## THE FREE PSEUDOSPACE IS N-AMPLE, BUT NOT (N + 1)-AMPLE

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**Abstract.** We give a uniform construction of free pseudospaces of dimension *n* extending work in [1]. This yields examples of  $\omega$ -stable theories which are *n*-ample, but not n + 1-ample. The prime models of these theories are buildings associated to certain right-angled Coxeter groups.

§1. Introduction. In the investigation of geometries on strongly minimal sets the notion of ampleness plays an important role. Algebraically closed fields are n-ample for all n and it is not known whether there are strongly minimal sets which are n-ample for all n and do not interpret an infinite field. Obviously, one way of proving that no infinite field is interpretable in a theory is by showing that the theory is *not* n-ample for some n.

In [1], Baudisch and Pillay constructed a free pseudospace of dimension 2. Its theory is  $\omega$ -stable (of infinite rank) and 2-ample. F. Wagner posed the question whether this example was 3-ample or not.

In Section 2 we give a uniform construction of a free pseudospace of dimension n and show that it is n-ample, but not n + 1-ample. It turns out that the theory of the free pseudospace of dimension n is the first order theory of a Tits-building associated to a certain Coxeter diagram and we will investigate this connection in Section 4.

In the final section we determine the orthogonality classes of regular types.

The construction given here is quite similar to the one given by Evans in [2] for a stable theory which is n-ample for all n, but does not interpret an infinite group. In contrast to the examples constructed by Evans, our theory is trivial and no infinite group is definable.

§2. Construction and results. Fix a natural number  $n \ge 1$ . Let  $L_n$  be the language for n + 1-colored graphs containing predicates  $V_i, i = 0, ..., n$  and an edge relation E. If  $x \in V_i$  we also say that x is of level i.

By an  $L_n$ -graph we mean an n + 1-colored graph with vertices of types  $V_i$ , i = 0, ..., n and an edge relation  $E \subseteq \bigcup_{i=1,...,n} V_{i-1} \times V_i$ . We say that a path in this graph is of type  $E_i$  if all its vertices are in  $V_{i-1} \cup V_i$  and of type  $E_i \cup ... \cup E_{i+j}$  if all its vertices are in  $V_{i-1} \cup ... \cup V_{i+j}$ .

© 2014, Association for Symbolic Logic 0022-4812/14/7902-0003/\$2.90 DOI:10.1017/jsl.2014.42

Received October 5, 2011.

Key words and phrases. Stable theory, ampleness, right-angled buildings.

The free pseudospaces will be modeled along the lines of a projective space as a simplicial complex, i.e., we will think of vertices of type  $V_i$  as *i*-dimensional spaces in a free pseudospace. Therefore we extend the notion of incidence as follows:

DEFINITION 2.1.

- 1. We say that vertices  $x_i, x_j$  of type  $V_i$  and  $V_j$ , respectively, are *incident* (for  $i \leq j$ ) if there are vertices  $x_\ell$  of type  $V_\ell, \ell = i + 1 \dots j$ , such that  $E(x_{\ell-1}, x_\ell)$  holds. In this case the sequence  $(x_i, \dots x_j)$  is called a *dense flag*. A *flag* is a sequence of vertices  $(x_1, \dots x_k), k \geq 0$  in which any two vertices are incident (and hence no two vertices of a flag have the same level). (The levels of a flag may be increasing or decreasing.) In particular, a vertex x is incident with itself.
- 2. The *residue* R(x) of a vertex x is the set of vertices incident with x. We write  $R_{<}(x)$  (and  $R_{>}(x)$ , respectively) for the elements in R(x) of level less (greater, respectively) than the level of x. Similarly for  $R_{<}(x)$  and  $R_{>}(x)$
- 3. We say that two vertices x and y intersect in the vertex z, and write  $z = x \land y$ , if  $R_{\leq}(x) \cap R_{\leq}(y) = R_{\leq}(z)$ . If  $R_{\leq}(x) \cap R_{\leq}(y) = \emptyset$ , we say that x and y intersect in the empty set.
- 4. Similarly, we say that two vertices x and y generate the vertex z, and write  $z = x \lor y$ , if  $R_{\geq}(x) \cap R_{\geq}(y) = R_{\geq}(z)$ . If  $R_{\geq}(x) \cap R_{\geq}(y) = \emptyset$ , we say that x and y generate the empty set.
- 5. A simple cycle is a cycle without repetitions.

We now give an inductive definition of a free pseudospace of dimension *n*:

DEFINITION 2.2. A free pseudospace of dimension 0 is an infinite set of vertices. Assume that a free pseudospace of dimension n - 1 has been defined. Then a free pseudospace of dimension n is an  $L_n$ -graph such that the following holds:

- $(\Sigma 1)_n$  (a) The set of vertices of type  $V_0 \cup \ldots \cup V_{n-1}$  is a free pseudospace of dimension n-1.
  - (b) The set of vertices of type  $V_1 \cup \ldots \cup V_n$  is a free pseudospace of dimension (n-1).
- $(\Sigma 2)_n$  (a) For any vertex x of type  $V_0$ ,  $R_>(x)$  is a free pseudospace of dimension (n-1).
  - (b) For any vertex x of type  $V_n$ ,  $R_{<}(x)$  is a free pseudospace of dimension (n-1).
- (Σ3)<sub>n</sub> (a) Any two vertices x and y intersect in a vertex or the empty set.
  (b) Any two vertices x and y generate a vertex or the empty set.
- $(\Sigma 4)_n$  (a) If *a* is a vertex of type  $V_0$  and  $b, b' \in R_>(a)$  with  $b' \notin R(b)$  are connected by a path  $\gamma$  of length *k* such that for some dense flags  $f = (a, \ldots b)$  and  $f' = (b', \ldots, a)$  the concatenation of these paths  $f \circ \gamma \circ f'$  is a simple cycle, there is a path  $\gamma'$  of length at most *k* from *b* to *b'* in  $R_>(a)$  containing some interior vertex of  $\gamma$  such that  $f \circ \gamma' \circ f'$  is a simple cycle.
  - (b) If a is a vertex of type  $V_n$  and  $b, b' \in R_{<}(a)$  are connected by a path  $\gamma$  of length k such that for some dense flags  $f = (a, \ldots b)$  and  $f' = (b', \ldots, a)$  the concatenation  $f \circ \gamma \circ f'$  of these paths is

a simple cycle, there is a path  $\gamma'$  of length at most k from b to b' in  $R_{\leq}(a)$  containing some interior vertex of  $\gamma$  such that  $f \circ \gamma' \circ f'$  is a simple cycle.

REMARK 2.3. Note that by ( $\Sigma 4$ ) any path  $\gamma = (b = x_0, \dots, x_m = b')$  with  $b, b' \in \mathbb{R}_{<}(a)$  and  $m \geq 2$  contains an interior vertex which lies in  $\mathbb{R}_{<}(a)$  unless  $b' \in \mathbb{R}(b)$  (and dually for  $\mathbb{R}_{>}(a)$ ).

Note that a free pseudospace of dimension 1 is a free pseudoplane, i.e., an  $L_1$ -graph which by  $(\Sigma 4)_1$  does not contain any cycles and such that any vertex has infinitely many neighbors.

Let  $T_n$  denote the  $L_n$ -theory expressing these axioms.

Note that the inductive nature of the definition immediately has the following consequences:

- 1. The induced subgraph on  $V_j \cup \ldots \cup V_{j+m}$  is a free pseudospace of dimension m.
- 2. The notion of a free pseudospace of dimension *n* is *self-dual*: if we put  $W_i = V_{n-i}$ , i = 0, ..., n, then  $W_0, ..., W_n$  with the same set of edges is again a free pseudospace of dimension *n*.
- 3. If  $a \in V_i$ , then  $R_{\leq}(a)$  is a free pseudospace of dimension i 1 (and dually for  $R_{\geq}(a)$ .)

Our first goal is to show that  $T_n$  is consistent<sup>1</sup> and complete.

DEFINITION 2.4. Let A be a finite  $L_n$ -graph. The following extensions are called minimal strong extensions of A:

- 1. Add a vertex of any type to A which is connected to at most one vertex of A of an appropriate type.
- 2. If  $(x_0, x_1, \ldots, x_k, x_{k+1})$  is a dense flag in A, add vertices  $y_1, \ldots, y_k$  such that  $(x_0, y_1, \ldots, y_k, x_{k+1})$  is again a dense flag.

We say that *B* is a strong extension of *A*, written  $A \le B$ , if *B* arises from *A* by a sequence of finitely many minimal strong extensions.

DEFINITION 2.5. Let  $\mathcal{K}_n$  be the class of finite  $L_n$ -graphs A such that the following holds

- 1. The empty graph is strong in A.
- 2. If  $a, a' \in A$  with  $R_{\leq}(a) \cap R_{\leq}(a') \neq \emptyset$ , then there is a unique vertex  $z \in A$  with  $R_{\leq}(a) \cap R_{\leq}(a') = R_{\leq}(z)$ .
- 3. If  $a, a' \in A$  with  $R_{\geq}(a) \cap R_{\geq}(a') \neq \emptyset$ , then there is a unique vertex  $z \in A$  with  $R_{\geq}(a) \cup R_{\geq}(a') = R_{\geq}(z)$ .
- 4. If  $a \in V_i$  and  $b, b' \in R(a)$  are connected by a path  $\gamma$  of length k contained in  $V_{i-m} \cup \ldots \cup V_i$  such that for some dense flags  $f = (a, \ldots b)$  and  $f' = (b', \ldots, a)$  the concatenation of these paths  $f \circ \gamma \circ f'$  is a simple cycle, there is a path  $\gamma'$  from b to b' of length at most k containing some interior vertex of  $\gamma$ and contained in  $R(a) \cap (V_{i-m} \cup \ldots \cup V_{i-1})$  such that  $f \circ \gamma \circ f'$  is a simple cycle.
- 5. If  $a \in V_i$  and  $b, b' \in R(a)$  are connected by a path  $\gamma$  of length k contained in  $V_i \cup \ldots \cup V_{i+m}$  such that for some dense flags  $f = (a, \ldots b)$  and  $f' = (b', \ldots, a)$

<sup>&</sup>lt;sup>1</sup>Readers familiar with buildings will notice that any building of type  $A_{\infty,n+1}$  with infinite valencies is a model of  $T_n$ .

the concatenation of these paths is a simple cycle, there is a path  $\gamma'$  from b to b' of length at most k containing some interior vertex of  $\gamma$  and contained in  $R(a) \cap (V_{i+1} \cup \ldots \cup V_{i+m})$  such that  $f \circ \gamma' \circ f'$  is a simple cycle.

6. If  $a \in V_i, c \in V_{i+j}$  and  $b, b' \in R(a) \cap R(c)$  are connected by a path  $\gamma \subset V_i \cup \ldots V_{i+j}$  of length k, then there is a path  $\gamma' \subset R_{>}(a) \cap R_{<}(c) \subseteq (V_{i+1} \cup \ldots V_{i+j-1})$  from b to b' of length at most k containing some interior vertex of  $\gamma$ .

Note that for the whole structure the property corresponding to 6. follows from  $T_n$ .

**REMARK 2.6.** Note that if  $A \in \mathcal{K}_n$ , then  $A \cap (V_i \cup \ldots \cup V_{i+j}) \in \mathcal{K}_j$ . Conversely, any  $A \in \mathcal{K}_j$  can be considered as a graph in  $\mathcal{K}_n$  for any  $n \ge j$ .

We next show that  $(\mathcal{K}, \leq)$  has the amalgamation property for strong extensions. This will be enough to obtain a strong limit which is well-defined up to isomorphism (see [8]).

For any finite  $L_n$ -graphs  $A \subseteq B, C$  we denote by  $B \otimes_A C$  the *free amalgam* of B and C over A, i.e., the graph on  $B \cup C$  containing no edges between elements of  $B \setminus A$  and  $C \setminus A$ .

LEMMA 2.7. If  $A \leq B$ , C are in  $\mathcal{K}_n$ , then  $D := B \otimes_A C \in \mathcal{K}_n$  and  $B, C \leq D$ .

**PROOF.** Clearly,  $B, C \leq D$ . To see that  $D \in \mathcal{K}_n$ , note that if  $B \in \mathcal{K}_n$  and B' is a minimal strong extension of B, then also  $B' \in \mathcal{K}_n$ . This is clear for strong extensions of type 1. For strong extensions of type 2 suppose that we have added a dense flag  $(y_1, \ldots, y_m)$  to B connecting  $y_0, y_{m+1} \in B$ . Conditions 1, 2, and 3 are clear. To see that Condition 4 holds, let  $a \in V_i, b, b' \in R(a)$ , and suppose that we have a path  $\gamma = (b, \dots, b') \subset B'$  of length k in  $V_{i-m} \cup \dots \cup V_i$  containing a new vertex  $y_i$  as in the assumptions of Condition 4. We may assume  $\gamma \not\subset B$ . Since the new vertices have exactly two neighbors, note that *a* cannot be one of the new vertices. Furthermore, either the entire flag is contained in  $\gamma$  and can be replaced by a flag of the same type in B or an initial or end segment of  $\gamma$  is contained in the flag. By symmetry assume that an end segment of  $\gamma$  is contained in the new flag, so  $b' = y_i$ for some  $i \in \{1, \ldots, m\}$ . In this case the entire flag  $(y_1, \ldots, y_m)$  is in  $R_{\leq}(a)$ , and so are  $y_0, y_{m+1}$ . By assumption we may replace the flag  $(y_1, \ldots, y_m)$  by a flag in  $B \cap R(a)$  connecting  $y_0$  and  $y_{m+1}$ . By Condition 4 we find a path from b to  $y_0$  inside  $R(a) \cap B$ . This can be extended to the required path from b to  $b' = y_i$  inside R(a). Conditions 5 and 6 are similar. -

This shows that the class  $(\mathcal{K}_n, \leq)$  has a Fraissé limit  $M_n$ .

**PROPOSITION 2.8.** The Hrushovski limit  $M_n$  is a model of  $T_n$ .

PROOF. This is clear for the case n = 1. By construction,  $M_n$  satisfies  $(\Sigma 3)_n$  and  $(\Sigma 4)_n$ . By Remark 2.6 and induction,  $V_i \cup \ldots \cup V_{i+j}$  is a model of  $T_j$  for any i + j < n. In particular,  $M_n$  satisfies  $(\Sigma 1)_n$ . It is left to show that  $M_n$  satisfies  $(\Sigma 2)_n$ . So let  $a \in V_n$ . We have to show that R(a) is a pseudospace of dimension n - 1. Clearly,  $(\Sigma 3)_{n-1}$  continues to hold. By induction applied to  $V_1 \cup \ldots \cup V_n$ ,  $R(a) \cap (V_1 \cup \ldots \cup V_{n-1})$  is a pseudospace of dimension n - 2, so  $(\Sigma 1)_{n-1}(b)$  holds. Similarly  $(\Sigma 2)_{n-1}(b)$  and  $(\Sigma 4)_{n-1}(b)$  hold. Note that  $(\Sigma 4)_{n-1}(a)$  holds for R(a) by condition 6. It is left to show that  $(\Sigma 1)_{n-1}(a)$  and  $(\Sigma 2)_{n-1}(a)$  hold for R(a), that

is we have to show that  $R(a) \cap (V_0 \cup \ldots \cup V_{n-2})$  and  $R(a) \cap R(c)$  for  $c \in V_0$  are pseudospaces of dimension n-2.

 $(\Sigma 1)_{n-1}(a)$ : To see that  $R(a) \cap (V_0 \cup \ldots V_{n-2})$  is a pseudospace of dimension n-2note that by induction  $(\Sigma 1)_{n-2}(b), (\Sigma 2)_{n-2}(b), (\Sigma 3)_{n-2}(a)$  and (b) and  $(\Sigma 4)_{n-2}(a)$ and (b) hold and it is left to show that  $R(a) \cap (V_0 \cup \ldots \cup V_{n-2})$  satisfies  $(\Sigma 1)_{n-2}(a)$ and  $(\Sigma 2)_{n-2}(a)$ . That is we have to show that  $R(a) \cap (V_0 \cup \ldots \cup V_{n-3})$  and  $R(a) \cap (V_1 \cup \ldots V_{n-2}) \cap R(c)$  for  $c \in V_0$  are pseudospaces of dimension n-3. In this way we reduce to show that  $R(a) \cap (V_0 \cup V_1)$  and  $R(a) \cap (V_1 \cup V_2) \cap R(c)$  for  $c \in V_0$ are pseudospaces of dimension 1, and this is obvious.

 $(\Sigma 2)_{n-1}(a)$ : To see that  $R(a) \cap R(c)$  for  $c \in V_0$  is a pseudospace of dimension n-2 note that  $(\Sigma 2)_{n-2}(a)$  and (b),  $(\Sigma 3)_{n-2}(a)$  and (b), and  $(\Sigma 4)_{n-2}(a)$  and (b) hold. It is therefore left to show that  $R(a) \cap (V_1 \cup \ldots V_{n-2}) \cap R(c)$  and  $R(a) \cap (V_2 \cup \ldots V_{n-1}) \cap R(c)$  are pseudospaces of dimension n-3. The first part was shown above in the proof for  $(\Sigma 1)_{n-1}(a)$  for R(a), the second follows by reduction to  $R(a) \cap (V_2 \cup V_3) \cap R(c)$ , which is a pseudospace of dimension 1.  $\dashv$ 

Note that  $(\Sigma 4)$  implies the following:

LEMMA 2.9. Let M be a model of  $T_n$ ,  $a, c \in M$  with  $a \notin R(c)$ . Let  $\gamma_1 = (x_0 = a, \ldots, x_s), \gamma_2 = (y_0 = a, \ldots, y_t)$  be paths with  $x_s, y_t \in R(c)$  and let i, j be minimal with  $x_i, y_j \in R(c)$ . Suppose that  $\gamma_1, \gamma_2$  do not contain any element of higher level than the level of c. Then  $(x_i, y_j)$  is a flag.

**PROOF.** Since  $x_i, y_j \in R_{<}(c)$  the claim follows directly from ( $\Sigma 4$ ) as otherwise the path  $(x_i, x_{i-1}, \ldots, x_0 = y_0, y_1, \ldots, y_j)$  would have to contain an element of R(c) in its interior, which it doesn't.

COROLLARY 2.10. Let M be a model of  $T_n$ ,  $a, c \in M$  with  $a \notin R(c)$ . There is a flag  $C \subset R_{\leq}(c)$  such that for any  $b \in R_{\leq}(c)$  and any path from a to b not containing an element of higher level than the level of c enters  $R_{\leq}(c)$  via an element of C.

From now on we work inside models of  $T_n$  unless specified otherwise.

Let us say that a path  $\gamma = (a = x_0, ..., x_m = b)$  changes direction in  $x_i$  if  $x_i \in V_j$ and either  $x_{i-1}, x_{i+1} \in V_{j-1}$  or  $x_{i-1}, x_{i+1} \in V_{j-1}$  for some *j*. Clearly, a path which never changes direction is a dense flag.

DEFINITION 2.11. Let  $\gamma = (y_0, \dots, y_1, \dots, y_{k+1})$  be a path changing direction exactly in  $y_1, \dots, y_k$ . We say that  $\gamma$  is reduced if for all  $i = 0, \dots, k-1$  we have

$$y_i \lor y_{i+2} = y_{i+1}$$
 or  $y_i \land y_{i+2} = y_{i+1}$ .

Clearly, any part of a reduced path is again reduced.

LEMMA 2.12. If  $\gamma = (a, ..., b) \subseteq V_j \cup ... \cup V_{j+m}$  has length s, there is a reduced path from a to b inside  $V_j \cup ... \cup V_{j+m}$  of length at most s.

PROOF. Let  $\gamma = (a = y_0, \dots, y_1, \dots, y_{k+1} = b) \subseteq V_j \cup \dots \cup V_{j+m}$  change direction exactly in  $y_1, \dots, y_k$ . We can reduce  $\gamma$  by putting  $z_0 = y_0$  and replacing for  $i = 1, \dots, k$  inductively  $y_i$  by  $z_i = z_{i-1} \vee y_{i+1}$  (or  $z_i = z_{i-1} \wedge y_{i+1}$ , respectively) filling the path with flags between  $z_i$  and  $z_{i+1}$ . The new path  $\gamma'$  has length at most s, is still in  $V_j \cup \dots \cup V_{j+m}$ , and changes direction *at most* in  $z_1, \dots, z_k$ . Note that if  $\gamma$  was not reduced, then the length must go down. We repeat the procedure with those vertices  $z_i$  in which  $\gamma'$  changes direction. This process stops after finitely many reductions with a reduced path of length at most s. We need the following lemmas:

LEMMA 2.13. If  $a \in R(b)$ , then any reduced path from a to b is a flag. If  $a, b \in R_{\leq}(c)$  are connected by a reduced path  $\gamma$ , then  $\gamma \subseteq R_{\leq}(c)$ .

PROOF. We prove both claims by induction on the length k of a reduced path from a to b. For k = 1 both statements are clear.

Now suppose that both statements are proved for paths of length at most k - 1. Let  $\gamma = (a = x_0, \ldots, x_k = b)$  be a reduced path with  $a \in R_{\leq}(b)$ . By induction assumption we have  $(x_1, \ldots, x_k = b) \subset R_{\geq}(a)$ , and hence  $x_{k-1} \in R(a) \cap R(b)$ . Hence again by induction  $(a = x_0, x_1, \ldots, x_{k-1})$  is a flag. Since  $\gamma$  is reduced, it cannot change direction in  $x_{k-1}$ . It follows that  $\gamma$  is a flag.

For the second assertion let  $\gamma = (a = x_0, \dots, x_k = b)$  be a reduced path with  $a, b \in R_{\leq}(c)$ . If  $x_i \in R_{\leq}(c)$  for some  $1 \leq i \leq k - 1$ , then  $\gamma \subset R_{\leq}(c)$  by induction assumption. So suppose there is no  $1 \leq i \leq k - 1$  with  $x_i \in R_{\leq}(c)$ . Then by Remark 2.3  $\gamma$  contains some element whose level is higher than the level of c. Let  $\gamma' = (a, y_1, \dots, c, \dots, b)$  be a path consisting of flags  $(a, \dots, c)$  and  $(c, \dots, b)$ . So  $\gamma \cap \gamma' = \{a, b\}$ . Let  $1 \leq i \leq k - 1$  be such that the level of  $x_i$  is maximal. Let  $j_0, j_1$  be minimal (maximal, respectively) such that  $x_{j_0}, x_{j_1} \in R(x_i)$  and consider

$$\gamma'' = (x_{j_0}, \dots, x_1, x_0 = a, y_1, \dots, c \dots, b = x_k, x_{k-1}, \dots, x_{j_1}).$$

Again by Remark 2.3, the path  $\gamma''$  contains in its interior some vertex  $y \in R_{<}(x_i)$ . By choice of  $j_0, j_1$  we must have  $y \in (a, \ldots, c \ldots, b)$ . If y lies between a and c, then we have  $a \in R_{<}(x_i)$  and hence  $j_0 = 0$  and  $(a = x_0, \ldots, x_i)$  is a flag by induction. But then a is not an interior point of  $\gamma''$ . Similarly if y lies between c and b, and this finishes the proof.

Note that if  $a \in R_{\leq}(c)$ ,  $b \in R_{>}(c)$ , then a reduced path from a to b is a flag, but need not contain c.

LEMMA 2.14. Let  $\gamma = (x_0, \dots, x_s)$  be a reduced path and let  $(y_0 = x_0, \dots, y_k)$  be a flag. Let  $m \leq s$  be minimal such that  $y_k \notin R(x_m)$ . Then for some  $i \geq m - 1$  the path  $(y_k, \dots, x_i, x_{i+1}, \dots, x_j, \dots, x_s)$  is reduced and changes direction in  $x_i$ .

**PROOF.** First suppose  $(y_k, \ldots, x_0, \ldots, x_s)$  does not change direction in  $x_0$  and, say,  $y_k \in R_{\leq}(x_0)$ . Let  $j \ge m$  be minimal such that  $\gamma$  changes direction in  $x_j$ . Then  $z = y_k \lor x_j \in R(x_{m-1})$ . Hence we may assume that  $z \in \{x_s : s = m - 1, \ldots, j\}$  and the claim follows.

Now assume that  $(y_k, \ldots, x_0, \ldots, x_s)$  changes direction in  $x_0$  with  $y_k \in R_>(x_0)$ . Then  $x_{m-1} = y_k \wedge x_m \in R_<(x_m)$  and  $(y_k, \ldots, x_{m-1}, \ldots, x_s)$  is reduced and changes direction in  $z = x_{m-1}$ .

As in the theory of buildings we can show here that residues are *gated*, i.e., the following holds in any model M of  $T_n$ :

LEMMA 2.15. Let M be a model of  $T_n$ . For all  $a, c \in M$  with  $a \notin R(c)$  there is a flag  $f = (b_1, \ldots, b_k) \subset R(c)$  such that any reduced path from a to c enters R(c) via some element of f.

PROOF. The proof is by induction on the rank of the pseudospace. The claim is clear if the rank is 1. Assume now that the claim is proved for rank less than n. For rank n, it suffices to prove the following:

If  $\gamma_1 = (x_0 = a, \dots, x_s = c), \gamma_2 = (y_0 = a, \dots, y_t = c)$  are reduced paths and *i*, *j* are minimal with  $x_i, y_j \in R(c)$ , then  $x_i = y_j$  or  $(x_i, y_j)$  is a flag.

Note that by the definition of a reduced path and by Lemma 2.13,  $x_i$ ,  $y_j$  are exactly the last vertices where  $\gamma_1$ ,  $\gamma_2$  change direction.

We do induction on min{i, j}, by symmetry we may assume  $i \le j$ . If i = 1, then  $a, c \in R_{>}(x_1)$  (up to duality) and  $\gamma_1, \gamma_2 \in R_{\geq}(x_1)$  by Lemma 2.13. In particular,  $y_j \in R \ge (x_1)$ .

Now let i > 1. If  $\gamma = (x_1, a, y_1, \dots, y_t = c)$  is reduced we apply the induction hypothesis to the paths starting at  $x_1$ . So we may assume that  $(x_1, a, y_1, \dots, y_t = c)$  is not reduced.

First assume that  $\gamma$  changes direction in a. Let m be minimal such that  $\gamma_2$  changes direction in  $y_m$ , so  $m \leq j$ . Since  $\gamma$  is not reduced, we have  $x_1 \in R(y_m)$ . So we may assume that  $y_1 = x_1$  and apply the induction hypothesis to the paths starting from  $x_1$ .

If  $\gamma$  does not change direction in a, let us assume for fixing notation that  $x_1 \in R_{<}(a)$ . Let m be minimal with  $x_1 \notin R(y_m)$  and such that  $\gamma_2$  changes direction in  $y_m$  and put  $z = x_1 \lor y_m$ . Then  $(x_1, \ldots, z, \ldots, y_m, \ldots, c)$  is a reduced path. If  $j \ge m$ , we may apply the induction hypothesis to  $x_1$ . If j < m, then  $x_1 \in R_{<}(y_j)$ . If  $c \in R_{>}(y_j)$ , then  $x_1 \in R(c)$ , contradicting i > 1. Hence  $x_1, c \in R_{<}(y_j)$  and we finish by Lemma 2.13 as  $(x_1, x_2, \ldots, c) \subseteq R_{\le}(y_j)$ .

As a first step toward showing that  $T_n$  is a complete theory, we show that any finite subset of a model of  $T_n$  is contained in a *nice* subset.

We call two reduced paths  $\gamma_1$ ,  $\gamma_2$  from *a* to *b* equivalent,  $\gamma_1 \sim \gamma_2$  if they have the same set of vertices in which they change direction. Note that if  $\gamma_1 \sim \gamma_2$ , then  $\gamma_1$  is reduced if and only  $\gamma_2$  is.

DEFINITION 2.16. Following [1] we call a subset A of a model M of  $T_n$  nice (in M) if A is in  $\mathcal{K}_n$  and the following holds:

- 1. if  $a, b \in A$  are connected (in M) by a reduced path  $\gamma$  of length k contained in  $V_{i-m} \cup \ldots \cup V_i$  in M there is an equivalent path  $\gamma'$  from a to b inside A.
- 2. If  $a, a' \in A$ , then  $a \lor a' \in A$  if  $a \lor a'$  exists in M.
- 3. If  $a, a' \in A$ , then  $a \wedge a' \in A$  if  $a \wedge a'$  exists in M.

Remark 2.17.

- 1. Note that (the proof of) Lemma 2.12 implies that if A is nice in a model M and  $a, b \in A$  are connected by a path in M, then they are connected by a reduced path in A. Therefore conditions 4, 5, and 6 for graphs in  $\mathcal{K}_n$  automatically hold for a nice set A.
- 2. If A is nice,  $a, b \in A$ , and  $\gamma = (a = x_0, ..., x_m = b)$  is a reduced path changing direction in  $y_1, ..., y_k$ , then  $y_1, ..., y_k \in A$ .
- 3. If A is nice, then also A ∩ (V<sub>j</sub> ∪ ... ∪ V<sub>j+m</sub>) is nice in the sense of the pseudospace (V<sub>j</sub> ∪ ... ∪ V<sub>j+m</sub>) for all j + m ≤ n. In particular, since E<sub>i</sub>-paths between elements are unique, any E<sub>i</sub>-path between elements of A lies entirely in A and if a, b ∈ A are contained in a dense flag in M, they are contained in a dense flag of A. Furthermore, if A is nice and b ∈ A, then by Lemma 2.13 R≥(b) ∩ A, R≤(b) ∩ A and hence R(b) ∩ A are nice. Note also that R>(b) ∩ A, R<(b) ∩ A are nice in the sense of the pseudospace R<(b), R>(b), respectively.

LEMMA 2.18. Let M be a model of  $T_n$  and let A, B be finite subsets of M with  $A \leq B$ . If  $a, b \in A$  are connected by a path  $\gamma \subset B$ , then there is an equivalent path  $\gamma' \subset A$ . In particular, if B is nice, then so is A.

**PROOF.** Write  $B = \bigcup_{i < k} B_i$  with  $B_0 = A$  and such that  $B_{i-1} \leq B_i$  is a minimal strong extension for i = 1, ..., k. Let  $\gamma = (a = x_0, ..., x_m = b)$  and let j be minimal with  $\gamma \subseteq B_j$ . Then  $B_j$  must be a strong extension of  $B_{j-1}$  of type 2, and we may replace the new flag of  $B_j$  by a flag in  $B_{j-1}$  to obtain an equivalent path in  $B_{j-1}$ . Continuing in this way we eventually find an equivalent path in A.

We will show that any finite set is contained in a nice strong finite set. To simplify inductive proofs we define a *pointed pseudospace (of dimension n)* as a free pseudospace V of dimension n together with a new vertex  $x \in V_{n+1}$  or  $x \in V_{-1}$ incident exactly with the elements of  $V_n$  (or of  $V_0$ ). Thus, if V is a free pseudospace of dimension n and  $b \in V_i$ , then  $R_{\leq}(b)$  is a pointed pseudospace of dimension i-1and  $R_{\geq}(b)$  is a pointed pseudospace of dimension n-i-1. We say that a subset of a pointed pseudospace is nice if it satisfies the conditions in Definition 2.16.

LEMMA 2.19. Let M be a pseudospace or a pointed pseudospace. If  $A \subset M$  is finite and nice in M and a is arbitrary, then there is a nice finite set B containing  $A \cup \{a\}$ such that  $A \leq B$ .

PROOF. The proof is by induction on the dimension n of the (pointed) pseudospace. The claim is clear for n = 0, so assume it has been proved for all (pointed) pseudospaces of dimension less than n. If M is a pointed pseudospace, let  $x \in M$  be the additional point. Of course we may assume  $a \notin A$ . We may also assume that there is some reduced path  $\gamma = (a = x_0, \dots b)$  for some  $b \in A$  and  $\gamma \cap A = \{b\}$  as otherwise  $A \cup \{a\}$  is nice. It therefore suffices to prove the claim for the case where  $a \in V_i$  has a neighbor  $b \in A$  of type  $V_{i+1}$  (the other case being dual to this one). If b = x and any reduced path from a to an element of A passes through x, then  $A \cup \{a\}$  is nice. Otherwise we may assume that  $b \neq x$  so that we may apply the induction hypothesis to  $R_{<}(b)$ .

By Remark 2.17 we know that  $R_{\leq}(b) \cap A$  is nice. By induction hypothesis we find a nice finite set  $B \subseteq R_{\leq}(b)$  containing *a* such that  $(R_{\leq}(b) \cap A) \leq B$ . We claim that  $A \leq A \cup B$  and  $A \cup B$  is nice. To see this we write  $B = \bigcup_{j < r} B_j$  as a union of minimal strong extensions over  $B_0 = A \cap R_{\leq}(b)$  and show inductively that  $A \cup B_{j-1} \leq A \cup B_j$  and  $A \cup B_j$  is nice. Note that by Lemma 2.18 each  $B_j$  is nice. **Case I** If  $B_j = B_{j-1} \cup \{c\}$  is a strong extension of  $B_{j-1}$  of type 1 by some  $c \in R_{<}(b)$ , then since  $B_j$  is nice and  $c \in R_{<}(b)$ , we must have  $c \in R_{<}(d)$  for a unique neighbor d of c with  $d \in B_{j-1}$ . Since  $c \in R_{<}(b)$  (and  $B_j$  is nice in the sense of  $R_{<}(b)$ ), we have  $R_{<}(c) \cap (A \cup B_j) = R_{<}(c) \cap B_j = \emptyset$ .

If *c* had another neighbor  $d' \in A \setminus B_{j-1}$ , then *d*, *d'* could not have the same level since otherwise by induction  $c \in A \cup B_{j-1}$ . So (d, c, d') must be a flag, contradicting our assumption. So *c* has no other neighbor in  $A \cup B_{j-1}$ , and hence  $A \cup B_{j-1} \leq A \cup B_j$ .

To see that  $A \cup B_j$  is nice, note that for  $x \in A \setminus B_j$  we must have  $c \lor x = d \lor x$  since otherwise the path  $(d, c, \ldots c \lor x, \ldots, x)$  is reduced and changes direction in c, yielding  $c \in B_{j-1}$  by induction assumption. Hence Condition 2. holds.

To see that Condition 3 holds, note that for  $x \in A$  we have  $d \wedge x \in B_{j-1}$ . If  $\emptyset \neq c \wedge x \neq d \wedge x$ , then the path  $(c, \ldots, c \wedge x, \ldots, d \wedge x)$  is reduced and changes direction in  $c \wedge x$ . Since  $B_j$  is nice we have  $c \wedge x \in B_j$ . So by assumption we must have  $c = c \wedge x$ , i.e.,  $c \in R_{\leq}(x)$ .

Now suppose that *c* is connected to  $x \in A \cup B_{j-1}$  by a reduced path  $\gamma = (c = x_0, \ldots, x_m = x)$ . Consider the extension  $\gamma' = (d, c = x_0, \ldots, x_m)$ .

If  $\gamma'$  did not change direction in c, then  $\gamma'$  could not be reduced since otherwise  $x_i \in R_{<}(c) \cap B_{j-1}$  for some 0 < i < m, a contradiction. So let i, j be minimal such that  $\gamma'$  changes direction in  $x_i$  and  $x_j$ . It is easy to see that we can reduce  $\gamma'$  to a path  $(d, \ldots, y = d \land x_j, \ldots, x_j, x_{j+1}, \ldots, x_m = x)$  changing direction in  $y \in B_{j-1}$ . But then  $x_i = c \land y \in R_{<}(c) \cap B_j$  by Condition 3, again a contradiction.

Hence  $\gamma'$  changes direction in c. Then  $\gamma'$  cannot be reduced since otherwise by niceness of  $A \cup B_{j-1}$  we have  $c \in B_{j-1}$ . We can reduce  $\gamma'$  to a path  $(d, \ldots, x_i, \ldots, x_m = x)$  where i is minimal such that  $\gamma$  changes direction in  $x_i$ . Then the path  $(c, d, \ldots, x_i, \ldots, x_m = x)$  is equivalent to  $\gamma$  and by induction assumption we may replace the path  $(d, \ldots, x_i, \ldots, x_m = x)$  by an equivalent one inside  $A \cup B_{j-1}$ .

**Case II** If  $B_j$  arises from  $B_{j-1}$  by a strong extension of type 2, then there are  $b_1, b_2 \in B_{j-1}$  which are connected by a flag inside  $B_{j-1}$  and  $B_j = B_{j-1} \cup \{x_1, \ldots, x_k\}$  where  $(b_1, x_1, \ldots, x_k, b_2)$  is a flag,  $b_1 \in R_{\leq}(b_2)$ , say. If for some  $i \in \{1, \ldots, k\}$  the vertex  $x_i$  has a neighbor  $d \in A \setminus B_{j-1}$ , then since  $d \notin R_{\leq}(b)$ , we must have  $x_i = d \land b_2 \in B_{j-1}$  by induction assumption. This is impossible since  $B_j$  is a minimal strong extension of  $B_{j-1}$ . Hence  $A \cup B_j$  is a minimal strong extension of  $A \setminus B_{j-1}$ .

To see that  $A \cup B_j$  is nice, let  $c, d \in A \cup B_j$ . If  $c, d \in A \cup B_{j-1}$  or  $c, d \in B_j$ there is nothing to show. So assume  $c \in A \setminus B_j$  and  $d = x_i \in B_j \setminus B_{j-1}$  for some  $1 \le i \le k$ . For Condition 2, we claim that  $c \lor d = c \lor b_2$ . Suppose to the contrary that  $c \lor d \ne c \lor b_2$  and consider the path  $(b_2, \ldots, d, \ldots, d \lor c, \ldots, c)$ . If this path is reduced, then since it changes direction in d, by induction assumption we have  $d \in A \cup B_{j-1}$  and hence  $d \in B_{j-1}$ , a contradiction. Hence we find a reduced path  $(b_2, \ldots, y, \ldots, d \lor c, \ldots, c)$  with  $y = b_2 \land (d \lor c) \in B_{j-1}$ . Since  $B_j$  is a minimal strong extension of  $B_{j-1}$ , we must have  $y = b_2$  and hence  $b_2 \in R_< (d \lor c)$ . Therefore  $b_2 \lor c = d \lor d$ , contradicting the assumption.

Similarly, for Condition 3, we claim that  $c \wedge d = c \wedge b_1$ , the argument being exactly dual to the one for Condition 2.

For Condition 1, assume now that c, d are connected by a reduced path  $\gamma = (d = y_0, \dots, y_\ell = c)$  and extend  $\gamma$  by the flag  $(b_2, \dots, x_i = d = y_0)$  to a path  $\gamma'$ . Then since  $c \notin R_<(b_2) \subseteq R_<(b)$ , there is a minimal t such that  $y_t \notin R_<(b_2)$ .

By Lemma 2.14 there is a reduced path from  $b_2$  to c changing direction in  $y_s$  for some  $s \ge t - 1$ . Then  $y_s \in A \cup B_{j-1}$ , and so  $y_s \in R_{<}(b)$  implies  $y_s \in B_{j-1}$  by induction assumption. By induction assumption and niceness of  $B_j$  we find inside  $A \cup B_j$  a path from d to c equivalent to  $\gamma$ .

REMARK 2.20. The construction shows that for a path  $\gamma = (x_0, \ldots, x_k)$  changing direction in  $y_1, \ldots, y_k$  there is a nice finite set  $A \supset \gamma$  with  $A \subseteq \bigcup_{i=1,\ldots,k} R_{\delta}(y_i)$  where  $\delta \in \{\leq, \geq\}$  depending on the change of direction in  $y_i$ .

COROLLARY 2.21.

- 1. If A is a finite subset of a model M of  $T_n$ , there is a nice finite set B in M containing A.
- 2. For any two vertices a and b, up to equivalence there are only finitely many reduced paths from a to b.
- 3. Reduced paths  $\gamma = (a = x_0, \dots, x_k = b), \gamma' = (a = y_0, \dots, y_k = b)$  from a to *b* are equivalent if (and only if)  $x_i, y_i$  have the same level for all  $i = 0, \dots, k$ .

**PROOF.** Part 1. is clear and Part 2. follows directly from 1. The proof of the nontrivial direction for 3. is by induction on the number of vertices in which  $\gamma$  (and  $\gamma'$ ) change direction. If they change direction exactly once, the claim is clear. By Lemma 2.15 the last change of direction of  $\gamma$  and  $\gamma'$  is in the same element of R(b). Hence the claim follows from the induction hypothesis.  $\dashv$ 

We also note the following corollary:

COROLLARY 2.22. If  $A_0 \subseteq M_n$  is a nice finite set, then we can write  $M_n = \bigcup_{i < \omega} A_i$  with  $A_i \leq A_{i+1}$ .

We say that a model M of  $T_n$  is  $\mathcal{K}_n$ -saturated if for all nice finite sets  $A \subset M$  and strong extensions C of A with  $C \in \mathcal{K}_n$  there is a nice embedding of C into M fixing A elementwise. The following lemma shows that  $M_n$  is  $\mathcal{K}_n$ -saturated:

LEMMA 2.23. Write  $M_n = \bigcup_{i < \omega} A_i$  with  $A_i \in \mathcal{K}_n$  and  $A_i \leq A_{i+1}$ . Then each  $A_i$  is nice in  $M_n$ .

**PROOF.** Fix *i* and let  $a, b \in A_i$ . If  $\gamma = (a = x_0, ..., x_k = b)$  is a path, then  $\gamma \subseteq A_{i+s}$  for some s > i. By Lemma 2.18 there is an equivalent path inside  $A_i$ , so  $A_i$  is nice.

LEMMA 2.24. An  $L_n$ -structure M is an  $\omega$ -saturated model of  $T_n$  if and only if M is  $\mathcal{K}_n$ -saturated.

PROOF. Let M be an  $\omega$ -saturated model of  $T_n$ . To show that M is  $\mathcal{K}_n$ -saturated, let  $A \subset M$  be a nice finite set and  $A \leq B \in \mathcal{K}_n$ . By induction we may assume that B is a minimal strong extension of A. By  $\omega$ -saturation it is easy to see that B can be nicely embedded over A into M. Conversely, assume that M is  $\mathcal{K}_n$ -saturated. Since any finite subset A of M is contained in a nice finite set  $B \subseteq M$  we see that M is back-and-forth equivalent to  $M_n$  and so is a model of  $T_n$ . Choose an  $\omega$ -saturated  $M' \equiv M$ . Then by the above M' is  $\mathcal{K}_n$ -saturated. So M' and M are also back-and-forth equivalent, which implies that M is  $\omega$ -saturated.

COROLLARY 2.25. The theory  $T_n$  is complete.

PROOF. Let M be a model of  $T_n$ . In order to show that M is elementarily equivalent to  $M_n$  choose an  $\omega$ -saturated  $M' \equiv M$ . By Lemma 2.24, M' is  $\mathcal{K}_n$ -saturated. Now M' and  $M_n$  are back-and-forth equivalent and therefore elementarily equivalent.

REMARK 2.26. Note that this implies also that for nice finite sets A the quantifier free type of A already determines the complete type of A.

We will see in Section 4 that  $T_n$  is the theory of the building of type  $A_{\infty,n+1}$  with infinite valencies.

COROLLARY 2.27. The algebraic closure  $\operatorname{acl}(A)$  of a finite set A is the intersection of all nice sets containing A.

PROOF. Clearly, the intersection of all nice sets is contained in acl(A) by Lemma 2.19. For the converse, assume *B* is a nice set containing *A* and  $x \notin B$ . Let  $D \supset B \cup \{x\}$  be a nice set. Since  $M_n$  is  $\mathcal{K}_n$ -saturated and the free amalgam of any finite number of copies of *D* over *B* is again in  $\mathcal{K}_n$ , we can find infinitely many copies of *D* over *B* in  $\mathcal{K}_n$ .

Using the fact that  $M_n$  is  $\omega$ -saturated we can now give an explicit description of the algebraic closure:

**PROPOSITION 2.28.** A vertex  $c \neq a, b$  is in acl(ab) if and only if there is a reduced path from a to b changing direction in c. Hence  $acl(ab) = \{a, b\}$  if and only if a, b is a flag or a and b are not connected.

**PROOF.** By Corollary 2.27 any *c* for which there is a reduced path from *a* to *b* that changes direction in *c* is in acl(ab).

Now let  $c \neq a, b$  and suppose that there is no reduced path between a and b changing direction in c. If there is no (reduced) path from a to b, the set  $\{a, b\}$  is already nice. So let  $\gamma = (a = x_0, \ldots, x_s = b)$  be a reduced path. We may assume  $c \notin \gamma$ . We construct a nice set containing a, b but not c. We may also assume that  $|\gamma \cap R(c)|$  is minimal. For rank n = 1 this is easy. So we may assume that we have proved the claim for all pseudospaces of rank less than n. Also, if R(c) does not intersect  $\gamma$ , then Remark 2.20 shows that we can find a nice set containing  $\gamma$  but not c. So let k be maximal such that for some  $i \leq s$  we have  $x_i, x_{i+k} \in R(c)$ . Note that  $\gamma$  must change direction in  $x_i$  and  $x_{i+k}$  as otherwise we could replace  $x_i, x_{i+k}$  by other elements to minimize  $|\gamma \cap R(c)|$ . Suppose that  $c \in V_m$ .

First suppose that k > 0 and both  $x_i, x_{i+k} \in R_{<}(c)$  (or by symmetry both in  $R_{>}(c)$ ) so that  $R_{<}(c) \cap \gamma = (x_i, \ldots, x_{i+k})$  by Lemma 2.13. Note that this implies that  $x_i, x_{i+k}$  have level at most m - 2. Then  $\gamma$  must change direction at some place between  $x_i$  and  $x_{i+k}$  as this path cannot be a flag. If  $x_i \lor x_{i+k} = c$  then we claim that we can obtain a reduced path from a to b changing direction in c: namely if a is the last place where  $\gamma$  changes direction before  $x_i$ , then  $z_1 = a \land c \in R_{<}(a)$ . Similarly if b is the first place where  $\gamma$  changes direction after  $x_{i+k}$ , then  $z_2 = b \land c \in R_{<}(b)$ . Replacing in  $\gamma$  the path  $(a, \ldots, b)$  by  $(a, \ldots, z_1, \ldots, c, \ldots, z_2, \ldots, b)$ , we obtain a path, which is easily seen to be reduced, changing direction in c.

Hence  $z = x_i \lor x_{i+k} \in R_{<}(c)$ . We may assume that the path from  $x_i$  to  $x_{i+k}$  consists of flags  $(x_i, \ldots, z)$  and  $(z, \ldots, x_{i+k})$ . By induction assumption and the fact that  $z \in R_{<}(c)$  we find an absolutely nice set  $D_1 \subseteq R_{<}(c)$  containing  $(x_i, \ldots, x_{i+k})$  (and not containing c). So  $D_1 \subseteq V_0 \cup \ldots \cup V_{m-1}$ . Since  $x_{i-1}, x_{i+k+1} \in V_0 \cup \ldots \cup V_{m-1}$  we can extend  $D_1$  to a nice set  $D_2 \subset V_0 \cup \ldots \cup V_{m-1}$  containing  $x_{i-1}, x_{i+k+1}$ . We extend  $D_2$  to a nice set D containing  $(x_0, \ldots, x_{i-2})$  and  $(x_{i+k+2}, \ldots, x_s)$  as described in Lemma 2.19. Since  $c \notin R(x_j), j = 0, \ldots, i-2, i+k+2, \ldots, s$  we have  $c \notin D$ .

Now suppose that k = 0 and  $x_i \in R_{<}(c)$ . Let a', b' be the last place before (and the first place after)  $x_i$  where  $\gamma$  changes direction, so  $a', b' \in R_{>}(x_i)$ .

**Claim:** There is no reduced path from a' to b' changing direction in c. Suppose otherwise and let  $\gamma' = (a' = y_0, \dots, y_s = b')$ . Then c cannot be the first or last place where  $\gamma'$  changes direction since k = 0. If  $y_1 \in R_{<}(a'), y_{s-1} \in R_{<}(b')$ ,

then replacing  $(a', \ldots, x_i, \ldots, b')$  by  $\gamma'$  in  $\gamma$  yields a reduced path changing direction

in c, contradicting the assumption. If  $y_1 \in R_>(a')$  let  $d_0 = a', d_1, d_2, \ldots \in \gamma'$  an enumeration of the places where  $\gamma'$  changes direction and let  $a_1$  be the last place before a' where  $\gamma$  changes direction. Put  $z_1 = a_1 \lor d_2 \in R_{\leq}(d_1)$ , so  $z_1 \land d_3 = d_2$ . We can now reduce the path starting from  $z_1$  from right to left as in Lemma 2.12. If necessary we can do the same on the right end of  $\gamma'$  and thus obtain a reduced path from a to b changing direction in c, a contradiction proving the claim.

By the claim we can apply the induction assumption to  $R_>(x_i)$ . So let  $D \subset R_>(x_i)$ be a nice set in the sense of  $R_>(x_i)$  containing  $x_{i-1}, x_{i+1}$ , but not containing c. Then  $D_1 = D \cup \{x_i\}$  is nice in the sense of the  $V_0 \cup \ldots \cup V_n$  and we can extend  $D_1$  to a nice set  $D_2$  as in Lemma 2.19. Since  $c \notin R(x_j), j = 0, \ldots, i - 2, i + 2, \ldots, s$  we have  $c \notin D_2$ .

Finally suppose that  $x_i \in R_{\leq}(c)$ ,  $x_{i+k} \in R_{>}(c)$  (the other case is similar), so that  $(x_i, \ldots, x_{i+k})$  is a flag and hence a nice set. If  $x_{i+k} \in V_k$  for k < n, then we can extend the nice set  $\{x_i, \ldots, x_{i+k}\}$  by induction assumption to a nice set containing  $x_{i-1}, x_{i+k+1}$ , but not containing c and we continue as in Lemma 2.19. If  $x_{i+k} \in V_n$  and every reduced path from a to b passes through  $x_{i+k}$ , then we can use the induction assumption to construct a nice set in  $V_0 \cup \ldots \cup V_{n-1}$  containing  $(x_{i-1}, \ldots, x_{i+k-1})$  and not c. This can be extended to a nice set in  $V_0 \cup \ldots \cup V_n$  by adding the necessary elements in  $V_n$ . We can also construct a nice set containing  $(x_{i+k+1}, \ldots, b)$ . The union of these two sets together with  $x_{i+k}$  will be a nice set since every path from one of these sets to the other set has to pass through  $x_{i+k}$ .

Now suppose that there is a path from  $x_{i+k-1}$  to  $x_{i+k+1}$  not passing through  $x_{i+k}$ . Since  $x_{i+k-1}, x_{i+k+1} \in R_{<}(x_{i+k})$ , by ( $\Sigma$ 4) there is such a path  $\gamma' = (x_{i+k-1} = y_0, \ldots, y_t = x_{i+k+1})$  in  $R_{<}(x_{i+k})$ . We can now apply the induction assumption to find a nice set in  $V_0 \cup \ldots \cup V_{n-1}$  containing  $(x_{i-1}, x_i, \ldots, x_{i+k+1})$ , but not containing c. We extend this to a nice set D in  $V_0 \cup \ldots \cup V_n$  by adding the necessary elements of  $V_n$  and continue as in Lemma 2.19.

**PROPOSITION 2.29.** For any set A we have acl(A) = dcl(A).

**PROOF.** It suffices to prove the proposition for finite sets A. If |A| = 1, we have  $\operatorname{acl}(A) = A$  by  $\omega$ -saturation of  $M_n$ . Now assume that the statement is true for A and let  $a \notin \operatorname{acl}(A) = \operatorname{dcl}(A)$ . By Corollary 2.21 it suffices to prove the following:

CLAIM 2.30. We have  $c \in \operatorname{acl}(aA) \setminus \operatorname{acl}(A)$  if and only if there is a reduced path  $\gamma = (a = x_0, \ldots, x_s = b)$  for some  $b \in \operatorname{acl}(A)$  changing direction in c.

Clearly, if there is a reduced path  $\gamma = (a = x_0, ..., x_s = b)$  for some  $b \in acl(A)$  changing direction in c, then  $c \in dcl(aA)$  by Proposition 2.31 and Corollary 2.21. For the converse, it suffices to construct a nice set D containing  $A \cup \{a\}$  but not c. So let  $D_0$  be a nice finite set containing A and not c and let  $\gamma = (a = x_0, ..., x_m = b)$  be a reduced path with  $b \in acl(A)$ . Exactly as in Proposition 2.28, we construct a nice set  $D \ge D_0$  containing a but not c.

In Section 4 we will see that in the prime model the algebraic closure will be described by reduced words in the Coxeter group associated to the building.

We next show that algebraically closed sets are *weakly gated* in the following sense:

**PROPOSITION 2.31.** For any set A and  $a \notin \operatorname{acl}(A)$ , there is a flag  $C \in \operatorname{acl}(A)$  such that for any  $b \in \operatorname{acl}(A)$  and any reduced path from a to b there is an equivalent one

entering acl(A) through one of the elements of C and this is the last vertex where the path changes direction.

The proposition follows from the following lemma:

LEMMA 2.32. Let  $b_1, b_2 \in acl(A)$ ,  $a \notin acl(A)$  and let  $\gamma_1 = (a = x_0, \dots, x_s = b_1)$ and  $\gamma_2 = (a = y_0, \dots, y_t = b_2)$  be reduced paths. Assume that  $\gamma_1, \gamma_2$  are chosen in their equivalence classes such that i, j are minimal with  $x_i, y_j \in acl(A)$ . Then  $x_i \in R(y_j)$ .

**PROOF.** Let  $\gamma$  denote the extension of  $\gamma_2$  by  $(x_i, \ldots, x_1, a)$ . If  $\gamma$  is reduced, then  $(x_i, \ldots, a, y_1, \ldots, y_j)$  is a flag by minimality of i, j and we are done. Now assume that  $\gamma$  is not reduced.

**Case I:** Suppose that  $\gamma$  does not change direction in a. To fix notation let  $x_1 \in R_{<}(a)$ . We do induction on min $\{i, j\}$ . By symmetry we may assume  $i \leq j$ . First assume i = 1. Let m be minimal such that  $x_1 \notin R(y_m)$ . Note that in this case  $\gamma_2$  must change direction in  $y_{m-1}$ . If j < m we are done. Otherwise as in Remark 2.14 we obtain a reduced path  $(x_1, \ldots, y_{m-1}, y_m, \ldots, y_t)$  containing a flag from  $x_1$  to  $y_{m-1}$ . Then  $y_{m-1} \in acl(A)$ , contradicting the minimality of j.

So now assume that i > 1 and consider the path  $\gamma' = (x_1, a, y_1, \ldots, y_t)$ . If  $\gamma'$  is reduced we may apply the induction hypothesis to  $\gamma'$  and the path from  $x_1$  to  $b_1$ . Otherwise we reduce  $\gamma'$  as in Remark 2.14: let *m* be minimal such that  $x_1 \notin R(y_m)$ , so  $y_{m-1} = x_1 \lor y_m$ . If j > m we obtain again a reduced path  $(x_1, \ldots, y_{m-1}, y_m, \ldots, y_t)$ containing a flag from  $x_1$  to  $y_{m-1}$  and use the induction hypothesis on this path and  $\gamma_1$  starting from  $x_1$ . If j < m we have  $x_1 \in R(y_j)$ . Let *k* be minimal such that  $x_k \notin R(y_j)$ . If k > i we are done. Otherwise we obtain a reduced path  $(y_{j-1}, \ldots, x_{k-1}, x_k, \ldots, x_s = b_1)$  and apply the induction hypothesis to this and  $\gamma_2$ starting from  $y_{j-1}$ .

**Case II:** Suppose that  $\gamma$  changes direction in a, say  $x_1 \in R_{>}(a)$ . We do induction on min $\{i, j\}$ . By symmetry we may assume  $i \leq j$ . Let m be minimal such that  $\gamma_2$  changes direction in  $y_m$ .

We start with i = 1. Since  $\gamma$  is not reduced, we have  $x_1 \in R(y_m)$  since we must have  $x_1 \wedge y_m = x_1$ . Since  $a \notin acl(A)$ , the path  $(x_1, a, y_1, \dots, y_j)$  is not reduced as it changes direction in a. But  $(a, y_1, \dots, y_j)$  is reduced and we have  $a, y_m \in R(x_1)$ . Since  $x_1, y_1$  have the same level, we conclude that  $x_1 = y_1 \in acl(A)$  and we finish.

Now assume i > 1 and consider the path  $\gamma' = (x_1, a, y_1, \dots, y_t)$ . If  $\gamma'$  is reduced, we may apply the induction hypothesis to the paths starting at  $x_1$ . Otherwise we have as before  $x_1 \in R(y_m)$ . If  $x_1 \notin R(y_j)$ , then by the proof of Lemma 2.14 we see that the path  $(y_j, \dots, y_1, a, x_1, \dots, x_i)$  is reduced and changes direction in a implying that  $a \in acl(A)$ . Hence  $x_1 \in R(y_j)$ . In  $\gamma_2$  we may therefore replace  $(a, y_1, \dots, y_j)$  by  $(a, x_1, \dots, y_j)$  and apply the induction hypothesis to  $\gamma_1, \gamma_2$  starting from  $x_1 = y_1$ .  $\dashv$ 

COROLLARY 2.33. For any type tp(a/A), if C is a flag in acl(A) such that for any  $b \in acl(A)$  and any reduced path from a to b there is an equivalent one entering acl(A) through one of the elements of C and if no proper subset of C has this property, then C is uniquely determined.

**PROOF.** Let  $D \subset \operatorname{acl}(A)$  be another flag with the same property and let  $d \in D \setminus C$ . If  $\gamma = (a, \ldots, d)$  is a reduced path, we may assume  $\gamma$  to enter  $\operatorname{acl}(A)$  via an element  $c \in C \subset \operatorname{acl}(A)$ . Thus we may assume that  $\gamma = (a, \ldots, c, \ldots, d)$ . Now by assumption there is an equivalent path from a to c entering  $\operatorname{acl}(A)$  via an element d' of D. So we may assume that  $\gamma = (a, \ldots, d', \ldots, c, \ldots, d)$  is a reduced path. Thus any reduced path from a to an element of acl(A) which contains d is equivalent to one that enters acl(A) via some element of  $D \setminus \{d\}$ , contradicting the minimality of D.

The flag C is called the *projection* from a to A and we write C = proj(a/A). Note that  $\text{proj}(a/A) = \emptyset$  if and only if a is not connected to any vertex of acl(A).

It is now easy to show the following:

THEOREM 2.34. The theory  $T_n$  is  $\omega$ -stable.

PROOF. Let M be a countable model and let  $\overline{d}$  be a tuple from  $\overline{M}$ . Let  $C \in M$  be the finite set of projections from  $\overline{d}$  to M. Then the type  $tp(\overline{d}/M)$  is determined by  $tp(\overline{d}/C)$ . By Lemma 2.19,  $\overline{d} \cup C$  is contained in a nice finite subset of  $M_n$  and for such subsets the quantifier-free type determines the type by Remark 2.26. Hence there are only countably many types over a countable model.

In fact, using the characterization of forking independence in stable theories (see [6] Ch. 8) it is easy to see directly without counting types that  $T_n$  is superstable. Note that since for any set A we have acl(A) = dcl(A) this is an example of a stationary independence relation as defined in [7].

**THEOREM 2.35.** Forking independence in models of  $T_n$  is given by

$$A \bigcup_{C} B$$

if and only if for all  $a \in acl(AC)$ ,  $b \in acl(BC)$  and any reduced path from a to b there is an equivalent path passing through an element of acl(C).

PROOF. It is easy to see that this notion of independence satisfies the characterizing properties of forking in stable theories (see [6] Ch. 8) and hence agrees with the usual one. The existence of nonforking extensions follows from the construction of  $M_n$  as a Hrushovski limit. Since we have just seen that for any type tp(a/A) there is a finite set  $A_0 \subseteq acl(A)$  such that  $a 
otherwise the shows directly (without counting types) that <math>T_n$  is superstable.

For convenience we also state the following special case:

COROLLARY 2.36. The vertex *a* is independent from *A* over *C* if  $proj(a/AC) \subseteq acl(C)$ , i.e., if for any  $b \in acl(AC)$  connected to *a* and any reduced path from *a* to *b* there is an equivalent path passing through an element of acl(C). In particular, *a* is independent from *A* over  $\emptyset$  if and only if *a* is not connected to any vertex of acl(A) by *a* path.

This characterization of forking now directly implies the following:

COROLLARY 2.37. The free pseudospace has weak elimination of imaginaries, i.e., any type has a canonical basis consisting of a finite set of real elements: for any vertex a in a saturated model  $\overline{M}$  of  $T_n$ , the projection  $\operatorname{proj}(a/M)$  of a onto a saturated elementary submodel M is a canonical basis for the type of a over M.

**§3. Ampleness.** We now recall the definition of a theory being *n*-ample given by Pillay and Evans in [2, 5], in slightly more symmetrized form:

DEFINITION 3.1. A theory T is called *n*-ample if (possibly after naming parameters) there are tuples  $a_0, \ldots a_n$  in M such that the following holds:

1. for i = 0, ..., n - 1 we have

$$\operatorname{acl}(a_0,\ldots,a_{i-1},a_i) \cap \operatorname{acl}(a_0,\ldots,a_{i-1},a_{i+1}) = \operatorname{acl}(a_0,\ldots,a_{i-1});$$

- 2.  $a_i \not\perp a_j$  for  $i, j = 1, \ldots n$ , and
- 3.  $a_0 \ldots a_{i-1} \, \bigcup_{a_i} a_{i+1}, \ldots, a_n$  for  $i = 1, \ldots n 1$ .

**REMARK** 3.2. If  $a_0, \ldots a_n$  witness *n*-ampleness over some parameters *A*, then for any  $i = 1, \ldots n$  we have

$$\operatorname{acl}(a_0, a_i) \cap \operatorname{acl}(a_0, a_{i-1}) \subseteq \operatorname{acl}(a_0 A).$$

**THEOREM 3.3.** The theory  $T_n$  is n-ample, but not n + 1-ample.

**PROOF.** Any maximal flag  $(x_0, \ldots x_n)$  in  $M_n$  is a witness for *n*-ampleness. This follows immediately from the description of acl in Proposition 2.28 and of forking in Corollary 2.36.

To see that the free pseudospace of dimension n is not n + 1-ample, suppose toward a contradiction that  $A_0, \ldots A_{n+1}$  are witnesses for  $T_n$  being n + 1-ample over some set of parameters A. We have

$$A_{n+1} \underbrace{\not\downarrow}_{A} A_0,$$
$$A_0 \dots A_{i-1} \underbrace{\downarrow}_{AA_i} A_{i+1} \dots A_{n+1}, i = 0, \dots n.$$

By the first condition there are vertices  $a_0 \in \operatorname{acl}(A_0A) \setminus \operatorname{acl}(A)$  and  $a_{n+1} \in \operatorname{acl}(A_{n+1}A) \setminus \operatorname{acl}(A)$  which are in the same connected component and not independent over A. We may choose  $a_0$  and  $a_{n+1}$  at minimal distance with this property, i.e., in such a way that no reduced path from  $a_0$  to  $a_{n+1}$  contains an element  $b \in \operatorname{acl}(A_0A)$  with  $b \not\perp_A a_{n+1}$ .

Since  $a_0 
int_{A_nA} a_{n+1}$  and  $a_0 
int_A a_{n+1}$  by the characterization of independence in Corollary 2.36 there is a flag  $C_n \subset \operatorname{acl}(A_nA)$ ,  $C_n \not\subset \operatorname{acl}(A)$  such that for any reduced path from  $a_0$  to  $a_{n+1}$  there is an equivalent one passing through an element of  $C_n$ , so  $a_0 
int_{C_n} a_{n+1}$ . Clearly, we may assume that  $C_n$  is minimal with this property. Let  $\gamma$  be a reduced path from  $a_0$  to  $a_{n+1}$  not equivalent to a path containing an element of  $\operatorname{acl}(A)$ . Let  $a_n \in C_n$  be such that  $\gamma$  or an equivalent path passes through  $a_n$ , then  $a_0 
ot \ A_a a_n$ . If  $\gamma$  changes direction in some b between  $a_0$  and  $a_n$ or in  $a_n$ , then by the previous remark  $b \in \operatorname{acl}(a_0, a_{n+1}) \cap \operatorname{acl}(a_0, a_n) \subseteq \operatorname{acl}(A_0A)$ and  $b 
ot \ A_a a_{n+1}$ , contradicting the choice of  $a_0$ . Therefore  $(a_0, a_n)$  form a flag with  $a_n \notin V_n \cup V_0$  as otherwise  $\gamma$  changes direction in  $a_n$ . Now there is some flag  $C_{n-1} \in \operatorname{acl}(A_{n-1}A), C_{n-1} \not\subset \operatorname{acl}(A)$  such that  $a_0 
othermal contradicting the choice of <math>i = 1, \ldots n-1$ . This is impossible if  $a_n \notin V_0 \cup V_n$ .

The proof shows that in fact the following stronger ampleness result holds:

COROLLARY 3.4. If  $a_0, \ldots a_n$  are witnesses for  $T_n$  being *n*-ample, then there are vertices  $b_i \in \operatorname{acl}(a_i)$  such that  $(b_0, \ldots b_n)$  is a flag.

The following was pointed out by Itay Ben Yacov:

REMARK 3.5. If we let  $T_{\omega}$  denote the theory of  $\omega$ -colored graphs with vertices of type  $\bigcup_{i < \omega} V_i$  such that for all  $i, j \in \omega$  the restriction to  $V_i \cup \ldots V_{i+j}$  is a model

of  $T_j$ , we see that  $T_{\omega}$  is  $\omega$ -stable, *n*-ample for all  $n < \omega$  and does not interpret an infinite group.

§4. Buildings and the prime model of  $T_n$ . Let M be a model of  $T_n$ . We say that maximal flags  $\zeta_1, \zeta_2$  are *totally connected* if for i = 0, ..., n - 1 the  $V_i$ -vertices of  $\zeta_1, \zeta_2$  are  $E_i$ - and  $E_{i+1}$ -connected (whenever this makes sense). Given a flag  $\zeta$  in M we let  $M^0(\zeta)$  be its totally connected component. Since  $M_n$  is homogeneous for maximal flags, all totally connected components of  $M_n$  are isomorphic.

The uniqueness part of Proposition 4.4 follows directly from the following theorem and Proposition 5.1 of [4] which states that this type of building is uniquely determined by its associated Coxeter group and the cardinality of the residues.

THEOREM 4.1.  $M_n^0$  is a building of type  $A_{\infty,n+1}$  all of whose residues have cardinality  $\aleph_0$ .

Recall the following definitions (see, e.g., [3]). Let W be the Coxeter group

$$W = \langle t_0, \dots, t_n : t_i^2 = (t_i t_k)^2 = 1, i, k = 0 \dots n, |k - i| \ge 2 \rangle,$$

whose associated diagram we call  $A_{\infty,n+1}$ .

DEFINITION 4.2. A building of type  $A_{\infty,n+1}$  is a set  $\Delta$  with a *Weyl distance function*  $\delta : \Delta^2 \to W$  such that the following two axioms hold:

- 1. For each  $s \in S := \{t_i, i = 0, ..., n\}$ , the relation  $x \sim_s y$  defined by  $\delta(x, y) \in \{1, s\}$  is an equivalence relation on  $\Delta$  and each equivalence class of  $\sim_s$  has at least 2 elements.
- 2. Let  $w = r_1 r_2 \dots r_k$  be a shortest representation of  $w \in W$  with  $r_i \in S$  and let  $x, y \in \Delta$ . Then  $\delta(x, y) = w$  if and only if there exists a sequence of elements  $x, x_0, x_1, \dots, x_k = y$  in  $\Delta$  with  $x_{i-1} \neq x_i$  and  $\delta(x_{i-1}, x_i) = r_i$  for  $i = 1, \dots, k$ .

A sequence as in 2. is called a *gallery of type*  $(r_1, r_2, ..., r_k)$ . The gallery is called reduced if the word  $w = r_1 r_2, ..., r_k$  is reduced, i.e., a shortest representation of w. We now show how to consider  $M_n^0$  as a building of type  $A_{\infty,n+1}$ .

**PROOF.** (of Theorem 4.1) We extend the set of edges of the n + 1-colored graph  $M_n^0$  by putting edges between any two vertices that are incident in the sense of Definition 2.1.2.1. In this way, flags of  $M_n^0$  correspond to a complete subgraph of this extended graph, which thus forms a simplicial complex. A maximal simplex consists of n + 1 vertices each of a different type  $V_i$ . (Such a simplex is called a *chamber*.) Let  $\Delta$  be the set of maximal simplices in this graph. Define  $\delta : \Delta^2 \to W$  as follows:

Put  $\delta(x, y) = t_i$  if and only if the flags x and y differ exactly in the vertex of type  $V_i$ . Extend this by putting  $\delta(x, y) = w$  for a reduced word  $w = r_1 r_2 \dots r_k$  if and only if there exists a sequence of elements  $x = x_0, x_1, \dots, x_k = y$  in  $\Delta$  with  $x_{i-1} \neq x_i$  and  $\delta(x_{i-1}, x_i) = r_i$  for  $i = 1, \dots, k$ .

Clearly, with this definition of  $\delta$ , the set  $\Delta$  satisfies the first condition of Definition 4.2. In fact, for all  $s \in S$  every equivalence class  $\sim_s$  has cardinality  $\aleph_0$ . We still need to show that  $\delta$  is well-defined, i.e., we have to show the following for any  $x, y \in \Delta$ : if there are reduced galleries  $x_0 = x, x_1, \ldots, x_k = y$  and  $y_0 = x, y_1, \ldots, y_m = y$  of type  $(r_1, r_2, \ldots, r_k)$  and  $(s_1, \ldots, s_m)$ , respectively, then

in W we have  $r_1r_2...r_k = s_1...s_m$ . Equivalently, we will show the following, which completes the proof of Theorem 4.1:

**Claim:** There is no reduced gallery  $a_0, a_1, \ldots, a_k = a_0$  for k > 0 in  $M_n^0$ .

PROOF OF CLAIM. Suppose otherwise. Let  $a_0, a_1, \ldots, a_k = a_0$  be a reduced gallery of type  $(r_1, \ldots, r_k)$  for some k > 0. Note that the flags  $a_{i-1}$  and  $a_i$  contain the same vertex of type  $V_i$  as long as  $r_i \neq t_i$ .

Now consider the sequence of vertices of type  $V_n$  and  $V_{n-1}$  occurring in this gallery. Since  $V_n \cup V_{n-1}$  contains no cycles, the sequence of vertices of type  $V_n$  and  $V_{n-1}$  occurring in this gallery will be of the form

$$(x_1, y_1, x_2, y_2, \dots, x_i, y_i, x_i, y_{i-1}, \dots, y_1, x_1)$$
(1)

with  $x_i \in V_n$ ,  $y_i \in V_{n-1}$  and  $x_i$  a neighbor of  $y_i$  and  $y_{i-1}$  in the original graph. This implies that at some place in the gallery type there are two occurrences of  $t_n$  which are not separated by an occurrence of  $t_{n-1}$  (or conversely). Since  $t_n$  commutes with all  $t_i$  for  $i \neq n - 1$  and the word  $r_1 \dots r_k$  is reduced, the first case cannot happen. Hence there are two occurrences of  $t_{n-1}$  which are not separated by an occurrence of  $t_n$ , say  $r_j$ ,  $r_{j+m} = t_{n-1}$  with  $r_{j+1}, \dots, r_{j+m-1} \neq t_n$ .

We now consider the gallery  $a_j, \ldots a_{j+m}$  of type  $(r_j = t_{n-1}, r_{j+1}, \ldots, r_{j+m} = t_{n-1})$ . Notice that by (1), the flags  $a_j$  and  $a_{j+m}$  have the same  $V_n$  and the same  $V_{n-1}$  vertex. Since  $M_n^0$  does not contain any  $E_{n-1}$ -cycles, the sequence of  $V_{n-1}$ - and  $V_{n-2}$ -vertices appearing in this sequence must again be of the same form as in (1). Exactly as before we find two occurrences<sup>2</sup> of  $t_{n-2}$  in the gallery type of  $a_j, \ldots a_{j+m}$  which are not separated by an occurrence of  $t_{n-1}$ . Continuing in this way, we eventually find two occurrences of  $t_1$  which are not separated by any  $t_i$ . Since  $t_1^2 = 1$  this contradicts the assumption that the gallery be reduced.

The proof shows in fact the following:

COROLLARY 4.3. A model of  $T_n$  is a building of type  $A_{\infty,n+1}$  if and only if it is  $E_i$ -connected for all i and every vertex is contained in a maximal flag.

**PROPOSITION 4.4.**  $M^0(\zeta)$  is a model of  $T_n$ . Furthermore  $M_n^0$  is the unique countable model of  $T_n$  which is  $E_i$ -connected for i = 1, ... n and such that every vertex is contained in a maximal flag.

(The corresponding Remark 3.6 of [1] uses Lemma 3.2, which is not correct as phrased there:  $M_n^0$  and  $M_n^0 \cup \{a\}$  with *a* an isolated point are not isomorphic, but satisfy the assumptions of Remark 3.6.)

THEOREM 4.5. The building  $M_n^0$  is the prime model of  $T_n$ .

**PROOF.** To see that  $M_n^0$  is the prime model of  $T_n$  note that for any flags  $C_1, C_2 \in M_n^0$  and gallery  $C_1 = x_0, \ldots, x_k = C_2$  the set of vertices occurring in this gallery is  $E_i$ -connected for all *i*. Hence by Remark 2.26 its type is determined by the quantifier-free type.

Thus, given a maximal flag M in any model of  $T_n$  and a maximal flag  $c_0$  of  $M_n^0$  we can embed  $M_n^0$  into M by moving along the galleries of  $M_n^0$ .

<sup>&</sup>lt;sup>2</sup>If  $t_{n-2}$  does not occur in the type of the gallery, this would contradict the assumption that the type is reduced since  $t_{n-1}$  commutes with all  $t_i$  for  $i \neq n, n-2$ .

§5. Ranks and types. Recall that for vertices  $x, y \in M_n^0$  with  $x \in V_i, y \in V_j$ the *Weyl-distance*  $\delta(x, y)$  equals  $w \in W$  if there are flags  $C_1, C_2$  containing x, y, respectively, with  $\delta(C_1, C_2) = w'$  and such that w is the shortest representative of the double coset  $\langle t_k : k \neq i \rangle w' \langle t_k : k \neq j \rangle$  (where as usual  $\langle X \rangle$  denotes the subgroup of W generated by X).

Note that the Weyl distance between two vertices describes exactly the reduced paths between these vertices. Therefore we have the following:

**PROPOSITION 5.1.** The theory  $T_n$  has quantifier elimination in a language containing predicates  $\delta_w^{i,j}$  for Weyl distances between vertices of type  $V_i$  and of type  $V_j$ .

**PROOF.** Since the predicates  $\delta_w^{i,j}$  describe the reduced paths between vertices, in this language two tuples have the same quantifier-free type if and only if they are contained in isomorphic nice sets. For nice sets the quantifier-free type determines the type, whence the claim.  $\dashv$ 

Using the description of forking given in Theorem 2.35 it is easy to give a list of regular types such that any nonalgebraic type is nonorthogonal to one of these. This is entirely similar to the list given in [1] and we omit the details. It is also clear from this description of forking that the geometry on these types is trivial.

For any small set A in a large saturated model we have the following kinds of regular types:

- (I)  $\operatorname{tp}(a/A)$  where  $a \in V_i$  is not connected to any element in  $\operatorname{acl}(A)$
- (II) tp(a/A) where  $a \in V_i$  is incident with some  $b \in acl(A) \cap V_j$  but not connected in R(b) to any vertex in  $acl(A) \cap R(b)$ .
- (III)  $\operatorname{tp}(a/A)$  where  $a \in V_i$  is incident with some  $x, y \in \operatorname{acl}(A)$  such that (x, a, y) is a flag with  $x \in V_k, y \in V_j$ ; and as a special case of this we have
- (IV)  $\operatorname{tp}(a/A)$  where  $a \in V_i$  has neighbors  $x, y \in \operatorname{acl}(A)$  such that (x, a, y) is a (necessarily dense) flag.

By quantifier elimination any of these descriptions determines a complete type. Using the description of forking in Corollary 2.36 one sees easily that each of these types is regular and trivial.

Clearly, any type in (IV) has *U*-rank 1 and in fact Morley rank 1 by quantifier elimination. It also follows easily that  $MR(a/A) = \omega^n$  if tp(a/A) is as in (I). In case (II) we find that  $MR(a/A) = \omega^{n-j-1}$  or  $MR(a/A) = \omega^{j-1}$  depending on whether or not i < j. In case (III) we have  $MR(a/A) = \omega^{|k-j|-2}$ .

Just as in [1] we obtain:

LEMMA 5.2. Any regular type in  $T_n$  is nonorthogonal to a type as in (I), (II), or (III).

**PROOF.** Let  $p = \operatorname{tp}(b/\operatorname{acl}(B))$ . If b is not connected to  $\operatorname{acl}(B)$ , then p is as in (I), so we may assume that  $\operatorname{proj}(b/B) = C \neq \emptyset$ . Let a be a vertex on a short path from b to C incident with an element of C. Then by Corollary 2.36 we see that p is nonorthogonal to  $\operatorname{tp}(a/C)$  and  $\operatorname{tp}(a/C)$  is of type (II) or (III).

§6. Acknowledgment. I would like to thank the referee for her/his very careful work and for pointing out several mistakes in an earlier version of this paper and Linus Kramer for providing reference [4].

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