

# MOST PERMUTATIONS POWER TO A CYCLE OF SMALL PRIME LENGTH

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*Abstract* We prove that most permutations of degree  $n$  have some power which is a cycle of prime length approximately  $\log n$ . Explicitly, we show that for  $n$  sufficiently large, the proportion of such elements is at least  $1 - 5/\log \log n$  with the prime between  $\log n$  and  $(\log n)^{\log \log n}$ . The proportion of even permutations with this property is at least  $1 - 7/\log \log n$ .

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## 1. Introduction

The symmetric and the alternating groups  $S_n$  and  $A_n$  of degree  $n$  have been viewed as probability spaces with the uniform distribution since the seminal work of Gruder [9], Gončarov [7] and Erdős-Turán [4]. There is an analogy between the disjoint cycle decomposition of a permutation and the prime factorization of an integer whereby the cycles correspond to prime numbers (see [5, 8] for a description). The probability that a permutation  $g \in S_n$  is an  $n$ -cycle is  $1/n$ , while the probability that a number  $p \leq x$  is prime is  $1/\log x$ , and according to this analogy,  $n$  corresponds to  $\log x$ . Counting integers without small prime factors is like counting permutations without small cycles: both lead to limits reminiscent of Mertens' third theorem:

$$\lim_{n \rightarrow \infty} \left( \log n \prod_{p \leq n} \left( 1 - \frac{1}{p} \right) \right) = e^{-\gamma}, \quad \text{where } \gamma = 0.5772 \dots$$

is the Euler–Mascheroni constant. Statistical methods have been invaluable for proving theoretical results. For example, the number of cycles of a uniformly distributed random

element  $g \in S_n$  (or  $g \in A_n$ ) behaves as  $n \rightarrow \infty$  like a normal distribution  $N(\mu, \sigma^2)$  (c.f. [6, Theorem 1]) with mean and variance  $\mu = \sigma^2 = \log n$ , see [7, 15]. This paper focuses on proving the following theorem. We abbreviate  $\log(n)$  and  $\log(\log(n))$  as  $\log n$  and  $\log \log n$ , respectively. All logarithms are to the natural base  $e = 2.718 \dots$ .

**Theorem 1.** *Suppose that  $G \in \{A_n, S_n\}$ . Let  $\rho_G$  be the proportion of permutations in  $G$  for which some power is a cycle of prime length  $p$ , where  $p$  lies in the open interval  $(\log n, (\log n)^{\log \log n})$ . Then, for  $n$  sufficiently large,  $\rho_G \geq 1 - c/\log \log n$  where  $c = 5$  when  $G = S_n$ , and  $c = 7$  when  $G = A_n$ .*

A permutation  $g \in S_n$  having a single cycle of length greater than 1 is called a  $k$ -cycle, where the cycle has length  $k$ . Their theoretical importance has long been recognised: the presence of a  $k$ -cycle in a primitive subgroup  $G$  of  $S_n$  was shown to imply that  $G$  is  $A_n$  or  $S_n$  by Jordan in 1873 in the case where  $k$  is prime and  $k \leq n - 3$  ([12], or see [19, Theorem 13.9]). The same conclusion also holds for arbitrary  $k < n/2$  by a result in Marggraf's dissertation [14], see also [11, Corollary 1.3], [19, Theorem 13.5] and P. M. Neumann's Mathematical Review MR0424912.

To use Jordan's result for deciding whether a given primitive subgroup  $G$  of  $S_n$  is indeed  $A_n$  or  $S_n$ , one needs to locate a  $p$ -cycle with  $p$  a prime, say by choosing random elements from  $G$ . It is inefficient to do this directly since the proportion of such elements in  $A_n$  or  $S_n$  is  $O(n^{-1})$  (see § 1.1). Instead one searches for a 'pre- $p$ -cycle', a permutation which powers to a  $p$ -cycle. It is shown in [17, Lemma 10.2.3] that the proportion of elements in  $A_n$  or  $S_n$  that power to a  $p$ -cycle with  $n/2 < p \leq n - 3$  is asymptotically  $\log 2/\log n$ , while an application of the main result Theorem 1 of [1] shows that considering only pre- $p$ -cycles with  $p$  bounded, say  $p \leq m$ , produces a proportion  $c(m)/n^{1/m}$ . Thus, to approach Seress's asymptotic proportion, the primes  $p$  must be allowed to grow unboundedly with  $n$ .

Quite decisively, and perhaps surprisingly, the third author recently showed [18, Theorem 2] that *almost all* permutations in  $G \in \{A_n, S_n\}$  are pre- $p$ -cycles for some prime  $p$ . The purpose of this paper is to prove an even stronger statement: namely that, asymptotically, almost all permutations in  $G$  power to a  $p$ -cycle where the prime  $p$  is roughly  $\log n$  (Theorem 1) and we may derive, from Proposition 6 and the proof of Theorem 1, explicit values for  $n_G$  such that the bounds of Theorem 1 hold for all  $n \geq n_G$ . On the one hand, we quantify the asymptotic results [18], giving a precise analysis with explicit bounds rather than asymptotics. In addition, we show (which came as a surprise to the authors) that the prime  $p$  can be chosen in a very small interval of length logarithmic in  $n$ . Theorem 1 relies on Proposition 6 which strengthens a key technical result [4, Theorem VI] of Erdős and Turán. Theorem 10 shows that the proportion of pre- $p$ -cycles in  $G$  is at least  $1 - c' \log \log n / \log(n - 3)$  for  $p$  in the range  $2 \leq p \leq n - 3$ . We prove in Remark 11 that for all  $n \geq 5$ , the proportion of pre- $p$ -cycles in  $S_n$  with  $2 \leq p \leq n - 3$  is at least  $1/19$ . In the remainder of this section, we comment on uses and proofs of these results.

### 1.1. Recognising finite symmetric and alternating groups

We call a permutation  $g \in S_n$  a *pre- $k$ -cycle* if it powers to a  $k$ -cycle where  $k > 1$ . For example,  $g = (1, 2, 3, 4)(5, 6)(7, 8, 9)$  is a pre-3-cycle as  $g^4 = (7, 8, 9)$ , but  $g$  is not a pre-2-cycle. A simple argument (Lemma 3) shows that the disjoint cycle decomposition of a

pre- $k$ -cycle contains exactly one  $k$ -cycle and all other cycles (if any) have lengths coprime to  $k$ . The number of  $k$ -cycles in  $S_n$  equals  $n!/c(k)$  where  $c(k) = k(n - k)!$ . Hence the proportion  $\rho_n$  of cycles in  $S_n$  equals  $\sum_{k=2}^n 1/c(k) = \sum_{j=0}^{n-2} (1/(n - j)j!)$ . For  $n$  even,

$$n\rho_n - \sum_{j=0}^{n-2} \frac{1}{j!} = \sum_{j=0}^{n-2} \frac{j}{(n - j)j!} = \sum_{j=0}^{n/2} \frac{j}{(n - j)j!} + \sum_{j=n/2}^{n-2} \frac{j}{(n - j)j!} < \frac{2}{n} \sum_{j=0}^{n/2} \frac{j}{j!} + \sum_{j=n/2}^{\infty} \frac{j}{j!}.$$

Thus,  $|n\rho_n - e| \rightarrow 0$  as  $n \rightarrow \infty$ , and so  $\rho_n = O(n^{-1})$ . By contrast, the proportion of permutations in  $S_n$  which are pre- $k$ -cycles (for  $2 \leq k \leq n$ ) approaches 1 as  $n \rightarrow \infty$ .

Let  $X \subseteq S_n$  and suppose that the subgroup  $\langle X \rangle$  of  $S_n$  generated by  $X$  is primitive. Then a probabilistic algorithm for testing whether  $\langle X \rangle$  contains  $A_n$  involves a random search for a pre- $p$ -cycle, where  $p$  is a prime less than  $n - 2$ , and it turns out that such elements have density 1 as  $n \rightarrow \infty$  (Theorem 1). If the subgroup  $\langle X \rangle$  does not contain  $A_n$ , we want to limit the (fruitless) search for pre- $p$ -cycles, and for this we need explicit lower bounds on their density in both  $A_n$  and  $S_n$ , see Remark 2.

There are two important tools in our proof. The first is to quantify the sum  $\sum_{a < p \leq b} 1/p^2$  as a function of the real numbers  $a$  and  $b$ . (Here, and henceforth,  $p$  denotes a prime.) The second is to estimate the number of permutations whose cycle lengths do not have certain ‘forbidden lengths’. Erdős-Turán [4, Theorem VI] proved that the probability  $\rho$  that the cycle lengths of  $g \in S_n$  do not lie in a subset  $\mathcal{A} \subseteq \{1, \dots, n\}$  is at most  $1/\mu$  where  $\mu = \sum_{a \in \mathcal{A}} 1/a$ . Applying a result of Ford [5, Theorem 2.9] gives the bound  $\rho \leq e^{1-\mu}$  (take  $r = 1$ ,  $T_1 = \mathcal{A}$  and  $k_1 = 0$  in Theorem 2.9). We prove in Proposition 6 that  $\rho < e^{\gamma-\mu}$ , where  $\gamma = 0.5772 \dots$ . Our bound is less than 2/3 the size of bounds of Erdős-Turán and Ford. This small improvement is needed to obtain practically useful bounds in Theorem 10. Manstavičius [13, p. 39] claimed that a slightly weaker bound, *viz.*  $\rho < e^{\gamma-\mu}(1 + 1/n)$ , followed from a pre-print of his. Unfortunately, we could not locate his pre-print.

After describing the cycle structure of pre- $k$ -cycles in §2, we prove our bound on forbidden cycle lengths in §3. The sum  $\sum_{a < p \leq b} 1/p^2$  is estimated in §4, and the main theorems are proved in §5.

**Remark 2.** Suppose that the proportion of pre- $p$ -cycles in  $A_n$  or  $S_n$  is at least  $c_0 = c_0(n)$ . Given a primitive permutation group  $G \leq S_n$ , then the probability that  $A_n \leq G$  and  $m$  independent random selections from  $G$  do *not* find a pre- $p$ -cycle for some  $p$  with  $2 \leq p \leq n - 3$ , is at most  $(1 - c_0)^m$ . This upper bound is at most a prescribed ‘failure probability’  $\varepsilon$  if and only if  $m \geq \log(\varepsilon)/\log(1 - c_0)$ . Using Theorem 1, we can take  $c_0(n) = 1 - c/\log \log n$ , and a larger value is given in Theorem 10. Therefore, we may take  $c_0$  to be an absolute constant thus requiring only  $m = m(\varepsilon)$  random selections. We prove in Remark 11 that  $c_0 = 0.05$  works for  $S_n$  and this gives  $m(\varepsilon) \geq 20 \log(\varepsilon^{-1})$ . By contrast, the standard analysis of this algorithm [17, pp. 226–227] requires  $m = O(\log(n) \log(\varepsilon^{-1}))$  random selections (because the proportion of pre- $p$ -cycles in the range  $n/2 < p < n - 2$  is approximately  $\log 2/\log n$ ). Thus, our analysis gives an algorithm that is faster by a factor of  $\log(n)$ .

### 2. Precycles

We begin with some notation. Fix  $g \in S_n$  and let  $\lambda = \lambda(g)$  be the partition of  $n$  induced by the (disjoint) cycle lengths of  $g$ . We write  $\lambda \vdash n$  and  $\lambda = \langle 1^{m_1} 2^{m_2} \dots n^{m_n} \rangle$  where  $m_k$  is the multiplicity of the part  $k$ , so  $n = \sum_{k=1}^n km_k$ . The vector  $(m_1, \dots, m_n)$  is called the *cycle type* of  $g$ . A pre- $k$ -cycle  $g \in S_n$  can be recognized by its *cycle partition*  $\lambda(g)$ .

**Lemma 3.** *Let  $g \in S_n$  be a pre- $k$ -cycle with  $\lambda(g) = \langle 1^{m_1} 2^{m_2} \dots n^{m_n} \rangle$ . Then  $m_k = 1$  and for  $i \neq k$ ,  $m_i > 0$  implies  $\gcd(k, i) = 1$ . That is,  $g$  has a unique  $k$ -cycle, and its other cycles (if any) have lengths coprime to  $k$ .*

**Proof.** Suppose  $g$  has disjoint cycle decomposition  $g_1 \dots g_r$  where  $g_i$  is a cycle of length  $\lambda_i$  and  $\lambda_1 + \dots + \lambda_r = n$ . Note that  $g^\ell = g_1^\ell \dots g_r^\ell$  and  $g_i^\ell$  is a product of  $\gcd(\lambda_i, \ell)$  cycles each of length  $\lambda_i / \gcd(\lambda_i, \ell)$ . Suppose  $g^\ell$  is a  $k$ -cycle. Then there exists an  $i$  for which  $\lambda_i / \gcd(\lambda_i, \ell) = k$  and  $\gcd(\lambda_i, \ell) = 1$ . Also,  $\lambda_j / \gcd(\lambda_j, \ell) = 1$  for  $j \neq i$ . Thus,  $\lambda_i = k, \gcd(\lambda_i, \ell) = 1$ . For  $j \neq i, \lambda_j \mid \ell$ . Hence  $\gcd(k, \lambda_j) = 1$ . □

**Corollary 4.** *For a prime  $p$ , the set of pre- $p$ -cycles in  $X \subseteq S_n$  equals*

$$\left\{ g \in X \mid m_p(g) = 1 \text{ and } \sum_{i \geq 2} m_{ip}(g) = 0 \right\}.$$

A group element whose order is coprime to  $p$ , for a prime  $p$ , is called a  *$p'$ -element*. The density of  $p'$ -elements in  $S_n$  is  $\sigma_n = \prod_{i=1}^{\lfloor n/p \rfloor} (1 - 1/(ip))$  by [4, Lemma I]. We are grateful to John Dixon for pointing out to us that the density of pre- $p$ -cycles in  $S_n$  is  $(1/p)\sigma_{n-p}$ . The density in  $A_n$  can be calculated similarly using [2, Theorem 3.3].

### 3. Forbidden cycle lengths

Given a set  $\mathcal{A} \subseteq \{1, \dots, n\}$  of ‘forbidden’ cycle lengths, let  $\rho$  be the proportion of elements of  $S_n$  having no cycle length in  $\mathcal{A}$ . Erdős and Turán prove  $\rho \leq \mu^{-1}$  in [4, Theorem VI] where  $\mu = \sum_{a \in \mathcal{A}} 1/a$ . To see that his bound is not optimal, consider the case when  $\mathcal{A} = \{1\}$  and  $\mu = 1$ . In this case,  $\rho$  is the proportion of derangements and the Erdős–Turán bound is unhelpful. However, an inclusion–exclusion argument shows that  $\rho = \sum_{k=0}^n (-1)^k / k!$ . As  $e^{-1} < \rho$  when  $n$  is even and  $\rho < e^{-1}$  otherwise, we expect for large  $n$  a bound of the form  $c/e$  where  $c = 1 + o(1)$ . We prove in Proposition 6 that there is a bound of the form  $\rho \leq ce^{-\mu}$  where  $c = 1.781\dots$ . The constant  $c$  equals  $e^\gamma$  where  $\gamma = 0.5772\dots$  is the Euler–Mascheroni constant. Our proof uses exponential generating functions, and has similarities to calculations used by Gruder [9], 15 years before Erdős and Turán [4].

If  $\mu < 1$ , then the Erdős–Turán bound gives no information. On the other hand, if  $\mu \geq 1$ , then we will show that the upper bound we obtain in Proposition 6, namely  $e^{\gamma-\mu}$ ,

is strictly less than the bound  $\mu^{-1}$  of Erdős and Turán. We use the fact that  $e^\mu \geq e\mu$ , since the tangent to the curve  $y = e^x$  at  $x = 1$  is  $y = ex$ . Hence

$$e^{\gamma-\mu} = \frac{c}{e^\mu} < \frac{1.782}{e^\mu} \leq \frac{1.782}{e\mu} < \frac{2}{3\mu}.$$

Similarly, our bound improves the bound  $e^{1-\mu}$  of Ford [5, Theorem 2.9] by a factor of  $e^{1-\gamma} < 2/3$ . Note that Ford’s bound can be sharpened using estimations in the proof of Proposition 6.

The centralizer of  $g$  in  $S_n$  has order  $C(\lambda(g)) := \prod_{k=1}^n k^{m_k} m_k!$ , and hence the conjugacy class  $g^{S_n}$  has size  $n!/C(\lambda(g))$ . Since the conjugacy classes in  $S_n$  are parameterized by the partitions of  $n$ , it follows that  $\sum_{\lambda \vdash n} n!/C(\lambda) = n!$ , or equivalently that  $\sum_{\lambda \vdash n} 1/C(\lambda) = 1$ .

The reader when trying to compare our analysis with that in [4] may wish to know that the quantity  $L(P)$  defined on [4, p. 159] for a permutation  $P$  should be the number of cycles of  $P$  with length equal to some  $a_\nu \in \mathcal{A}$ , and not the number of  $a_\nu \in \mathcal{A}$  with the property that  $P$  has a cycle of length  $a_\nu$ .

**Remark 5.** Given a set  $X$  and a property  $P$ , let  $\rho(X)$  denote the proportion of elements  $x \in X$  that have property  $P$ . If  $Y \subseteq X$ , then  $|Y|\rho(Y) \leq |X|\rho(X)$ . To see this let  $P(X)$  be the subset of  $x \in X$  that have property  $P$ . Since  $|P(X)| = |X|\rho(X)$  and  $P(Y) \subseteq P(X)$ , we have  $|Y|\rho(Y) \leq |X|\rho(X)$ .

**Proposition 6.** *The proportion  $\rho(S_n)$  of elements of  $S_n$  having no cycle length in  $\mathcal{A} \subseteq \{1, \dots, n\}$  satisfies  $\rho(S_n) < e^{\gamma-\mu}$  where  $\mu = \sum_{a \in \mathcal{A}} 1/a$  and  $e^\gamma \approx 1.781$ . The proportion  $\rho(A_n)$  of elements of  $A_n$  with no cycle length in  $\mathcal{A}$  satisfies  $\rho(A_n) < 2e^{\gamma-\mu}$ .*

**Proof.** For  $k \in \{1, \dots, n\}$ , let  $p_k$  equal 0 if  $k \in \mathcal{A}$ , and 1 otherwise. Set

$$P(z) = \sum_{k=1}^n \frac{p_k z^k}{k} \quad \text{and} \quad Q(z) = e^{P(z)}.$$

As  $P(z)$  is a polynomial,  $Q(z)$  is a differentiable function of the complex variable  $z$ , so its Taylor series  $\sum_{m=0}^\infty q_m z^m$  at  $z = 0$  converges to  $Q(z)$  for all  $z$ . We consider the coefficient  $q_n$ . Since

$$Q(z) = \prod_{k \notin \mathcal{A}, k \leq n} e^{z^k/k} = \prod_{k \notin \mathcal{A}, k \leq n} \sum_{m_k=0}^\infty \frac{(z^k/k)^{m_k}}{m_k!},$$

the term  $q_n z^n$  is a sum whose summands have the form

$$\prod_{k \notin \mathcal{A}, k \leq n} \frac{z^{km_k}}{k^{m_k} (m_k)!} = \frac{z^n}{C(\lambda)} \quad \text{for each partition } \lambda = \langle k^{m_k} \rangle_{k \notin \mathcal{A}} \text{ with } n = \sum_{k \notin \mathcal{A}} km_k.$$

It follows that  $q_n = \sum_{\lambda \vdash n} 1/C(\lambda)$  where the parts of  $\lambda$  do not lie in  $\mathcal{A}$  (equivalently  $m_k = 0$  for all  $k \in \mathcal{A}$ ). This is precisely the proportion  $\rho(S_n)$  of elements of  $S_n$  none of whose cycle lengths lie in  $\mathcal{A}$ . Alternatively, [9, Eq. (18)] shows that  $q_n = \rho(S_n)$ .

We now compute an upper bound for  $q_n$ . Define the  $n$ th harmonic number to be

$$H_n = \sum_{k=1}^n \frac{1}{k}, \quad \text{and let } \gamma = \lim_{n \rightarrow \infty} (H_n - \log n) = 0.5772 \dots$$

be the Euler–Mascheroni constant. The error term  $E(n) = H_n - \log n - \gamma$  satisfies  $0 < E(n) < 1/(2n)$  by [10, p. 75]. Since  $e^x < 1 + x + x^2 + \dots = 1/(1 - x)$  for  $0 < x < 1$ , it follows that  $e^{1/(2n)} < 1 + 1/(2n - 1) \leq 1 + 1/n$ . Hence

$$1 < e^{E(n)} < 1 + \frac{1}{n} \quad \text{and} \quad P(1) = \sum_{k \notin A, k \leq n} \frac{1}{k} = H_n - \mu = \log n + \gamma - \mu + E(n).$$

The previous display implies

$$Q(1) = e^{P(1)} = e^{\log n} e^{\gamma - \mu} e^{E(n)} < n e^{\gamma - \mu} \left(1 + \frac{1}{n}\right) = (n + 1) e^{\gamma - \mu}. \tag{1}$$

Differentiating the equation  $Q(z) = e^{P(z)}$  gives

$$Q'(z) = P'(z)Q(z), \quad \text{that is, } \sum_{k=0}^{\infty} k q_k z^{k-1} = \left(\sum_{k=0}^{n-1} p_{n-k} z^{n-1-k}\right) \left(\sum_{k=0}^{\infty} q_k z^k\right).$$

Equating the coefficients of  $z^{n-1}$  on both sides gives

$$n q_n = \sum_{k=0}^{n-1} p_{n-k} q_k \leq \sum_{k=0}^{n-1} q_k \leq \left(\sum_{k=0}^{\infty} q_k\right) - q_n = Q(1) - q_n.$$

It follows from (1) that

$$(n + 1) q_n \leq Q(1) < (n + 1) e^{\gamma - \mu}.$$

Cancelling  $n + 1$  gives  $q_n < e^{\gamma - \mu}$ , that is,  $\rho(S_n) < e^{\gamma - \mu} < 1.782 e^{-\mu}$ . Finally, Lemma 5 implies that  $\rho(A_n) \leq 2\rho(S_n) < 2e^{\gamma - \mu} < 3.563 e^{-\mu}$ . □

#### 4. Estimating sums

In this section, we bound the quantity  $\mu$  in Proposition 6. One approach is to use the Stieltjes integral [16, p. 67]. Instead, we take an elementary approach using finite sums, which we now briefly review. In the sequel,  $k$  denotes an integer and  $p$  denotes a prime. The *backward difference* of a function  $f(k)$  is defined by  $(\nabla f)(k) = f(k) - f(k - 1)$ . An

easy calculation shows

$$\sum_{k=a+1}^b (\nabla f)(k) = f(b) - f(a) \quad \text{and} \tag{2}$$

$$\nabla(f(k)g(k)) = (\nabla f)(k)g(k-1) + f(k)(\nabla g)(k). \tag{3}$$

Rearranging (3) and using (2) gives

$$\sum_{k=a+1}^b f(k)(\nabla g)(k) = f(b)g(b) - f(a)g(a) - \sum_{k=a+1}^b (\nabla f)(k)g(k-1). \tag{4}$$

Let  $\pi(x) = \sum_{p \leq x} 1$ . Observe that  $(\nabla \pi)(k)$  equals 1 if  $k$  is prime, and 0 otherwise. Thus,  $\sum_{a < p \leq b} f(p) = \sum_{a < k \leq b} f(k)(\nabla \pi)(k)$  and the latter can be estimated using (4). The following bounds hold for real  $x \geq 17$  by [16, Theorem 1]. In order to get a practically useful bounds in Theorem 10, we note that (5) holds for all integers  $x \geq 11$ :

$$\frac{x}{\log x} \leq \pi(x) \leq \frac{x}{\log x} \left( 1 + \frac{3}{2 \log x} \right) \quad \text{for } x \in \{11, 12, 13, \dots\}. \tag{5}$$

**Lemma 7.** Suppose  $a$  and  $b$  are real numbers with  $12 \leq a \leq b$ . Then

$$\sum_{a < p \leq b} \frac{1}{p^2} \leq \frac{2.22}{[a] \log [a]} - \frac{1.61}{[b] \log [b]}.$$

**Proof.** Since the primes in the range  $a < p \leq b$  coincide with the primes in the range  $[a] < p \leq [b]$ , we henceforth may (and will) assume that  $a$  and  $b$  are integers satisfying  $12 \leq a \leq b$ . Also, if  $a = b$  the inequality holds, so we assume also that  $a + 1 \leq b$ . Applying (4) and (5) gives

$$\begin{aligned} \sum_{a < p \leq b} \frac{1}{p^2} &= \sum_{k=a+1}^b \frac{1}{k^2} (\nabla \pi)(k) = \frac{\pi(b)}{b^2} - \frac{\pi(a)}{a^2} - \sum_{k=a+1}^b \frac{-2k+1}{k^2(k-1)^2} \pi(k-1) \\ &\leq \frac{1}{b \log b} \left( 1 + \frac{3}{2 \log b} \right) - \frac{1}{a \log a} \\ &\quad + \sum_{k=a+1}^b \frac{2k-1}{k^2(k-1) \log(k-1)} \left( 1 + \frac{3}{2 \log(k-1)} \right). \end{aligned}$$

Since  $12 \leq a \leq k-1$ , we have  $1 + 3/(2 \log(k-1)) \leq 1 + 3/(2 \log(12)) < 1.61$ . As  $k - \frac{1}{2} < k$ , the  $\Sigma$ -term above is less than

$$\begin{aligned} \frac{1.61}{\log a} \sum_{k=a+1}^b \frac{2(k-\frac{1}{2})}{k(k-\frac{1}{2})(k-1)} &= \frac{3.22}{\log a} \sum_{k=a+1}^b \frac{1}{k(k-1)} = \frac{3.22}{\log a} \sum_{k=a+1}^b \nabla \left( -\frac{1}{k} \right) \\ &= \frac{3.22}{\log a} \left( \frac{1}{a} - \frac{1}{b} \right) \\ &\leq \frac{3.22}{a \log a} - \frac{3.22}{b \log b}. \end{aligned}$$

Combining the previous two displays gives the desired inequality:

$$\sum_{a+1 < p < b} \frac{1}{p^2} < (-1 + 3.22) \frac{1}{a \log a} + (1.61 - 3.22) \frac{1}{b \log b} = \frac{2.22}{a \log a} - \frac{1.61}{b \log b}. \quad \square$$

**Lemma 8.** *Suppose  $a, b$  are real numbers with  $2 \leq a \leq b$  and  $p$  is prime. Then*

$$\log \left( \frac{\log b}{\log a} \right) - \frac{1}{2(\log b)^2} - \frac{1}{(\log a)^2} < \sum_{a < p \leq b} \frac{1}{p} < \log \left( \frac{\log b}{\log a} \right) + \frac{1}{(\log b)^2} + \frac{1}{2(\log a)^2}.$$

**Proof.** The proof uses the following bounds [16, Theorem 5, Corollary]

$$\log \log x + M - \frac{1}{2(\log x)^2} < \sum_{p \leq x} \frac{1}{p} < \log \log x + M + \frac{1}{(\log x)^2} \quad \text{for } x > 1,$$

where  $M = 0.26149 \dots$  is the Meissel–Mertens constant. Thus,

$$\begin{aligned} \sum_{a < p \leq b} \frac{1}{p} &= \sum_{p \leq b} \frac{1}{p} - \sum_{p \leq a} \frac{1}{p} \\ &< \left( \log \log b + M + \frac{1}{(\log b)^2} \right) - \left( \log \log a + M - \frac{1}{2(\log a)^2} \right) \\ &= \log \left( \frac{\log b}{\log a} \right) + \frac{1}{(\log b)^2} + \frac{1}{2(\log a)^2}. \end{aligned}$$

Similarly,

$$\begin{aligned} \sum_{a < p \leq b} \frac{1}{p} &= \sum_{p \leq b} \frac{1}{p} - \sum_{p \leq a} \frac{1}{p} \\ &> \log \log b + M - \frac{1}{2(\log b)^2} - \left( \log \log a + M + \frac{1}{(\log a)^2} \right) \\ &= \log \left( \frac{\log b}{\log a} \right) - \frac{1}{2(\log b)^2} - \frac{1}{(\log a)^2}. \quad \square \end{aligned}$$

### 5. Proof of Theorem 1

Suppose  $G \in \{A_n, S_n\}$ . Given certain functions  $a = a(n)$  and  $d = d(n)$ , our strategy is to find a lower bound for the proportion  $\rho_G$  of pre- $p$ -cycles in  $G$  with  $a(n) < p < a(n)^{d(n)}$ . It is shown in [18] that  $\rho_G \rightarrow 1$  as  $n \rightarrow \infty$ . Theorem 1 was motivated by the need to analyse certain Las-Vegas algorithms for permutation groups, we shall adapt a combinatorial argument in [18] to quantify this result. Henceforth,  $p$  denotes a prime, and  $k$  denotes an integer.



Fix  $G \in \{A_n, S_n\}$ . Recall  $\lambda(g) = \langle 1^{m_1(g)} 2^{m_2(g)} \dots n^{m_n(g)} \rangle$  is the partition of  $n$  whose parts are the cycle lengths of  $g$ , and part  $k$  has multiplicity  $m_k(g)$ . Define

$$\begin{aligned} \mathcal{P}_n &= \{p \mid a(n) < p \leq a(n)^{d(n)}\}, \\ T(G) &= \{g \in G \mid m_p(g) \geq 1 \text{ for some } p \in \mathcal{P}_n\}, \\ U_p(G) &= \left\{ g \in G \mid m_p(g) \geq 1 \text{ and } \sum_{k \geq 1} m_{kp}(g) \geq 2 \right\}, \text{ and} \\ U(G) &= \bigcup_{p \in \mathcal{P}_n} U_p(G). \end{aligned}$$

Note that  $g \in T(G) \setminus U_p(G)$  has  $m_p(g) \geq 1$  and  $\sum_{k \geq 1} m_{kp}(g) = 1$ . Hence  $m_p(g) = 1$  and  $g$  is a pre- $p$ -cycle by Lemma 3. Thus,  $T(G) \setminus U(G)$  is precisely the set of pre- $p$ -cycles in  $G$  for some  $p \in \mathcal{P}_n$ . In the following proposition, we view  $G \in \{A_n, S_n\}$  as a probability space with the uniform distribution, and we seek a lower bound for

$$\begin{aligned} \rho_G &= \text{Prob}(g \in G \text{ is a pre-}p\text{-cycle for some } p \in \mathcal{P}_n) = \text{Prob}(g \in T(G) \setminus U(G)) \\ &= \text{Prob}\left(g \in G \mid m_p(g) = 1 \text{ and } \sum_{k \geq 2} m_{kp}(g) = 0 \text{ for some } p \in \mathcal{P}_n\right). \end{aligned}$$

**Proposition 9.** *Let  $a(n), d(n)$  be functions satisfying  $a(n) \geq 12$ ,  $d(n) > 1$  and  $a(n)^{d(n)} \leq n$  for all  $n$ . Using the preceding notation and  $\delta = |S_n : G|$ , we have*

$$\rho_G = \frac{|T(G)| - |U(G)|}{|G|} \geq 1 - \frac{2.287\delta}{d(n)} - \frac{2.22(\log n - 1)}{[a(n)] \log[a(n)]} - \frac{4.4\delta \log n}{a(n)(\log a(n))n}.$$

**Proof.** Let  $\mu = \sum_{p \in \mathcal{P}_n} \frac{1}{p}$ . Write  $a$  and  $d$  instead of  $a(n)$  and  $d(n)$ , and set  $b = a^d$ . As  $a < p \leq b \leq n$ , we have  $\mathcal{P}_n \subseteq \{1, \dots, n\}$ . Thus, Proposition 6 gives

$$\frac{|T(G)|}{|G|} = 1 - \text{Prob}(g \in G \mid m_p(g) = 0 \text{ for all } p \in \mathcal{P}_n) \geq 1 - \delta e^{\gamma - \mu}. \tag{6}$$

We seek a lower bound for (6). This means finding a lower bound for  $\mu$ . Since  $a \geq 12$ , Lemma 8 gives

$$\mu = \sum_{a < p \leq a^d} \frac{1}{p} > \log d - \frac{1}{2(d \log a)^2} - \frac{1}{(\log a)^2} > \log d - 0.25. \tag{7}$$

Thus, it follows from (6) and (7) that

$$\frac{|T(G)|}{|G|} \geq 1 - \frac{\delta e^\gamma}{e^\mu} \geq 1 - \frac{\delta e^\gamma}{e^{\log d - 0.25}} \geq 1 - \frac{2.287\delta}{d}. \tag{8}$$

We seek an upper bound for  $|U_p(G)|$ . We (over)count the number of permutations  $g_1 g_2 g_3 \in U_p(G)$  where  $g_1, g_2, g_3$  have disjoint supports,  $g_1$  is  $p$ -cycle and  $g_2$  is a  $kp$ -cycle for some  $k \geq 1$ . We may choose  $g_1$  in  $\binom{n}{p}(p-1)!$  ways, because we take the first element to

be the smallest in the cycle. Next, we choose a  $kp$ -cycle in  $\binom{n-p}{kp}(kp-1)!$  ways. Observe that  $g_1$  is even:  $p$  is odd as  $p > \log 12 > 2$ . Thus, when  $G = A_n$  we must be able to choose  $g_3$  to have the same parity as  $g_2$ , so that  $g_1g_2g_3$  is even. In the generic case when  $(k+1)p \leq n-2$  this is possible. The number of choices for  $(g_1, g_2, g_3)$  in the generic case is

$$\binom{n}{p}(p-1)! \binom{n-p}{kp}(kp-1)! \frac{(n-p-kp)!}{\delta} = \frac{n!}{\delta kp^2}.$$

This is an upper bound for the number of products  $g_1g_2g_3$  in the generic case, and we may halve this upper bound if  $k = 1$ .

Consider the special case when  $n-1 \leq (k+1)p \leq n$  has a solution for  $k$  and  $p$ . Given  $p$ , this happens for at most one value of  $k$ , namely  $k = n/p - 1$  or  $k = (n-1)/p - 1$ , where  $p$  divides  $n$  or  $n-1$ , respectively. We will show that there are very few primes  $p$  for which  $(k+1)p$  lies in  $\{n-1, n\}$  for some  $k$ . Suppose that  $(k+1)p = n-1$  and  $n-1$  has  $r$  (not necessarily distinct) prime divisors greater than  $a$ . Then  $n-1 > a^r$  and hence there are at most  $r < \log_a(n-1)$  choices for  $p$ . Similarly, if  $(k+1)p = n$  there are less than  $\log_a(n)$  choices for  $p$ . Thus, the special case has less than  $2\log_a(n)$  choices for  $p$ . Consequently, the contribution in the special case is small, and our estimations need not be so careful.

In this special case, arguing as above, the number of choices of  $(g_1, g_2)$  is  $n!/(kp^2)$ , and the number  $n_3$  of choices for  $g_3$  is 1, unless  $G = A_n$  and  $k$  is even, in which case  $n_3 = 0$ . Thus, the number of products  $g_1g_2g_3$  is at most  $n!/(kp^2)$  and we bound the denominator as follows:

$$kp^2 > kap \geq \frac{ka(n-1)}{k+1} \geq \frac{a(n-1)}{2}.$$

In the special case, the number of choices for  $g_1g_2g_3$  is at most

$$\frac{n!}{kp^2} = \frac{\delta|G|}{kp^2} \leq \frac{2\delta|G|}{a(n-1)} = \varepsilon|G| \quad \text{where } \varepsilon = \frac{2\delta}{a(n-1)}.$$

For each prime  $p$ , let  $\varepsilon_p$  be  $\varepsilon$  if  $n-1 \leq (k+1)p \leq n$  has a solution for  $k$ , and 0 otherwise. In the generic case we have  $(k+1)p \leq n-2$  and  $k \leq m := \lfloor (n-2)/p \rfloor - 1$ . The bound  $\sum_{k=1}^m 1/k < 1 + \int_1^m dt/t = 1 + \log m$  is problematic if  $m = 0$ , so we replace  $m$  with  $m+1$  to get

$$\frac{|U_p(G)|}{|G|} \leq \left( \sum_{k=1}^{m+1} \frac{1}{kp^2} \right) + \varepsilon_p < \frac{1 + \log(m+1)}{p^2} + \varepsilon_p.$$

However,  $1 + \log(m+1) \leq 1 + \log(\lfloor n/p \rfloor) < \log n - 1$  as  $\log p > \log 12 > 2$ . Therefore

$$\frac{|U_p(G)|}{|G|} \leq \frac{\log n - 1}{p^2} + \varepsilon_p.$$

As  $\varepsilon_p = \varepsilon \neq 0$  for less than  $2\log_a(n)$  choices of  $p$ , we have

$$\frac{|U(G)|}{|G|} \leq \sum_{a < p \leq b} \frac{|U_p(G)|}{|G|} \leq (\log n - 1) \left( \sum_{a < p \leq b} \frac{1}{p^2} \right) + 2\varepsilon \log_a(n).$$

Applying the bound for  $\sum_{a < p \leq b} 1/p^2$  in Lemma 7 gives

$$\frac{|U(G)|}{|G|} \leq (\log n - 1) \left( \frac{2.22}{[a] \log[a]} - \frac{1.61}{[b] \log[b]} \right) + \frac{4\delta \log_a(n)}{a(n-1)}.$$

Since  $12 \leq a < n$  and  $4/(n-1) < 4.4/n$  holds for  $n \geq 13$ , we have

$$\frac{|U(G)|}{|G|} \leq (\log n - 1) \left( \frac{2.22}{[a] \log[a]} - \frac{1.61}{[b] \log[b]} \right) + \frac{4.4\delta \log_a(n)}{an}. \tag{9}$$

Now (8) and (9) give

$$\rho_G = \frac{|T(G)| - |U(G)|}{|G|} \geq 1 - \frac{2.287\delta}{d} - \frac{2.22(\log n - 1)}{[a] \log[a]} - \frac{4.4\delta \log n}{a(\log a)n}. \tag{10}$$

□

**Proof of Theorem 1.** Set  $a = \log n$  and  $d = \log \log n$ . Suppose that  $n \geq e^{12}$ . Then  $a \geq 12$ , and also  $a^d = (\log n)^{\log \log n} < n$ . Using Proposition 9 and the inequalities  $a - 1 < [a]$  and  $4.4/\log a < 2$  gives

$$\begin{aligned} \rho(G) &\geq 1 - \frac{2.287\delta}{d} - \frac{2.22(a-1)}{[a] \log[a]} - \frac{4.4\delta a}{a(\log a)n} \\ &> 1 - \frac{2.287\delta}{d} - \frac{2.22}{\log(a-1)} - \frac{2\delta}{n}. \end{aligned}$$

However,

$$\log(a-1) = \log a + \log\left(1 - \frac{1}{a}\right) \geq \log a + \log\left(\frac{11}{12}\right) \geq c' \log a,$$

where  $c' = 1 - \log(12/11)/\log(12) = \log(11)/\log(12)$  and  $2/n < 10^{-3}/\log \log n$ , so

$$\rho_G > 1 - \frac{2.287\delta}{\log \log n} - \frac{2.22}{c' \log \log n} - \frac{0.001\delta}{\log \log n} \geq 1 - \frac{2.288\delta}{\log \log n} - \frac{2.301}{\log \log n}.$$

This is at least  $1 - c/\log \log n$  where  $c = 4.6$  if  $\delta = 1$  and  $c = 6.9$  if  $\delta = 2$ . □

Allowing pre- $p$ -cycles with larger  $p$  gives us a sharper lower bound for  $\rho_G$ .

**Theorem 10.** Suppose that  $n \geq e^{12}$  and  $A_n \leq G \leq S_n$ . Let  $\rho_G$  be the proportion of permutations in  $G$  that power to a cycle with prime length  $p \leq n - 3$ . Then

$$\rho_G \geq 1 - \frac{(4.58\delta + 0.17) \log \log n}{\log(n-3)} \quad \text{where } \delta = \frac{n!}{|G|} \in \{1, 2\}.$$

**Proof.** Set  $a(n) = (\log n)^2$  in Proposition 9 and suppose that  $a(n)^{d(n)} = n - 3$ . (The hypotheses  $a(n) \geq 12$ ,  $d(n) > 1$  and  $a^d \leq n$  of Proposition 9 clearly hold.) Then  $d(n) = \log(n - 3)/\log a = \log(n - 3)/2 \log \log n$  and Proposition 9 gives

$$\begin{aligned} \rho_G &\geq 1 - \frac{2.287\delta}{d(n)} - \frac{2.22(\log n - 1)}{[(\log n)^2] \log[(\log n)^2]} - \frac{4.4\delta \log n}{a(\log a)n} \\ &> 1 - \frac{4.574\delta \log \log n}{\log(n - 3)} - \frac{2.22(\log n - 1)}{[(\log n)^2] \log[(\log n)^2]} - \frac{2.2\delta}{\log(n - 3)(\log \log n)n}. \end{aligned}$$

Note that  $[(\log n)^2] > (\log n)^2 - 1 = (\log n - 1)(\log n + 1)$  and since  $n \geq e^{12}$  we have

$$\frac{2.22(\log n - 1)}{[(\log n)^2]} < \frac{2.22}{\log n + 1} < \frac{2.05}{\log(n - 3)} \quad \text{and} \quad \frac{2.2\delta}{(\log \log n)n} \leq \frac{\delta \log \log n}{10^3}.$$

Using these inequalities, and combining the  $\delta$  terms, shows that

$$\rho_G \geq 1 - \frac{4.575\delta \log \log n}{\log(n - 3)} - \frac{2.05}{\log(n - 3) \log(12^2)} \geq 1 - \frac{(4.58\delta + 0.17) \log \log n}{\log(n - 3)}. \quad \square$$

**Remark 11.** Suppose that  $n \geq 5$ . We prove that the proportion  $\pi_n$  of elements of  $S_n$  that are pre- $p$ -cycles for some  $p$  with  $2 \leq p \leq n - 3$  is at least  $1/19$ . We know that  $\pi_n \geq \pi_0$  where  $\pi_0 := \sum_{n/2 < p \leq n-3} 1/p$ . A simple computation with MAGMA [3] shows that  $\pi_0 \geq 1/19$  for all  $n$  satisfying  $5 \leq n \leq 400,000$ . For  $n > 400,000 > e^{12}$  we have  $\pi_n \geq 1 - 4.75 \log \log n / \log(n - 3) > 1/19$  by Theorem 10. Precise computations of  $\pi_n$  for  $n \leq 50$  suggest that  $\pi_n > 1/3$  may even hold for all  $n \geq 5$ .

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