cluster, in the second case we can embed edges.

This motivates us to tag the reduced graphs with some additional information. For this purpose we colour an edge of the reduced graph white if the corresponding regular pair is sparse, grey if it is of medium density, and black if it is dense. Moreover, we colour a cluster white if it has a sparse subpartition and black otherwise. We call a cluster graph that is coloured in this way a *type*. The following definitions make this precise.

**Definition 2.6 (coloured graph, type).** A *coloured graph*  $R$  is a triple  $(V_R, E_R, \sigma)$  such that  $(V_R, E_R)$  is a graph and  $\sigma : V_R \cup E_R \rightarrow \{0, \frac{1}{2}, 1\}$  is a colouring of the vertices and the edges of this graph where  $\sigma(V_R) \subseteq \{0, 1\}$ . Vertices and edges with colour 0,  $\frac{1}{2}$ , and 1 are also called *white*, *grey*, and *black*, respectively.

Let  $G = (V, E)$  be a graph and let  $V = V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  be an  $\varepsilon$ -regular partition of  $G$  with reduced graph  $([k], E_R)$ . The  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  corresponding to the partition  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  is the coloured graph  $R = ([k], E_R, \sigma)$  with colouring

$$\sigma(\{i, j\}) = \begin{cases} 0 & \text{if } d(V_i, V_j) < d, \\ 1 & \text{if } d(V_i, V_j) > 1 - d, \\ \frac{1}{2} & \text{otherwise,} \end{cases}$$

for all  $\{i, j\} \in E_R$ , and

$$\sigma(i) = \begin{cases} 0 & \text{if } G[V_i] \text{ has a sparse } (\mu, \varepsilon', k')\text{-subpartition,} \\ 1 & \text{if } G[V_i] \text{ has a dense } (\mu, \varepsilon', k')\text{-subpartition,} \end{cases}$$

for all  $i \in [k]$ , where  $\mu$  is the constant from Lemma 2.5 for input  $k'$  and  $\varepsilon'$ . In this case we also simply say that  $G$  has  $(\varepsilon, \varepsilon', d, k')$ -type  $R$ .

By the discussion above a combination of the regularity lemma, Lemma 2.2, and Lemma 2.5 gives the following.

**Lemma 2.7 (type lemma).** For every  $\varepsilon, \varepsilon' \in (0, \frac{1}{2})$  and for all integers  $k', k_0$ , there are integers  $k_1$  and  $n_0$  such that, for every  $d > 0$ , every graph  $G$  on at least  $n_0$  vertices has an  $(\varepsilon, \varepsilon', d, k')$ -type  $R = ([k], E_R, \sigma)$  with  $k_0 \leq k \leq k_1$  and with at most  $\varepsilon k^2$  non-edges.

**Proof.** Given  $\varepsilon, \varepsilon'$  and  $k', k_0$ , we let  $k_1$  be the constant returned from Lemma 2.2 for input  $\varepsilon$  and  $k_0$ , and let  $\mu$  be the constant returned from Lemma 2.5 for input  $\varepsilon'$  and  $k'$ . Set  $n_0 := 2\mu^{-1}k_1$ .

Now let  $d$  be given and let  $G$  be a graph on at least  $n_0$  vertices. By Lemma 2.2 the graph  $G$  has an  $\varepsilon$ -regular partition  $V = V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  such that  $k_0 \leq k \leq k_1$ . By definition at most  $\varepsilon \binom{k}{2} \leq \varepsilon k^2$  pairs  $(V_i, V_j)$  are not  $\varepsilon$ -regular. Let  $R' = ([k], E_R)$  be the  $\varepsilon$ -reduced graph of  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$ .

It follows that  $R'$  has at most  $\varepsilon k^2$  non-edges. Let  $i \in [k]$ . Since  $|V_i| \geq (1 - \varepsilon)n/k \geq n/(2k_1) \geq \mu^{-1}$ , we can apply Lemma 2.5 and conclude that  $V_i$  either has a sparse or a dense  $(\mu, \varepsilon', k')$ -subpartition. Accordingly we obtain an  $(\varepsilon, \varepsilon', d, k')$ -type  $R = ([k], E_R, \sigma)$  for  $G$  by colouring the edges and vertices of  $R'$  as specified in Definition 2.6. □

For formulating our embedding lemma we need one last preparation. We generalize the concept of a graph homomorphism to the setting of coloured graphs.

**Definition 2.8 (coloured homomorphism).** Let  $F = (V_F, E_F)$  be a graph and let  $R = (V_R, E_R, \sigma)$  be a coloured graph. A *coloured homomorphism* from  $F$  to  $R$  is a mapping  $h: V_F \rightarrow V_R$  with the following properties.

- (a) If  $u, v \in V_F$  and  $h(u) \neq h(v)$  then  $\{h(u), h(v)\} \in E_R$ .
- (b) If  $\{u, v\} \in E_F$  then  $h(u) = h(v)$  and  $\sigma(h(u)) = 1$ , or  $h(u) \neq h(v)$  and  $\sigma(\{h(u), h(v)\}) \in \{\frac{1}{2}, 1\}$ .
- (c) If  $\{u, v\} \notin E_F$  then  $h(u) = h(v)$  and  $\sigma(h(u)) = 0$ , or  $h(u) \neq h(v)$  and  $\sigma(\{h(u), h(v)\}) \in \{0, \frac{1}{2}\}$ .

If there is a coloured homomorphism from  $F$  to  $R$  we also write  $F \xrightarrow{\sigma} R$ .

The following embedding lemma states that a graph  $F$  is an induced subgraph of a graph  $G$  with type  $R$  if there is a coloured homomorphism from  $F$  to  $R$ . This lemma is inherent in [4], for example. For completeness we provide its proof below.

**Lemma 2.9 (embedding lemma for induced graphs).** *For every pair of integers  $k, k'$  and for every  $d \in (0, 1)$  there are  $\varepsilon, \varepsilon' > 0$  such that the following holds. Let  $f \leq k'$  and  $G$  be a graph on  $n$  vertices with  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  on  $k$  vertices. Let  $F$  be an  $f$ -vertex graph such that there is a coloured homomorphism from  $F$  to  $R$ . Then  $F$  is an induced subgraph of  $G$ .*

**Proof.** Let  $k, k' \in \mathbb{N}$  and  $d \in (0, 1)$  be given. Let  $\varepsilon'$  be given by Lemma 2.4 for input  $\frac{d}{2}$  and  $k'$ . Let  $\mu$  be the constant from Lemma 2.5 for input  $k'$  and  $\varepsilon'$ . Set  $\varepsilon := \min\{d/2, \mu\varepsilon'\}$ .

Let  $G, R = ([k], E_R, \sigma)$  and  $F$  be as required, let  $V_0 \dot{\cup} V_1 \dot{\cup} \dots \dot{\cup} V_k$  be an  $\varepsilon$ -regular partition of  $G$  corresponding to  $R$  and let  $h: F \xrightarrow{\sigma} R$  be a coloured homomorphism from  $F$  to  $R$ . For each  $i \in [k]$  we have by definition that if  $\sigma(V_i) = 0$  then  $V_i$  has a sparse  $(\mu, \varepsilon', k')$ -subpartition  $W_{i,1}, \dots, W_{i,k'}$ , and if  $\sigma(V_i) = 1$  then  $V_i$  has a dense  $(\mu, \varepsilon', k')$ -subpartition  $W_{i,1}, \dots, W_{i,k'}$ .

Observe that  $\bigcup_{i \in [k], j \in [k']} W_{i,j}$  has the following properties, since  $\varepsilon' \geq \varepsilon/\mu$  and  $\varepsilon \leq d/2$ . If  $\{i, i'\} \in E_R$ , then for all  $j, j' \in [k']$  the pair

$$(W_{i,j}, W_{i',j'}) \text{ is } \varepsilon'\text{-regular,} \tag{2.5}$$

and has density

$$d(W_{i,j}, W_{i',j'}) \in \begin{cases} [0, 2d) & \text{if } \sigma(\{i, i'\}) = 0, \\ (1 - 2d, 1] & \text{if } \sigma(\{i, i'\}) = 1, \\ (\frac{d}{2}, 1 - \frac{d}{2}) & \text{if } \sigma(\{i, i'\}) = \frac{1}{2}. \end{cases} \tag{2.6}$$

Moreover, for all  $i \in [k]$  and all  $j, j' \in [k']$  with  $j \neq j'$  the pair

$$(W_{i,j}, W_{i,j'}) \text{ is } \varepsilon'\text{-regular,} \tag{2.7}$$

and has density

$$d(W_{i,j}, W_{i,j'}) \in \begin{cases} [0, \frac{1}{2}] & \text{if } \sigma(i) = 0, \\ [\frac{1}{2}, 1] & \text{if } \sigma(i) = 1. \end{cases} \tag{2.8}$$

Now, we define an injective mapping  $g : V(F) \rightarrow [k] \times [k']$  as follows. For  $i \in [k]$  let  $F_i := \{x \in V(F) : h(x) = i\}$  and name the vertices in  $F_i$  arbitrarily by  $\{x_{i,1}, \dots, x_{i,f_i}\} = F_i$ . Set  $g(x_{i,j}) := (i, j)$  for all  $j \in [f_i]$ . This is well-defined since  $|F_i| = f_i \leq f \leq k'$ . Let  $I \subseteq [k]$  with  $|I| \leq f$  be the set of indices  $i \in [k]$  such that  $F_i \neq \emptyset$ .

We claim that  $G, \bigcup_{i \in I, j \in [f_i]} W_{i,j}, F$ , and  $g$  satisfy the conditions of Lemma 2.4 with parameters  $d/2, k', \varepsilon',$  and  $f$ . Indeed, by (2.7) each cluster pair  $(W_{i,j}, W_{i,j'})$  with  $i \in I$  and  $j, j' \in [k'], j \neq j'$  is  $\varepsilon'$ -regular. Moreover, for each  $i, i' \in I$  with  $i \neq i'$  we have that there are  $x \in F_i$  and  $y \in F_{i'}$ . By the definition of a coloured homomorphism (Definition 2.8) we have that  $\{V_i, V_{i'}\} = \{h(x), h(y)\} \in E_R$ . Hence (2.5) implies that  $(W_{i,j}, W_{i',j'})$  is  $\varepsilon'$ -regular too. It remains to show that if  $x, y$  are two vertices in  $V(F)$  and  $(i, j) = g(x)$  and  $(i', j') = g(y)$ , then  $d(W_{i,j}, W_{i',j'}) \geq d/2$  if  $\{x, y\} \in E(F)$  and  $d(W_{i,j}, W_{i',j'}) \leq 1 - d/2$  otherwise. To see this, assume first that  $\{x, y\} \in E(F)$ . Then, by the definition of a coloured homomorphism, either  $h(x) = h(y)$  and  $\sigma(h(x)) = 1$ , which implies  $d(W_{i,j}, W_{i',j'}) \geq \frac{1}{2} \geq d/2$  by (2.8). Or  $h(x) \neq h(y)$  and  $\sigma(\{h(x), h(y)\}) \geq \frac{1}{2}$ , and hence we have  $d(W_{i,j}, W_{i',j'}) \geq d/2$  by (2.6). If  $\{x, y\} \notin E(F)$ , on the other hand, then either  $h(x) = h(y)$  and  $\sigma(h(x)) = 0$ , and so  $d(W_{i,j}, W_{i',j'}) \leq \frac{1}{2} \leq 1 - d/2$  by (2.8). Or  $h(x) \neq h(y)$  and  $\sigma(\{h(x), h(y)\}) \leq \frac{1}{2}$  and thus  $d(W_{i,j}, W_{i',j'}) \leq 1 - d/2$  by (2.6).

It follows that we can indeed apply Lemma 2.4 and conclude that  $F$  is an induced subgraph of  $G$  as desired. □

**2.5. The upper bound of Theorem 1.2**

Now we are ready to prove the upper bound of Theorem 1.2, that is, we establish the following lemma.

**Lemma 2.10.** *For all  $c, \gamma \in (0, 1)$ , there is an  $n_0$  such that for all  $n \geq n_0$  we have*

$$|\mathcal{C}(n, c)| \leq 2^{h(c)\binom{n}{2} + \gamma \binom{n}{2}}.$$

The idea of the proof of Lemma 2.10 is as follows. We proceed in three steps. Firstly, as in the proof of (1.1) in Section 2.3, we start by applying the regularity lemma to all graphs in  $\mathcal{C}(n, c)$ . For each of the regular partitions obtained in this way there is a corresponding type, and in total we only get a constant number  $K$  of different types. Secondly, we continue with a structural analysis of the possible types  $R$  for graphs from  $\mathcal{C}(n, c)$  and infer from Lemma 2.9 that  $R$  cannot contain a triangle all of whose edges are grey (see Lemma 2.11). Thirdly, we prove that a coloured graph without such a grey triangle can only serve as a type for at most  $UB(n)$  graphs on  $n$  vertices (see Lemma 2.12). Multiplying  $UB(n)$  with  $K$  then gives the desired bound.

We start with the second step.

**Lemma 2.11.** *For every integer  $k' \geq 5$  and every  $d > 0$ , there exist  $\varepsilon_{L2.11}, \varepsilon'_{L2.11} > 0$  such that for every  $0 < \varepsilon \leq \varepsilon_{L2.11}$  and every  $0 < \varepsilon' \leq \varepsilon'_{L2.11}$  the following is true. If  $G$  is a graph whose  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  contains three grey edges forming a triangle, then  $G$  has an induced  $C_5$ .*

**Proof.** Given  $d > 0$  and  $k' \geq 5$ , set  $k := 3$  and let  $\varepsilon_{L2.11}, \varepsilon'_{L2.11} > 0$  be the constants given by Lemma 2.9 for  $d, k$ , and  $k'$ . Let positive constants  $\varepsilon \leq \varepsilon_{L2.11}$  and  $\varepsilon' \leq \varepsilon'_{L2.11}$  be given. Let  $G$  be a graph with  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  such that  $R$  contains a triangle  $T$  with three grey edges.

By Lemma 2.9 the graph  $G$  contains an induced  $C_5$  if there exists a coloured homomorphism from  $C_5$  to  $T$ , in which case we are done. We claim that such a coloured homomorphism  $h$  always exists (regardless of the colours of the vertices of  $T$ ). Indeed, if  $T$  has at least two black vertices  $V_1, V_2$  then we can construct  $h$  by mapping a pair of adjacent vertices of  $C_5$  to  $V_1$ , a disjoint pair of adjacent vertices of  $C_5$  to  $V_2$ , and the remaining vertex of  $C_5$  to the remaining vertex of  $T$ . If  $T$  has at least two white vertices  $V_1, V_2$ , on the other hand, then we can construct  $h$  by mapping one pair of non-adjacent vertices of  $C_5$  to  $V_1$ , a disjoint pair of non-adjacent vertices of  $C_5$  to  $V_2$  and the remaining vertex of  $C_5$  to the remaining vertex of  $T$ . □

Next, we show an upper bound on the number of graphs on  $n$  vertices with a fixed type  $R$ , where  $R$  does not contain a triangle with three grey edges. We use the following definition:

$$\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) := \{G \in \mathcal{G}(n, c) : G \text{ has } (\varepsilon, \varepsilon', d, k')\text{-type } R\}. \tag{2.9}$$

We stress that  $\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)$  and  $\mathcal{R}(R', \varepsilon, \varepsilon', d, k', n, c)$  may have non-empty intersection for  $R \neq R'$ .

**Lemma 2.12.** *For every  $c$  with  $0 < c \leq \frac{1}{2}$ , and every  $\gamma > 0$ , there exist  $\varepsilon_{L2.12}, d_0 > 0$  and integers  $n_{L2.12}, k_0$  such that for all positive  $d \leq d_0$ ,  $\varepsilon \leq \varepsilon_{L2.12}$ ,  $\varepsilon'$ , and all integers  $k \geq k_0$ ,  $k', n \geq \max\{k, n_{L2.12}\}$ , the following holds. If  $R$  is a coloured graph of order  $k$  which has at most  $ek^2$  non-edges and does not contain a triangle with three grey edges, then*

$$|\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)| \leq 2^{h(c)\binom{n}{2} + \gamma\binom{n}{2}}.$$

**Proof.** Let  $c, \gamma$  be given. Choose  $\varepsilon_{L2.12}, d_0, k_0$  such that

$$\max\left\{4\varepsilon_{L2.12}, H(d_0), \frac{1}{k_0}\right\} \leq \frac{\gamma}{5}.$$

Let  $n_{L2.12}$  be large enough to guarantee  $\log(n_{L2.12} + 1) \leq \frac{\gamma}{5}(n_{L2.12} - 1)$ . Let  $\varepsilon \leq \varepsilon_{L2.12}, \varepsilon', d \leq d_0, k \geq k_0, k', n \geq \max\{n_{L2.12}, k\}$  be given.

Let  $R = ([k], E_R, \sigma)$  be a coloured graph which has at most  $ek^2$  non-edges and does not contain a triangle with three grey edges. We shall count the graphs in  $\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)$  by estimating the number of equipartitions  $V_0 \dot{\cup} \dots \dot{\cup} V_k$  of  $[n]$ , the number of choices for edges with one end in the exceptional set  $V_0$  and edges in pairs  $(V_i, V_j)$  such that  $\{i, j\} \notin E_R$ , the number of choices for edges in clusters  $V_i$  such that  $i$  is white or black in  $R$ , and the number of choices for at most  $c\binom{n}{2}$  edges in pairs  $(V_i, V_j)$  such that  $\{i, j\}$  is a white, black, or grey edge of  $R$ .

The number of equipartitions  $V_0 \dot{\cup} \dots \dot{\cup} V_k$  of  $[n]$  is bounded by

$$(k + 1)^n = 2^{n \log(k+1)} \leq 2^{n \log(n+1)} \leq 2^{\frac{\gamma}{5}\binom{n}{2}}. \tag{2.10}$$

Let us now fix such an equipartition. There are at most  $\varepsilon n^2$  possible edges that have at least one end in  $V_0$  and at most  $\frac{\varepsilon}{2} n^2$  possible edges in pairs  $(V_i, V_j)$  such that  $\{i, j\} \notin E_R$ . Thus there are at

most

$$2^{\frac{3}{2}\varepsilon n^2} \leq 2^{4e\binom{n}{2}} \tag{2.11}$$

possible ways to distribute such edges. In addition, the number of ways to distribute edges in clusters  $V_i$  corresponding to white or black vertices of  $R$  is at most

$$2^{k\binom{n/k}{2}} \leq 2^{\frac{1}{k}\binom{n}{2}}. \tag{2.12}$$

By definition, white edges of an  $(\varepsilon, \varepsilon', d, k')$ -type correspond to pairs with density at most  $d$  and black edges correspond to pairs with density at least  $(1 - d)$ . Hence, by the symmetry of the binomial coefficient the number of ways to distribute edges in pairs  $(V_i, V_j)$  such that  $\{i, j\}$  is a white or a black edge of  $R$  is at most

$$\binom{\binom{n}{k}^2}{d\binom{n}{k}}^{\binom{k}{2}} \leq 2^{\binom{k}{2}\binom{n}{k}^2 H(d)} \leq 2^{H(d)\binom{n}{2}}. \tag{2.13}$$

For later reference we now sum up the estimates obtained so far. The product of (2.10)–(2.13) gives less than

$$2^{\left(\frac{\gamma}{5} + 4\varepsilon + H(d) + \frac{1}{k}\right)\binom{n}{2}} \leq 2^{\frac{4}{3}\gamma\binom{n}{2}} \tag{2.14}$$

choices for the partition  $V_0 \dot{\cup} \dots \dot{\cup} V_k$  and for the distribution of edges inside such a partition, except for the pairs  $(V_i, V_j)$  corresponding to grey edges of  $R$ .

It remains to take the grey edges  $E_g$  of  $R$  into account. By assumption,  $E_g$  does not contain a triangle. Hence, by Turán’s theorem (see, e.g., [23]) we have  $|E_g| \leq \frac{k^2}{4}$ . It follows that there are at most  $\frac{k^2}{4} \binom{n}{k}^2 = \frac{n^2}{4}$  possible places for edges in  $E_g$ -pairs  $(V_i, V_j)$ , i.e., pairs such that  $\{i, j\} \in E_g$ . Hence the number  $N_g$  of possible ways to distribute at most  $c\binom{n}{2}$  edges to  $E_g$ -pairs is at most  $c\binom{n}{2} \binom{n^2/4}{c\binom{n}{2}}$ . If  $c < \frac{1}{4}$ , then this gives

$$N_g \leq 2^{\frac{1}{2}\binom{n}{2}H(2c) + \gamma\binom{n}{2}}, \tag{2.15}$$

and if  $\frac{1}{4} \leq c \leq \frac{1}{2}$  then

$$N_g \leq 2^{\frac{n^2}{4}} \leq 2^{\frac{1}{2}\binom{n}{2} + \frac{\gamma}{5}\binom{n}{2}}. \tag{2.16}$$

Combining (2.15) and (2.16) and recalling the definition of  $h(c)$  in (1.2) gives

$$N_g \leq 2^{h(c)\binom{n}{2} + \frac{\gamma}{5}\binom{n}{2}}. \tag{2.17}$$

Multiplying (2.14) and (2.17) gives the desired upper bound

$$|\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)| \leq 2^{h(c)\binom{n}{2} + \frac{\gamma}{5}\binom{n}{2} + \frac{4}{3}\gamma\binom{n}{2}} = 2^{h(c)\binom{n}{2} + \gamma\binom{n}{2}}. \quad \square$$

With this we are in position to prove Lemma 2.10.

**Proof of Lemma 2.10.** Observe first that it suffices to prove the lemma for  $c \leq \frac{1}{2}$ , as the complement of a graph without induced  $C_5$  is induced  $C_5$ -free and hence  $|\mathcal{C}(n, c)| = |\mathcal{C}(n, 1 - c)|$ .

Now let  $c \in (0, \frac{1}{2}]$  and  $\gamma > 0$  be given. Set  $k' = 5$ . Lemma 2.12 with input  $c$  and  $\gamma/2$  provides constants  $\varepsilon_{L2.12}$ ,  $d_0$ ,  $n_{L2.12}$ ,  $k_0$ . Set  $d = d_0$ . From Lemma 2.11 with input  $d$  we obtain constants

$\varepsilon_{L2.11}$  and  $\varepsilon'_{L2.11}$ . Set  $\varepsilon := \min\{\varepsilon_{L2.12}, \varepsilon_{L2.11}\}$  and  $\varepsilon' := \varepsilon'_{L2.11}$ . The type lemma, Lemma 2.7, finally, with input  $\varepsilon, \varepsilon'$ , and  $k_0, k'$  gives constants  $k_1$  and  $n_{L2.7}$ . We set  $n_0 = \max\{n_{L2.7}, n_{L2.12}, \frac{3}{\sqrt{\gamma}}k_1\}$ .

Now, for each graph  $G \in \mathcal{C}(n, c)$  we apply the type lemma, Lemma 2.7, with parameters  $\varepsilon, \varepsilon', k_0, k'$  and  $d$  and obtain an  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  of  $G$  on  $k \leq k_1$  vertices and with at most  $\varepsilon k^2$  non-edges. Let  $\tilde{\mathcal{R}}$  be the set of types obtained from these applications of Lemma 2.7. It follows that  $|\tilde{\mathcal{R}}| \leq 4^{\binom{k_1}{2}} 2^{k_1} \leq 2^{k_1^2}$ . By Lemma 2.11 applied with  $d, \varepsilon$ , and  $\varepsilon'$ , no coloured graph in  $\tilde{\mathcal{R}}$  contains a triangle with three grey edges. Hence, by Lemma 2.12 applied with  $c, \gamma/2, \varepsilon, \varepsilon'$  and  $d$  we have  $|\mathcal{R}(R, \varepsilon, \varepsilon', d, 5, n, c)| \leq 2^{h(c)\binom{n}{2} + \frac{1}{2}\gamma\binom{n}{2}}$ . Since, by Lemma 2.7,

$$\mathcal{C}(n, c) \subseteq \bigcup_{R \in \tilde{\mathcal{R}}} \mathcal{R}(R, \varepsilon, \varepsilon', d, 5, c, n).$$

We conclude from the choice of  $n_0$  that

$$|\mathcal{C}(n, c)| \leq 2^{k_1^2} \cdot 2^{h(c)\binom{n}{2} + \frac{1}{2}\gamma\binom{n}{2}} \leq 2^{h(c)\binom{n}{2} + \gamma\binom{n}{2}}. \quad \square$$

### 3. The proof of Theorem 1.3

Our proof of Theorem 1.3 consists of the following steps. We start, as in the proof of Theorem 1.2, by constructing for each graph  $G$  in  $\text{Forb}_{n,\eta}^*(C_5, c)$  a type  $R$  of size independent of  $n$  with the help of the type lemma, Lemma 2.7. Next, we consider each cluster  $V_i$  of a partition of  $V(G)$  corresponding to  $R$  separately. We shall show that the fact that  $G$  does not contain homogeneous sets of size  $\eta n$  implies that  $G[V_i]$  has many vertex-disjoint induced copies of  $P_3$ , the path on three vertices, or many vertex-disjoint induced copies of the anti-path  $\bar{P}_3$ , the complement of  $P_3$  (see Lemma 3.1). Many induced copies of  $P_3$  or  $\bar{P}_3$  in two clusters  $V_i$  and  $V_j$ , however, limit the number of possibilities for inserting edges between  $V_i$  and  $V_j$  without inducing a  $C_5$  (see Lemma 3.2). Combining this with the proof strategy from Theorem 1.2 will give us an upper bound for the number of graphs from  $\text{Forb}_{n,\eta}^*(C_5, c)$  with type  $R$  (see Lemma 3.3). Finally, comparing this upper bound with the lower bound on  $|\text{Forb}_n^*(C_5, c)|$  from Theorem 1.2 will lead to the desired result.

We start by proving that graphs without big homogeneous sets contain many vertex-disjoint induced  $P_3$  or  $\bar{P}_3$ .

**Lemma 3.1.** *Let  $G$  be a graph of order  $n$  with  $\text{hom}(G) \leq n/6$ . Then one of the following is true:*

- (i)  $G$  contains  $n/6$  vertex-disjoint induced copies of  $P_3$ , or
- (ii)  $G$  contains  $n/6$  vertex-disjoint induced copies of  $\bar{P}_3$ .

**Proof.** Let  $G$  be an  $n$ -vertex graph with  $\text{hom}(G) \leq n/6$ . Select a maximal set of disjoint copies of  $P_3$ . If this set consists of less than  $n/6$  paths then there is a subgraph  $G' \subseteq G$  with  $v(G') = n/2$  that has no induced  $P_3$  and thus is a vertex-disjoint union of cliques  $Q_1, \dots, Q_\ell$ . We claim that in  $G'$  we can find  $n/6$  vertex-disjoint induced  $\bar{P}_3$ , which proves the lemma.

Indeed, since  $\text{hom}(G) \leq n/6$  we have  $\ell \leq n/6$ , and for each  $i \in [\ell]$  we have  $q_i := |Q_i| \leq n/6$ . This implies

$$\sum_{i \in [\ell]} \lfloor q_i/2 \rfloor \geq \frac{1}{2} \left( \frac{n}{2} - \frac{n}{6} \right) = n/6.$$

It follows that we can find a set of  $n/6$  vertex-disjoint edges  $E = \{e_1, \dots, e_{n/6}\}$  in these cliques in the following way. We first choose as many vertex-disjoint edges in  $Q_1$  as possible, then in  $Q_2$ , and so on, until we have chosen  $n/6$  edges in total. Let  $Q_k$  be the last clique used in this process. Then, for each clique  $Q_i$  with  $i < k$  at most one vertex was unused in this process, and in  $Q_k$  possibly several vertices were unused. Let  $X$  be the set of all these unused vertices together with all vertices from  $\bigcup_{k < i \leq \ell} Q_i$ . Clearly  $|X| = n/6$ .

We consider the auxiliary bipartite graph  $B = (X \cup E, E_B)$  with  $\{x, e\} \in E_B$  for  $x \in X$  and  $e \in E$  if and only if  $x$  and  $e$  do not lie in the same clique of  $G'$ . We verify Hall's condition for  $B$ . So let  $Y \subseteq X$ . If  $Y \not\subseteq Q_i$  for all  $i$  we have  $|N(Y)| = |E| = |X| \geq |Y|$ . Otherwise, if  $Y \subseteq Q_i$  for some  $i$ , then  $|N(Y)| \geq |E| - (|Q_i| - |Y|)/2 \geq n/6 - n/12 + |Y|/2 \geq |Y|$  since  $|Y| \leq n/6$ . It follows that  $B$  has a perfect matching, which means that there are  $n/6$  vertex-disjoint induced  $\bar{P}_3$  in  $G'$ , as claimed. □

Now suppose we are given a graph  $G$  with vertex set  $V_1 \dot{\cup} V_2$  and no edges between  $V_1$  and  $V_2$ . Further, let  $H_1$  and  $H_2$  be such that for  $i \in [2]$  the graph  $H_i$  induces a copy of  $P_3$  or  $\bar{P}_3$  in  $G[V_i]$ . Observe that, no matter which combination of  $P_3$  or  $\bar{P}_3$  we choose, we can create an induced  $C_5$  in  $G$  by adding appropriate edges between  $H_1$  and  $H_2$ . Since we are interested in graphs without induced  $C_5$ , this motivates us to call  $(H_1, H_2)$  a *dangerous pair* of  $(V_1, V_2)$ .

Our next goal is to use these dangerous pairs in order to derive an upper bound on the number of possibilities for inserting edges between  $V_1$  and  $V_2$  without creating an induced copy of  $C_5$  if we know that  $(V_1, V_2)$  contains many dangerous pairs. In order to quantify this upper bound, in Lemma 3.2 we use the following technical definition. We define  $R(c) = c^4(1 - c)^4$  and the function  $r : (0, 1) \rightarrow \mathbb{R}^+$  with

$$r(c) = \frac{1}{72} \begin{cases} R(2c) & \text{if } c < \frac{1}{4}, \\ (1/4)^4 & \text{if } c \in [\frac{1}{4}, \frac{3}{4}], \\ R(2c - 1) & \text{otherwise.} \end{cases} \tag{3.1}$$

Recall in addition the definition of the function  $h(c)$  from (1.2).

**Lemma 3.2.** *For every  $0 < c_0 \leq \frac{1}{2}$  there is an  $n_0$  such that, for all  $c$  with  $c_0 \leq 2c \leq 1 - c_0$  and  $n \geq n_0$ , the following holds. Let  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  be two  $n$ -vertex graphs, each of which contains  $n/6$  vertex-disjoint induced copies of  $P_3$  or  $n/6$  vertex-disjoint induced copies of  $\bar{P}_3$ . Let  $G = (V_1 \dot{\cup} V_2, E)$  be the disjoint union of  $G_1$  and  $G_2$ . Then there are at most*

$$2^{2n^2(h(c)-r(c))}$$

*ways to add exactly  $2cn^2$  edges to  $G$  that run between  $V_1$  and  $V_2$  without inducing a  $C_5$  in  $G$ .*

We remark that in the proof of this lemma we are going to make use of the following probabilistic principle: we can *count* the number of elements in a finite set  $X$  which have some

property  $P$ , by determining the *probability* that an element which is chosen from  $X$  uniformly at random has property  $P$ .

**Proof of Lemma. 3.2.** Given  $c_0 \in (0, \frac{1}{2}]$ , let  $n_0$  be sufficiently large that

$$n_0 e^{-2r(c_0/2)n_0^2} \leq 2^{-2r(c_0/2)n_0^2}. \tag{3.2}$$

Now let  $c$  be such that  $c_0 \leq 2c \leq 1 - c_0$ . Observe first that it suffices to prove the lemma for  $2c \leq \frac{1}{2}$ , since induced  $C_5$ -free graphs are self-complementary and  $P_3$  is the complement of  $\bar{P}_3$ . Hence, we assume from now on that  $2c \leq \frac{1}{2}$ . Observe moreover that (3.2) remains valid if  $c_0$  is replaced by  $c$ , since  $r(c)$  is monotone increasing in  $[0, \frac{1}{2}]$ . Let  $G_1, G_2$ , and  $G$  be as required.

Our first goal is to estimate the probability  $P^*$  of inducing no  $C_5$  in  $G$  when choosing uniformly at random exactly  $2cn^2$  edges between  $V_1$  and  $V_2$ . Instead of dealing with  $P^*$  directly, we consider the following binomial random graph  $\mathcal{G}(V_1, V_2, p)$  with  $p = 2c$ : we start with  $G$  and add each edge between  $V_1$  and  $V_2$  independently with probability  $p$ .

Now, let  $A$  be the event that  $\mathcal{G}(V_1, V_2, p)$  contains exactly  $2cn^2$  edges between  $V_1$  and  $V_2$ , and let  $B$  be the event that  $\mathcal{G}(V_1, V_2, p)$  contains no induced  $C_5$ . Observe that each graph with  $2cn^2$  edges between  $V_1$  and  $V_2$  is equally likely in  $\mathcal{G}(V_1, V_2, p)$ , and thus

$$P^* = \mathbb{P}[B|A] \leq \frac{\mathbb{P}[B]}{\mathbb{P}[A]}. \tag{3.3}$$

Hence it suffices to estimate  $\mathbb{P}[A]$  and  $\mathbb{P}[B]$ .

We first bound  $\mathbb{P}[B]$ . By assumption there are at least  $n^2/36$  dangerous pairs in  $(V_1, V_2)$ . Now fix such a dangerous pair  $(H_1, H_2)$ . The probability that  $(H_1, H_2)$  induces a  $C_5$  in  $\mathcal{G}(V_1, V_2, p)$  is at least  $p^2(1-p)^4$  unless  $H_1$  and  $H_2$  are both  $\bar{P}_3$ , and at least  $p^4(1-p)^2$  unless  $H_1$  and  $H_2$  are both  $P_3$ . Thus  $(H_1, H_2)$  induces a copy of  $C_5$  with probability at least

$$p^4(1-p)^4 = (2c)^4(1-2c)^4 \stackrel{(3.1)}{\geq} 72 \cdot r(c).$$

Since we can upper-bound the probability of  $B$  by the probability that none of the  $n^2/36$  dangerous pairs in  $(V_1, V_2)$  induces a  $C_5$  in  $\mathcal{G}(V_1, V_2, p)$ , we obtain

$$\mathbb{P}[B] \leq (1 - 72 \cdot r(c))^{n^2/36} \leq e^{-2r(c)n^2}.$$

Note that the number of edges between  $V_1$  and  $V_2$  in  $\mathcal{G}(V_1, V_2, p)$  is binomially distributed. Thus, by Stirling’s formula,  $\mathbb{P}[A] \geq 1/(\sqrt{2\pi p(1-p)}n) \geq 1/n$ . By the choice of  $n_0$ , combining this with (3.3) gives

$$P^* \leq n e^{-2r(c)n^2} \stackrel{(3.2)}{\leq} 2^{-2r(c)n^2}. \tag{3.4}$$

It remains to estimate the number  $N$  of ways to choose exactly  $2cn^2$  edges between  $V_1$  and  $V_2$ . We have

$$N \leq \binom{n^2}{2cn^2} \stackrel{(2.1)}{\leq} 2^{2h(c)n^2}. \tag{3.5}$$



This implies that the number of ways to add exactly  $2cn^2$  edges to  $G$  that run between  $V_1$  and  $V_2$  without inducing a  $C_5$  is

$$P^* \cdot N \stackrel{(3.4),(3.5)}{\leq} 2^{-2r(c)n^2} \cdot 2^{2h(c)n^2}. \quad \square$$

Next, we want to show that Lemma 3.2 allows us to derive an upper bound on the number of graphs  $G$  such that (a)  $G$  has no large homogeneous sets and (b)  $G$  has a fixed type  $R$  which does not contain a triangle with three grey edges. Our aim is to obtain an upper bound which is much smaller than the bound provided in Lemma 2.12 for the corresponding problem without restriction (a). Lemma 3.3 states that this is possible. Recall for this purpose the definition of  $\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)$  from (2.9).

**Lemma 3.3.** *For every  $c$  with  $0 < c \leq \frac{1}{2}$ , and every  $\gamma > 0$ , there exist  $\varepsilon_0, d_0 > 0$  and an integer  $k_0$  such that, for all integers  $k_1 \geq k_0, k'$ , there is an integer  $n_0$  such that, for all positive  $d \leq d_0, \varepsilon \leq \varepsilon_0, \varepsilon'$ , and all integers  $n \geq n_0$  and  $k_0 \leq k \leq k_1$ , the following holds. If  $R$  is a coloured graph of order  $k$  which has at most  $ek^2$  non-edges and does not contain a triangle with three grey edges, and  $\eta = 1/(6k_1)$ , then*

$$|\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) \cap \text{Forb}_{n,\eta}^*(C_5, c)| \leq 2^{(h(c)-r(c))\binom{n}{2} + \gamma\binom{n}{2}}.$$

In the proof of this lemma we combine the strategy of the proof of Lemma 2.12 with an application of Lemma 3.1 to all clusters of a partition corresponding to  $R$ , and an application of Lemma 3.2 to regular pairs of medium density. We shall make use of the following observation.

Using the definition of  $r(c)$  from (3.1), it is easy to check that  $f(c) := h(c) - r(c)$  is a concave function for  $c \in (0, 1)$ . Thus  $f$  enjoys the following property, which is a special form of Jensen's inequality (see, e.g., [13]).

**Proposition 3.4 (Jensen's inequality).** *Let  $f$  be a concave function,  $0 < c < 1, 0 < c_i < 1$  for  $i \in [m]$  and let  $\sum_{i=1}^m c_i = mc$ . Then*

$$\sum_{i=1}^m f(c_i) \leq m \cdot f(c). \quad \square$$

**Proof of Lemma 3.3.** Let  $c, \gamma > 0$  be given and choose  $\varepsilon_0, d_0, k_0$  and  $n_0$  as in the proof of Lemma 2.12. Let  $d \leq d_0$  and  $k_1$  be given, and possibly increase  $n_0$  so that  $n_0 \geq 2k_1 n_{L3.2}$ , where  $n_{L3.2}$  is the constant from Lemma 3.2 with parameter  $d$ , and so that

$$\frac{3}{4}k_1^2 \log n_0 + 2n_0 \leq \frac{\gamma}{10} \binom{n_0}{2}. \tag{3.6}$$

If necessary decrease  $\varepsilon_0$  so that

$$3\varepsilon_0 \log \frac{1}{c} \leq \frac{\gamma}{10}. \tag{3.7}$$

Let  $\varepsilon \leq \varepsilon_0, \varepsilon', n \geq n_0$ , and  $k$  with  $k_0 \leq k \leq k_1$  be given.

Let  $R = ([k], E_R, \sigma)$  be a coloured graph which has at most  $ek^2$  non-edges and does not contain a triangle with three grey edges. In the proof of Lemma 2.12 we counted the number of graphs in  $\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c)$  by estimating the number of equipartitions  $V_0 \dot{\cup} \dots \dot{\cup} V_k$  of  $[n]$ , the number of choices for edges with one end in the exceptional set  $V_0$  and edges in pairs  $(V_i, V_j)$  such that  $\{i, j\} \notin E_R$ , the number of choices for edges inside clusters  $V_i$  (such that  $i$  is white or black in  $R$ ), and the number of choices for at most  $c\binom{n}{2}$  edges in pairs  $(V_i, V_j)$  such that  $\{i, j\}$  is a white, black, or grey edge of  $R$ . Now we are interested in the number of graphs in  $\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) \cap \text{Forb}_{n,\eta}^*(C_5, c)$ . Clearly we can use the same strategy, and it is easy to verify that the estimates in (2.10)–(2.13) and thus in (2.14) from the proof of Lemma 2.12 remain valid in this setting. From now on, as in the proof of Lemma 2.12, we fix a partition  $V_0 \dot{\cup} \dots \dot{\cup} V_k$  of  $[n]$  and observe that also  $m_g := |E_g| \leq \frac{k^2}{4}$  still holds for the grey edges  $E_g$  in  $R$ . However, we shall now use Lemma 3.2 to obtain an improved bound on the number of possible choices for edges in  $E_g$ -pairs  $(V_i, V_j)$ , and use this to replace (2.17) by a smaller bound on the number  $N_g$  of possible ways to distribute at most  $c\binom{n}{2}$  edges to  $E_g$ -pairs. Since in the following we do not rely on any interferences between different  $E_g$ -pairs, clearly  $N_g$  will be maximal if  $m_g$  is maximal, and hence we assume from now on that

$$m_g = \frac{k^2}{4}. \tag{3.8}$$

Let  $s := |V_1| = \dots = |V_k|$  and observe that

$$\frac{n}{k} \geq s \geq (1 - \varepsilon) \frac{n}{k} \geq n_{\text{L3.2}}. \tag{3.9}$$

By Lemma 3.1, for each cluster  $V_i$  of a partition  $P$  of a graph in

$$\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) \cap \text{Forb}_{n,\eta}^*(C_5, c)$$

such that  $P$  corresponds to  $R$ , we have that  $V_i$  contains either  $s/6$  copies of  $P_3$  or  $s/6$  copies of  $\bar{P}_3$ . Hence we will assume from now on that in our fixed partition the clusters  $V_i$  have this property.

We now upper-bound  $N_g$  by multiplying the possible ways  $A$  to assign at most  $c\binom{n}{2}$  edges to one  $E_g$ -pair each, and the maximum number  $B$  of ways to choose all these assigned edges in the corresponding pairs, without inducing a  $C_5$ . First observe that we have

$$A \leq \left( c \binom{n}{2} \right)^{m_g+1} \leq n^{3m_g} \stackrel{(3.8)}{\leq} 2^{\frac{3}{4}k^2 \log n}. \tag{3.10}$$

For estimating  $B$ , we now assume that we have fixed an assignment in which each pair  $(V_i, V_j)$  with  $\{i, j\} \in E_g$  is assigned  $2c_{i,j}s^2$  edges. Let  $\hat{c}$  be such that

$$\sum_{\{i,j\} \in E_g} 2c_{i,j}s^2 =: \hat{c}n^2 \leq c \binom{n}{2}. \tag{3.11}$$

Observe further that, since  $\{i, j\} \in E_g$  is a grey edge of  $R$ , and we are interested in counting graphs with a partition corresponding to  $R$ , we can assume that  $d \leq 2c_{i,j} \leq (1 - d)$ . Hence, by (3.9) we can apply Lemma 3.2 with  $c_0 = d$  to infer that for each  $\{i, j\} \in E_g$  there are at most

$$B_{ij} \leq 2^{2s^2(h(c_{i,j})-r(c_{i,j}))} \tag{3.12}$$

possible ways to choose the  $2c_{i,j}s^2$  edges in  $(V_i, V_j)$  without inducing a  $C_5$ . Now let  $2\tilde{c} := \sum_{\{i,j\} \in E_g} 2c_{i,j}/m_g$  and observe that

$$\tilde{c} \stackrel{(3.11)}{=} \hat{c} \frac{n^2}{2s^2m_g} \stackrel{(3.9)}{\leq} \hat{c} \frac{k^2}{2(1-\varepsilon)^2m_g} \stackrel{(3.8)}{\leq} 2(1+3\varepsilon)\hat{c} \stackrel{(3.11)}{\leq} (1+3\varepsilon)c \leq \frac{3}{4}.$$

Therefore, since  $f(x) := h(x) - r(x)$  is a concave function for  $x \in (0, 1)$ , which is moreover non-decreasing for  $x \leq \frac{3}{4}$ , we can infer from Lemma 3.4 that

$$\sum_{\{i,j\} \in E_g} f(c_{i,j}) \leq m_g \cdot f(\tilde{c}) \leq m_g \cdot f(c(1+3\varepsilon)) \stackrel{(3.8)}{=} \frac{k^2}{4} f(c(1+3\varepsilon)). \tag{3.13}$$

As  $h(x)$  is a convex function with  $h'(x) \leq \log(1/x)$  and  $r(x)$  is non-decreasing for  $x \leq 3/4$ , we have

$$f(c(1+3\varepsilon)) \leq h(c+3\varepsilon) - r(c) \leq h(c) + 3\varepsilon h'(c) - r(c) \stackrel{(3.7)}{\leq} h(c) - r(c) + \frac{\gamma}{10}.$$

Together with (3.12) and (3.13), this implies

$$B = \prod_{\{i,j\} \in E_g} B_{ij} \leq 2^{2s^2(k^2/4)(h(c)-r(c)+\gamma/10)} \stackrel{(3.9)}{\leq} 2^{(n^2/2)(h(c)-r(c)+\gamma/10)},$$

which in turn, together with (3.10), gives

$$N_g \leq 2^{\frac{3}{4}k^2 \log n} \cdot 2^{(n^2/2)(h(c)-r(c)+\gamma/10)} \stackrel{(3.6)}{\leq} 2^{(h(c)-r(c))\binom{n}{2} + (\gamma/5)\binom{n}{2}}.$$

By multiplying this with (2.14) from the proof of Lemma 2.12, we obtain

$$|\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) \cap \text{Forb}_{n,1/6k}^*(C_5, c)| \leq 2^{(h(c)-r(c))\binom{n}{2} + (\gamma)\binom{n}{2}},$$

as claimed. □

Lemma 3.3 with the type lemma, Lemma 2.7, implies an upper bound on  $|\text{Forb}_{n,\eta}^*(C_5, c)|$ . Now we can combine this with the lower bound on  $|\text{Forb}_n^*(C_5, c)|$ , which follows from Lemma 2.1, in order to prove Theorem 1.3.

**Proof of Theorem 1.3.** Observe first that, since  $C_5$  is self-complementary, it suffices to prove Theorem 1.3 for  $c \leq 1/2$ . Hence we assume  $c \leq 1/2$  from now on.

We first need to set up some constants. Given  $c \in (0, \frac{1}{2}]$ , we choose  $\gamma > 0$  such that  $2\gamma < r(c)$ . For input  $c$  and  $\gamma$  Lemma 2.1 supplies us with a constant  $n_{L2.1}$ . We apply Lemma 3.3 with input  $c$  and  $\gamma/2$  to obtain  $\varepsilon_{L3.3}$ ,  $d_0$ , and  $k_0$ . Next, we apply Lemma 2.11 with input  $d_0$  and obtain constants  $\varepsilon_{L2.11}$  and  $\varepsilon'$ . Let  $\varepsilon := \min\{\varepsilon_{L3.3}, \varepsilon_{L2.11}\}$ . For input  $\varepsilon, \varepsilon'$ , and  $k_0$ , Lemma 2.7 returns constants  $k_1$  and  $n_{L2.7}$ . With this parameter  $k_1$  we continue the application of Lemma 3.3 and obtain  $n_{L3.3}$ . Choose  $n_0 := \max\{n_{L2.1}, n_{L2.7}, n_{L3.3}, \frac{3}{\sqrt{\gamma}}k_1\}$ , assume that  $n \geq n_0$ , and set  $\eta := 1/(6k_1)$ .

Now, for each graph  $G \in \text{Forb}_{n,\eta}^*(C_5, c)$  we apply the type lemma, Lemma 2.7, with parameters  $\varepsilon, \varepsilon', k_0, k'$  and  $d$ , and obtain an  $(\varepsilon, \varepsilon', d, k')$ -type  $R$  of  $G$  on  $k$  vertices with  $k_0 \leq k \leq k_1$  and with at most  $\varepsilon k^2$  non-edges. Let  $\tilde{\mathcal{R}}$  be the set of types obtained from these applications of Lemma 2.7. It follows that  $|\tilde{\mathcal{R}}| \leq 4^{\binom{k}{2}} 2^{k_1} \leq 2^{k_1^2}$ . By Lemma 2.11 applied with  $d_0, \varepsilon$ , and  $\varepsilon'$ , no coloured graph in  $\tilde{\mathcal{R}}$  contains a triangle with three grey edges. Hence, by Lemma 3.3 applied with  $c, \gamma/2, \varepsilon, \varepsilon'$ ,

$n$ , and  $k$  we have

$$|\mathcal{R}(R, \varepsilon, \varepsilon', d, k', n, c) \cap \text{Forb}_{n,\eta}^*(C_5, c)| \leq 2^{(h(c)-r(c))\binom{n}{2} + \frac{1}{2}\gamma\binom{n}{2}}.$$

Since, by Lemma 2.7,

$$\text{Forb}_{n,\eta}^*(C_5, c) \subseteq \bigcup_{R \in \tilde{\mathcal{R}}} (\mathcal{R}(R, \varepsilon, \varepsilon', d, k', c, n) \cap \text{Forb}_{n,\eta}^*(C_5, c)),$$

we conclude from the choice of  $n_0$  that

$$|\text{Forb}_{n,\eta}^*(C_5, c)| \leq 2^{k_1^2} \cdot 2^{(h(c)-r(c))\binom{n}{2} + \frac{1}{2}\gamma\binom{n}{2}} \leq 2^{(h(c)-r(c))\binom{n}{2} + \gamma\binom{n}{2}}.$$

On the other hand, by Lemma 2.1 and the choice of  $n_0$  we have

$$|\text{Forb}_n^*(C_5, c)| \geq 2^{h(c)\binom{n}{2} - \gamma\binom{n}{2}}.$$

Since  $2\gamma < r(c)$ , by the choice of  $\gamma$ , this implies that almost all graphs in  $\text{Forb}_n^*(C_5, c)$  satisfy  $\text{hom}(G) \geq \eta n$ . □

### 4. Concluding remarks

#### $\mathcal{G}(n, m)$ versus $\mathcal{G}(n, p)$

Our counting results can be interpreted as probabilities in the random graph model  $\mathcal{G}(n, m = c\binom{n}{2})$ . We can reformulate Theorem 1.2 as

$$\mathbb{P}\left[\mathcal{G}\left(n, m = c\binom{n}{2}\right) \in \text{Forb}_n^*(C_5)\right] = \frac{|\text{Forb}_n^*(C_5, c)|}{|\mathcal{G}(n, c)|} = 2^{(h(c)-H(c)+o(1))\binom{n}{2}}.$$

We now compare  $\mathcal{G}(n, m = c\binom{n}{2})$  to the standard Erdős–Rényi model studied by Marchant and Thomason in [17]. They showed that

$$\mathbb{P}[\mathcal{G}(n, p) \in \text{Forb}_n^*(C_5)] = 2^{c_p\binom{n}{2}},$$

where  $c_p = \frac{1}{2} \max\{\log p, \log(1 - p)\}$ . We can now derive the same estimate via Theorem 1.2. Obviously  $\mathbb{P}[\mathcal{G}(n, p) \in \text{Forb}_n^*(C_5)]$  equals

$$\max_c \mathbb{P}\left[e(\mathcal{G}(n, p)) = c\binom{n}{2}\right] \cdot \mathbb{P}\left[\mathcal{G}\left(n, m = c\binom{n}{2}\right) \in \text{Forb}_n^*(C_5)\right] \cdot 2^{o(n^2)}.$$

Setting  $g_p(c) = H(c) + c \log p + (1 - c) \log(1 - p)$ , we obtain

$$\mathbb{P}\left[e(\mathcal{G}(n, p)) = c\binom{n}{2}\right] = 2^{(g(c)+o(1))\binom{n}{2}},$$

and thus, by Theorem 1.2,

$$\mathbb{P}[\mathcal{G}(n, p) \in \text{Forb}_n^*(C_5)] = 2^{(\max_c g_p(c) + h(c) - H(c) + o(1))\binom{n}{2}}.$$

The maximum is attained at  $c = p/2$  for  $p < 1/2$  and  $c = (p + 1)/2$  for  $p > 1/2$ . For  $p = 1/2$  all values  $c \in [1/4, 3/4]$  are optimal. Inserting the optimal value for  $c$  shows that the exponent is indeed equal to  $c_p$  as computed in [22]. This indicates that, for example, a graph from  $\mathcal{G}(n, 1/4)$  that happens to be induced  $C_5$ -free will a.a.s. have  $(1/8 + o(1))\binom{n}{2}$  edges. Thus the  $\mathcal{G}(n, m)$  model can also be used to derive that a typical element having a certain property might be far from a typical element in  $\mathcal{G}(n, p)$ .

## Extensions

The final thing to consider is whether more can be said about  $\text{Forb}_n^*(F, c)$  in the  $\mathcal{G}(n, m)$  model. Indeed, there are at least three natural ways to enhance our results. First one might want to count  $\text{Forb}_n^*(F, c)$  for graphs  $F$  other than  $C_5$ . But already the case of a forbidden induced  $C_7$  is more challenging, as tight upper bounds as in Lemma 2.12 are not so easy to derive. Furthermore,  $h(c)$  does not seem to be unimodal for  $C_7$ .

Second, it would be interesting to obtain even sharper asymptotic bounds for the speed of  $\text{Forb}_n^*(F, c)$ . Alon, Balogh, Bollobás and Morris [2] determine the speed of some hereditary properties up to a subquadratic term in the exponent. We believe their techniques can be extended to the case of restricted density.

Finally one might want to prove a much better constant in the size of the linear homogeneous sets that can be found in almost all graphs in  $\text{Forb}_n^*(C_5, c)$ . It easily follows from the fact that almost all graphs in  $\text{Forb}_n^*(C_5)$  are generalized split graphs (see [19]) that almost all of them also have a homogeneous set of size  $(1/2 - o(1))n$ . We believe that the same is true for the density-restricted case, and that this can be proved as in Theorem 1.3 combined with a stability-type argument. We plan to return to this in the near future.

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